Reengineering the Booch Component Library

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Abstract

Component-based software construction heavily relies on the ability of reusing components from a library with as little effort as possible. Among others, valuable features for reusing from a component library are: adaptability to many contexts, extensibility, abstraction and high level of robustness with respect to changes in some of their components. In this paper we study one of the most widely used component library for Ada 95, the Booch's one, mainly in relation to these features. The hierarchy of the Booch component library is organised into three main super-classes: Containers, Support and Graphs, which have a common parent-class. Our study focuses on the Containers family, which present some drawbacks mainly due to the fact that some parent-classes depend on the concrete form of their children-classes. We state first the problems and we propose next a solution centred on changing the base class. This new version of the Containers class offers a new concept, namely shortcut, that allows not only to avoid the dependencies between parent-classes and their children classes, but also offers some additional advantages, remarkably improving the efficiency of components and increasing the possibilities of reusing from the library in contexts with strong efficiency requirements. Most of the ideas of this work can be generalised to other component libraries.

1 Introduction

Component-based software development [Jaz95, Sit94] heavily relies on the ability of reusing components from a library with as little effort as possible. Reusing from a library can only take place when the library fulfils some nice properties. To name a few:

- It should be versatile, offering a wide range of components with different functionalities, and also with the same functionality but different implementation strategies (each one suited for particular efficiency requirements).
- It should be open, in the sense that its users should be able to add easily other components to adapt the library to their own context.
- It should be stable, to avoid new future releases forcing changes on existing software.
- It should be abstract enough, to allow easy use.
- And of course, it should be correct and properly documented.

One of the best well-known component libraries for Ada 95 [Ada95] is the Grady Booch's one. This library was originally created for Ada 83 [Boo87] and reengineered first for C++ [BV90]

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and later on for Ada 95 [BWW99]. This library fulfills many of the properties listed above, and it has proven to be very useful in the development of component-based software. However, it also presents some drawbacks that make it not as powerful as it could be with respect to the following criteria:

- **Versatility.** It offers a wide range of components, but it is not as versatile as one could expect, because all its components have just a single implementation.

- **Efficiency.** Not only does the lack of many implementations provoke a loss of efficiency in some contexts, but some of these implementations are not efficient enough, because component interfaces do not allow direct access to certain parts of the implementation.

- **Openness.** New components and implementations for components could be added, but implementations for already existing components should look very similar to the existing ones, due to the internal structure of the library.

- **Stability.** Some feasible changes on implementations can damage existing programs, because component definitions make use in some places of features that are inherent of proper implementations.

- **Usage.** The mixture of specification and implementation characteristics in component definitions interferes with the easy understanding and use of the components therein.

The purpose of this paper is to make reengineering on the Booch library to solve these problems. The proposal is based mainly on changing a particular base class, *Containers*, in such a way that all the dependencies on a concrete implementation disappear. This new class allows easy integration in the library of new components and also of new arbitrary implementations for every component, even the already existing ones. Moreover, the new *Containers* class offers an enlargement in its interface that make its derived classes more efficient in a wider range of contexts. All the changes in the library do not affect existing programs, which are using it; it only needs a recompilation step. However, existing programs could be improved with this new version of *Containers* to be more efficient.

The rest of the paper is organised as follows. First, in Sect. 2, we analyse the Booch component library for Ada 95 describing the advantages and drawbacks that it presents. Then, in Sect. 3, we introduce the main features of a new *Containers* class that allows solving the problems that the Booch library presents. Section 4 explains the implementation details of the new *Containers* class. Next, in Sect. 5, we show the necessary modifications of the *Containers* children-classes and an example. Finally, in Sect. 6, we give the conclusions and future work.

### 2 Analysing the Booch component library for ADA 95

The Booch component library is one of the best-known component libraries for Ada 95 [Ada95]. This library was originally created for Ada 83 [Boo87] and reengineered first for C++ [BV90] and later on for Ada 95 [BWW99]. The Ada 95 version of the Booch components is organised into three main super-classes: *Containers*, *Support* and *Graphs*, which have a common parent-class *BC*. The base class *BC* has no functionality at all, it only provides the definition of the common exceptions. The *Containers* category of classes provides a wide range of structural abstractions (lists, bags, sets, collections, etc.) using many widespread
implementation techniques (chaining, hashing, search trees and so on). Figure 1 shows the main hierarchy of these components; their code is available at [BWW99].

Our study is centred on the `Container` family of classes, which presents some drawbacks that decrease the potential of reusability of this component library. Most of the problems arise because some parent-classes depend on the concrete implementation of their children-classes. To make it clearer, Fig. 2 shows a typical situation in this library, in which the `BC.Containers.Bags` class depends on the concrete form of its children classes `Bounded`, `Unbounded` and `Dynamic`, which are hashing tables. Notice that the type definition of `Bag_Iterator`\(^2\) forces all the `Bags` children to be implemented by means of a hashing table. This restriction interferes with the possibility of extending the class hierarchy or changing the concrete form of one of its children. All these dependencies exist because iterators are strongly dependent on the concrete container implementation. Similar dependencies can be also found between the classes `Maps`, `Sets`, `Queues`, etc., and their respective children. Therefore, to solve the problems we clearly need to make iterators independent of the specific container.

```verbatim
generic
package BC.Containers.Bags is
...
private

  type Bag is abstract new Container with null record;
...
  type Bag_Iterator (B : access Bag'Class)
  is new Actual_Iterator (B) with record
    Bucket_Index : Natural := 0;
    Index : Natural := 0;
  end record;
...
end BC.Containers.Bag;
```

Figure 2: Extract of the generic package `BC.Containers.Bags`.

\(^2\) An iterator is a separate abstract data type, found in most of the existing commercial libraries, that allows iterating through all the items in a structure. Its existence and good behaviour is critical for a library being competitive, and this is why all of the most important component libraries offer them in one way or another.
To sum up, the main problems in this library are:

- The hierarchy is not robust enough with respect to changes in some of their components: changes in a component require the modification of other components. This is due to the implementation dependency mentioned above. For instance, changing the current hashing implementation of the `BC.Containers.Bags.Bounded` by an implementation (bounded) with a binary search tree (for instance, because elements must be obtained in some order), the implementation of the type `Bag_Iterator`, defined in the `BC.Containers.Bags` class, must be changed, and the implementation of its operations as well.

- Moreover, this hierarchy restricts to a single set of possible implementations for some of the different structure abstractions. This is a serious drawback because some of the implementations provided therein can be inefficient in some contexts (we have already mentioned ordered traversal of `Bags`). For instance, it is not possible to have different implementations of the class `BC.Containers.Bags.Bounded`, because it is not implementation-independent, and hence it forces a concrete implementation strategy (hashing). This problem could be solved adding another level in the hierarchy, making the class abstract and defining their concrete children, but it is not possible without changing other parts of the hierarchy, because of the implementation dependency again.

- Low level of abstraction makes the usage of the implementation harder. This happens when dealing with iterators. The iterator type and its operations are strongly dependent on the concrete implementation of the underlying structure. As a consequence, for every concrete implementation of a children-class a new `Actual_Iterator` type must be defined and its operations must be overridden. This approach, which is different from many other libraries, prevents the easy usage of the iterator facility.

- It is not only the lack of multiple implementations for components that damage efficiency, but also some of the iterator operations have lineal cost in the worst case with respect to a certain parameter (although their amortised cost is constant). For instance, as shown in Fig. 3, the `Reset` operation could have linear cost with respect to the `Number_of_Buckets`. A similar problem occurs in the `Next` operation.

```plaintext
procedure Reset (It : in out Bag_Iterator) is
begin
  It.Index := 0;
  if Cardinality (It.B.all) = 0 then
    It.Bucket_Index := 0;
  else
    It.Bucket_Index := 1;
    while It.Bucket_Index <= Number_Of_Buckets (It.B.all) loop
      if Length (It.B.all, It.Bucket_Index) > 0 then
        It.Index := 1;
        exit;
      end if;
      It.Bucket_Index := It.Bucket_Index + 1;
    end loop;
  end if;
end Reset;
```

Figure 3: `Reset operation's code of the generic package BC.Containers.Bags.`
In many contexts in which components often encapsulate data structures, reusability can be damaged due to efficiency requirements: even if a component fulfils a required functionality, the time complexity of its operations may be inadequate given the context in which it should be integrated (either considering them individually or when combining them to build more complex components). The access by means of the operations offered by a component may be costly if the logical layout of the data structure is used; if fast access is required, it becomes necessary to look up the item using directly a reference to it.

It can be argued that these disadvantages are strong enough to rebound the use of the Booch library. However, this is not the case; this library offers several advantages that make it very useful, mainly:

- A large amount of robust and well-designed components with appropriate algorithms and data structures.
- It is a well-known library supported by documentation and books.
- It is freeware.
- It has several and complete testing packages for its components.

Given the above advantages we think it is worth improving this library by solving the above mentioned problems, instead of discarding it and looking for building a brand new one.

### 3 A new Containers class

The goal of this paper is to improve the containers-class family of Booch components library to avoid the problems shown in the previous section. Our approach consists on changing the base class `Containers` in such a way that the items will be stored in a generic container. Then iterators are implemented in this generic container and so they are independent of the concrete implementation of the children-classes. This generic container will offer efficient access paths to the items therein, which we call shortcuts. The `Containers` children-classes must store the shortcuts (that allow access to items) instead of the items themselves. In addition, the users can obtain an alternative, efficient and abstract way to access to the items stored in the structure.

The children classes of the new `Containers` class inherit from it the shortcuts and the operations. Therefore, the items stored therein can be accessed not only by using the operations given in the former specification, but also by means of other new ones that use the efficient paths. This results very useful whenever the operations are considered too expensive; on the other hand, since the addition of shortcuts does not affect the former behaviour of the existing operations, they can be used in the other cases as well.

Shortcuts solve the problems mentioned in the previous section:

- The items in the container can be accessed without knowing how they are stored in the container and, therefore, without knowing the underlying representation of the container (with arrays, pointers, linked, in tree-form, ...). Therefore, many implementations can exist together for the same type of container.
The access to the items in a container by means of a shortcut is achieved in constant time \( O(1) \), making it possible to reuse containers even with high efficiency constraints.

The addition of shortcuts to a container does not modify its functional behaviour. This is assured by incorporating the concept of shortcut into the formal specification of the container (see [FM00] for more details). Preservation of behaviour makes possible the substitution of old components by new ones.

Shortcuts are created and destroyed dynamically as items are inserted to and removed from a container; every item is bound to one (and only one) item. The shortcut bound to an item remains the same while it is inside the container, even if the underlying representation requires rearrangements. The access to the items by means of shortcuts is safe because meaningless access to them is avoided: our approach avoids dangling shortcuts or out-of-date ones.

4 Implementing the new Containers class with shortcuts

The essential point consists in implementing the generic container with a mapping from shortcuts to items. At the same time, items in the children components must be substituted by the shortcut that identifies them. The mapping from shortcuts may be implemented both using dynamic storage or an array; shortcuts are implemented then as pointers or cursors (i.e., array positions), respectively. Fields to obtain (in constant time) the shortcuts to the last item stored and to the first and the last item in the iterator order are necessary too; the field \( \text{first} \) to reset an iterator and the field \( \text{last} \) to add an item after the previous last item in the iterator order (which is not the same that the last item stored because it can have been added anywhere).

Released shortcuts must be available somehow to allow further reassignment, provided that there are not extra copies of it. When an item is removed from the container, it is marked as deleted. The corresponding shortcut can only be reused when there are no references to it. Therefore, memory management should be incorporated in the implementation. In this paper shortcuts are implemented by pointers reusing a particular memory manager offered by the Booch’s library itself.

To assure efficient and independent iterators we need to create a double-linked list of shortcuts, which means that we need \( 2 \times N \times \text{space(pointer)} + N \times \text{space(shortcut)} \) extra space (where \( N \) is the number of items in the container). Then, the iterators are implemented as shortcuts. The waste of this extra space offers a lot of benefits: the iterators are independent of the concrete form of the container, the efficiency of all the iterator operations is constant even in the worst case, and even this waste of space will generate a later saving, when shortcuts substitute identifiers (generally strings, which require more space than pointers or integers) in outer references from programs or other data structures.

Let’s now establish formally the equation that assures the saving of space. Let \( N \) be the total number of items in the container and \( R \) the total number of external references. Since generally,

\[
\text{space(item)} \geq \text{space(pointer)}
\]

then \( \exists k \geq 1 \text{ s.t. } \text{space(item)} \geq k \times \text{space(pointer)} \)
and since $space(\text{shortcut}) = 2 \times space(\text{pointer})$ (see below) space is really saved when the relationship

$$R \times space(\text{item}) \geq 4 \times N \times space(\text{pointer})$$

Holds, which is satisfied when the following relationship holds:

$$R \times k \geq 4 \times N.$$

Figure 4 shows the implementation of the types of the new class `BC.Containers`. The `Shortcut` type is implemented with a record with two fields: the first one is an access to the `Container` to which the shortcut is associated and the last one is a pointer (reused from `BC.Smart`) to the node where the item is stored. The `Node` type is a record with four fields: the first one is the item itself, the second and the third ones are pointers to maintain doubled-linked list of items and the last one is a boolean to mark the item as deleted when it has been logically removed but there are still some shortcuts that refer to it. The `Container` type is a record that maintains the corresponding fields for the last, first and last added items, and an additional field to obtain the number of items therein in constant time.

```ada
with BC.Smart;
generic
    type Item is private;
    with function "=" (L, R : Item) return Boolean is <>;
package BC.Containers is
    type Container is new Ada.Finalization.Controlled with private;
    type Shortcut is private;
    type Iterator is private;
    function "=" (L, R : Shortcut) return Boolean;
    ...
private
    type Node;
    type Access_Node is access Node'Class;
    type Access_Node_P is access Access_Node;
    package SP is new BC.Smart (T => Access_Node, P => Access_Node_P);
    type Shortcut_Pointer is new SP.Pointer;
    function Create_Node(C : Container) return Access_Node;
    type Access_Container is access all Container;
    type Shortcut is record
        Position: Shortcut_Pointer;
        For_The_Container: Access_Container := null;
        end record;
    type Node is tagged
        record
            Elem: Item;
            Next: Shortcut_Pointer;
            Previous: Shortcut_Pointer;
            Deleted: Boolean := FALSE;
            end record;
    type Iterator is new Shortcut;
    type Container is new Ada.Finalization.Controlled with record
        Cardinality: Natural := 0;
        First_Item: Shortcut_Pointer; -- For access to the first and the
        Last_Item: Shortcut_Pointer; -- last items by shortcuts.
        Last_Item_Added: Shortcut_Pointer;--To obtain the shortcut bound
        end record;
    end BC.Containers;
```

Figure 4: New `BC.Containers` class.
This scheme also works, without further considerations, in the case of the linear containers, such as stacks, lists, and so on. Figure 5, on the left, shows this situation with stacks. Now, the stack contains just shortcuts, and the objects are stored in a generic container, directly accessible by means of the shortcut. In the case of types accessed by some kind of key, like bags, maps, etc., we need to define the functions required for the concrete implementation over the shortcuts (for instance, the equality function and, in case of a hashing table, the hash function). These new functions will consist in applying the original function to the item associated to the shortcut. The scheme is shown in figure 5, on the right, and the Ada 95 code can be found in the appendix.

The new Containers class includes, in addition to the operations offered by the old one, seven new ones shown below. The first three ones are public while the others are private and, as a consequence, can only be used by the children classes (Fig. 6 shows the signature of these operations; see the appendix for a complete implementation). Their behaviour is as follows:

- **Add** adds the given item into the container in the last position (with respect to the iterator order). To be used in concrete containers implementations that do not require a special traversing order.

- **Add_After** adds the given item after the one bound to the given shortcut (with respect to the iterator order). If the shortcut is undefined, an exception arises. To use in concrete containers implementations that require a special traversing order.

- **Shortcut_To_The_Last_Item_Added** returns the shortcut that allows accessing to the last item added (using Add or Add_After) into the container. If there are no items in the container, it returns an undefined shortcut. This operation is required to obtain the shortcuts from the data structure; once obtained, it can be stored in other data structures and then coupling of structures without duplication of items is possible.

- **Item_Of** returns the item bound to the given shortcut. If the shortcut is undefined, an exception arises.

- **Defined** returns true if there is an item associated to the given shortcut, otherwise returns false.
• *Remove* removes from the container the item associated to the given shortcut. If the shortcut is undefined, an exception arises.

• *Node_Of* returns an access to the node where the item associated to the given shortcut is stored. If the shortcut is undefined, an exception arises. This operation is required in order to access to the values that have been stored together with the item, if it is the case. In the children classes, public efficient operations to access to these values by shortcuts can be defined. To allow this, the *Node* type is tagged and the children classes can define a new *Node* type with the necessary extra fields; in this case, the *Create_Node* function (see Fig. 4) has to be overridden.

```plaintext
... generic
... package BC.Containers is
... function Shortcut_To_The_Last_Item_Added
  (In_The_Container : Container) return Shortcut;
function Item_Of (The_Shortcut : Shortcut) return Item;
function Defined (The_Shortcut: Shortcut) return Boolean;
...
private
... procedure Add (In_The_Container : in out Container'Class; Elem: Item);
procedure Add_After (The_Shortcut: Shortcut;
  In_The_Container : in out Container;
  Elem: Item);
procedure Remove (The_Shortcut : Shortcut);
function Node_Of(S : Shortcut) return Access_Node;
...
end BC.Containers;
```

Figure 6: Signature of the new operations of the BC.Containers class.

We remark that the cost of all these operations as well as all the iterator operations is constant even in the worst case. Another important feature is that the operations *Shortcut_To_The_Last_Item_Added*, *Item_Of* and *Defined* are public and they are inherited by the children-classes of *Containers*. In consequence, all the children classes offer shortcuts and their users can take profit of them to improve the efficiency of their programs.

The complete implementation can be found not only in the appendix but also in ftp://ftp.lsi.upc.es/pub/users/jmarco/.

## 5 Using the new Containers class

In this section, we explain the steps required to modify the *Containers* children-classes to adapt them to the new *Containers* class. First, we explain the general changes required and then we applied the required changes to a concrete example, the class *BC.Containers.Bags.Bounded* and its ancestors.
In the general case, the changes required in the *Containers* children-classes are:

- Remove the specification and the implementation of the concrete iterators (i.e., the new `Actual_Iterator`).
- Modify the operations that take into account features that are tight to the concrete implementation of the new `Actual_Iterator`.
- In the class where the container is implemented we must:
  - Replace the items by shortcuts.
  - If there are elements to store together with the item (as in the example below), then define a new `Node` type and override the function `Create_Node` to initialise these elements.
  - If the existing implementation needs some specific functions (e.g., a hash function in the case of a container implemented by hashing), then redefine them working on shortcuts instead of items.
  - For every operation that works with items, add the indirection required to work with shortcuts instead of items.

We present below an example of the (few) changes to be done in the *Containers* children-classes. More precisely, we explain the required changes in the class `BC.Containers.Bags` and `BC.Containers.Bags.Bounded`, respectively. Similar changes should be done to the rest of children-classes.

The only required change in the package specification `BC.Containers.Bags` is to remove the iterators therein, i.e. the new `Actual_Iterator` type as well as its corresponding operations (Fig. 7 shows the part of the specification that must be removed).

```plaintext
... generic ...
package BC.Containers.Bags is ...
private ...
    type Bag_Iterator (B : access Bag'Class) is new Actual_Iterator (B) with record
        Bucket_Index : Natural := 0;
        Index : Natural := 0;
    end record;

    procedure Initialize (It : in out Bag_Iterator);
    -- Overriding primitive subprograms of the concrete actual Iterator.
    procedure Reset (It : in out Bag_Iterator);
    procedure Next (It : in out Bag_Iterator);
    function Is_Done (It : Bag_Iterator) return Boolean;
    function Current_Item (It : Bag_Iterator) return Item;
    function Current_Item (It : Bag_Iterator) return Item_Ptr;
end BC.Containers.Bags;
```

**Figure 7:** Types and operations of the `BC.Containers.Bags` class that must be removed.
The changes required in the package body are: removing the body of all the iterator operations and modifying the `Intersection` operation, because this operation takes into account that the elements in the hashing table can change their position when removals take place. Therefore the iterator to the next element remains the same. Figure 8 presents the new `Intersection` operation (the changes are marked with a comment `IS NEW`).

```pascal
... package body BC.Containers.Bags is ...
  procedure Intersection (B : in out Bag'Class; O : Bag'Class) is
    It : Iterator := New_Iterator (B);
    Aux : Iterator;
    begin
      -- XXX left out the optimisation which checks whether L, R are
      -- identical.
      while not Is_Done (It) loop
        declare
          This_Item : Item renames Current_Item (It);
          B_Count : Positive := Count (B, This_Item);
          begin
            if not Exists (O, This_Item) then
              Aux := It;                    -- IS NEW
              Next(Aux);                    -- IS NEW
              Detach (B, This_Item);
              It := Aux;                    -- IS NEW
            else
              declare
                O_Count : Positive := Count (O, This_Item);
                begin
                  if B_Count > O_Count then
                    Set_Value (B, This_Item, O_Count);
                  end if;
                end;
                Next (It);
              end if;
            end;
            Next (It);
          end loop;
        end; -- IS NEW
      end loop;
    end Intersection;
  end Intersection;
... end BC.Containers.Bags;
```

Figure 8: New version of `Intersection` operation of the BC.Containers.Bag.

The required changes of the `Containers.Bags.Bounded` are to create the hash table with shortcuts as items and as values (see Fig. 9) and to remove the specification of the functions `Item_At` and `Value_At` and of the type `Bounded_Bag_Iterator` in the specification package. Notice that we need to store the number of copies of an item (because a bag actually owns only one copy of each item and it counts the duplicates). We must then define a new `Node` type to store the number of occurrences together with the item. In consequence, the function `Create_Node` (introduced in Sect. 5) must be overridden. We must also define the function `"="` and the function `New_Hash` to take into account the necessary indirection provoked by storing the shortcuts instead of the items.
In the package body, we must add for every operation the corresponding indirection to access to the item besides of the corresponding implementation of Create_Node, "=" and New_Hash. Figure 10 shows the implementation of these three functions and, as an example of how indirections are added, the implementation of the operation Add. In the appendix, we present the whole new Containers class and the classes Containers.Bags and Containers.Bags.Bounded with all the required modifications.
... package body BC.Containers.Bags.Bounded is
...
function Create_Node(B : Bounded_Bag) return Access_Node is
A : Access_Node;
beginn
A := new Node_Bag;
Node_Bag(A.all).Value := 1;
return A;
end Create_Node;

function "=" (L, R : Shortcut) return Boolean is
begin
return Item_Of(L) = Item_Of(R);
end "=";

function New_Hash (S : Shortcut) return Positive is
begin
return Hash(Item_Of(S));
end New_Hash;
...

procedure Add (B : in out Bounded_Bag; I : Item; Added : out Boolean) is
sh : Shortcut;
A : Access_Node;
beginn
Containers.Add(B,I);
sh := Shortcut_To_The_Last_Item_Added(B);
if Tables.Is_Bound(B.Rep,sh) then
sh := Tables.Value_Of (B.Rep, sh);
A := Node_Of(sh);
Node_Bag(A.all).Value := Node_Bag(A.all).Value + 1;
Added := False;
Remove(Shortcut_To_The_Last_Item_Added(B));
B.Last_Item_Added := sh.Position;
else
Tables.Bind (B.Rep, sh, sh);
Added := True;
end if;
end Add;
...
end BC.Containers.Bags.Bounded;

Figure 10: Implementation of the functions Create_Node, "=". New_Hash and of the operation Add of BC.Containers.Bags.Bounded.

In the operation Add, we associate first a shortcut to the item, and then we ask if it is already bound to some item in the hash table. If it is already bound, we must increment the number of occurrences using the previous shortcut associated to the item and remove the temporal one.
6 Conclusions

Component libraries play currently a crucial role on software development. The existence of a wide variety of such libraries, both general-purpose and also field-specific, is becoming essential and its importance will increase in the future. However, nowadays a twofold phenomena occurs. On the one hand, most of the existing libraries present some features that can be improved, concerning quality factors as efficiency, usability, functionality, etc. On the other hand, potential clients of the library tend to simply discard it if it does not fit completely into their context. Our point of view is that both phenomena will be corrected little by little, when maturity in the field will be reached.

Our paper tries to be a contribution to this maturity process. We have presented here a case study of reengineering a good library, the Booch one, that presents a few drawbacks. First, we have enumerated the advantages and disadvantages it presents and then we have proposed a solution to all the drawbacks. The solution is based on designing a new Containers base class that avoids the problems and also offers a new type that implement the concept of shortcut as alternative access path to items in the container. Shortcuts are interesting because, besides of assuring fast access time to items in the container, they are abstract (independent of the implementation of the component), persistent (movements inside the data structure do not affect them), secure (meaningless accesses are not possible) and they preserve behaviour (the new component behaves as the old one). Therefore, both existing and new software can benefit from the nice properties that present the concept of shortcut. We have also shown the few necessary changes in the Containers children-classes to use the new Containers base class. The paper focuses in the technical details, to show the feasibility of the reengineering process and to make explicit the cost of this process.

The library with the modifications presented in this paper can be found in ftp://ftp.lsi.upc.es/pub/users/jmarco/.

We would like to point out the benefits of our approach:

- The Booch library is improved from many points of view: versatility, stability, openness, efficiency and ease of use.
- This improvement has been made in a very comfortable way; with only few changes needed. The core data structures and algorithms are the same without any modification at all.
- The concept of shortcut is general enough to be exported to other libraries. In fact, most of the main existing component libraries offer it in many different ways [MN99, MS96]. The advantage of our proposal is that it provides a systematic way of adding shortcuts to existing libraries, instead of putting them in the library from the very beginning.
- The reengineering process does not interfere with the previous behaviour of the library (both for functionality and for efficiency) and in consequence, existing software that use this library does not need to be modified; only recompilation is needed.
- But existing software could be modified in a methodical way (basically, changing the way of accessing to the structure) to take profit of the new version of the library. New
software, of course, will be built in general using the new layout of the structure, making
intensive use of shortcuts.

- The new library has been tested using the bag_test provided by the original Booch library
  without any modification.

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Conference (IRMA). To celebrate in May 2000.

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[Sit94] M. Sitaraman (coordinator). Special Feature: Component–Based Software using RESOLVE.
Appendix

Package specification of Containers

with Ada.Finalization;
with BC.Smart;
generic
  type Item is private;
  with function "=" (L, R : Item) return Boolean is <>;
package BC.Containers is

  -- This package specifies the common protocol of all Container classes.
  -- This common protocol consists of Iterators and Shortcuts.
  type Container is new Ada.Finalization.Controlled with private;
  -- Active iteration
  type Shortcut is private;
  type Iterator is private;

  function "=" (L, R : Shortcut) return Boolean;

  function Shortcut_To_The_Last_Item_Added
    (In_The_Container : Container) return Shortcut;

  function Item_Of (The_Shortcut : Shortcut) return Item;

  function Defined (The_Shortcut : Shortcut) return Boolean;

  function New_Iterator (For_The_Container : Container)
    return Iterator;

  procedure Reset (Obj : in out Iterator);
  -- Reset the Iterator to the beginning.

  procedure Reverse_Reset (Obj : in out Iterator);
  -- Reset the Iterator to the ending.

  procedure Next (Obj : in out Iterator);
  -- Advance the Iterator to the next Item in the Container.

  procedure Previous (Obj : in out Iterator);
  -- NEW NEW NEW

  function Is_Done (Obj : Iterator) return Boolean;
  -- Return True if there are no more Items in the Container.

  function Current_Item (Obj : Iterator) return Item;
  -- Return a copy of the current Item.

  generic
    with procedure Apply (Elem : in out Item);
    In_The_Iterator : in out Iterator;
    procedure Access_Current_Item;
  -- Call Apply for the Iterator's current Item.
-- Passive iteration

generic
with procedure Apply (Elem : in Item; OK : out Boolean);
procedure Visit (Using : in out Iterator);
-- Call Apply with a copy of each Item in the Container to which the
-- iterator Using is bound. The iteration will terminate early if Apply
-- sets OK to False.

generic
with procedure Apply (Elem : in out Item; OK : out Boolean);
procedure Modify (Using : in out Iterator);
-- Call Apply with a copy of each Item in the Container to which the
-- iterator Using is bound. The iteration will terminate early if Apply
-- sets OK to False.

procedure Add (In_The_Container : in out Container'Class;
   Elem: Item);

procedure Add_After (The_Shortcut: Shortcut;
   In_The_Container : in out Container;
   Elem: Item);

procedure Remove (In_The_Container : in out Container;
   The_Shortcut : Shortcut);
Undefined_Shortcut : exception;
Iterator_Is_Not_Bound : exception;
Iterator_Out_Of_Range : exception;

private

type Node;
type Access_Node is access Node'Class;
type Access_Node_P is access Access_Node;
package SP is new BC.Smart (T => Access_Node, P => Access_Node_P);
type Shortcut_Pointer is new SP.Pointer;
function Create_Node(C : Container) return Access_Node;
type Access_Container is access all Container;
type Shortcut is
record
   Position: Shortcut_Pointer;
   For_The_Container: Access_Container := null;
end record;
type Node is tagged
record
   Elem: Item;
   Next: Shortcut_Pointer;
   Previous: Shortcut_Pointer;
   Deleted: Boolean := FALSE;
end record;
type Iterator is new Shortcut;
type Container is new  Ada.Finalization.Controlled with
record
   Cardinality: Natural := 0;
   First_Item: Shortcut_Pointer;  -- For access to the first and the
   Last_Item: Shortcut_Pointer;   -- last items by shortcuts.
   Last_Item_Added: Shortcut_Pointer;-- To obtain the shortcut bound
end record;

procedure Finalize(C : in out Container);
procedure Adjust(C : in out Container);
function Cardinality (C : Container) return Natural;
procedure Purge (C : in out Container);
function Node_Of(S : Shortcut) return Access_Node;
end BC.Containers;
package body of Containers

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--
-- $Id: bc-containers.adb,v 1.6 1999/04/10 14:38:20 simon Exp $

with System.Address_To_Access_Conversions;

package body BC.Containers is

  function "=" (L, R : Shortcut) return Boolean is
  begin
    return L.Position = R.Position and L.For_The_container = R.For_The_container;
  end "=";

  function Create_Node(C : Container) return Access_Node is
  begin
    return new Node;
  end Create_Node;

  package Address_Conversions is new System.Address_To_Access_Conversions(Container);

  function Shortcut_To_The_Last_Item_Added (In_The_Container : Container) return
  Shortcut is
  begin
    p := Address_Conversions.Object_Pointer(Address_Conversions.To_Pointer(In_The_Container'Address));
    sh := Shortcut;
    begin
      sh.Position := In_The_Container.Last_Item_Added;
      sh.For_The_Container := Access_Container(p);
      return sh;
    end Shortcut_To_The_Last_Item_Added;

  procedure Add (In_The_Container : in out Container'Class; Elem: Item) is
  begin
    shp := Create(new Access_Node'(Create_Node(In_The_Container)));
    SP.Value(SP.Pointer(shp)).all.Elem := Elem;
    SP.Value(SP.Pointer(shp)).all.Previous := In_The_Container.Last_Item;
    SP.Value(SP.Pointer(shp)).all.Next := Dummy;
    SP.Value(SP.Pointer(shp)).all.Deleted := False;
    if In_The_Container.Cardinality = 0 then
      In_The_Container.First_Item := shp;
    else
      SP.Value(SP.Pointer(In_The_Container.Last_Item)).all.Next := shp;
    end if;
    In_The_Container.Last_Item := shp;
    In_The_Container.Last_Item_Added := shp;
    In_The_Container.Cardinality := In_The_Container.Cardinality + 1;
  end Add;
procedure Add_After (The_Shortcut: Shortcut; In_The_Container : in out Container;
Elem: Item) is
  shp : Shortcut_Pointer;
begin
  if not Defined(The_Shortcut) then
    raise Undefined_Shortcut;
  end if;
  if The_Shortcut.Position = In_The_Container.Last_Item then
    Add(In_The_Container,Elem);
  else
    shp := Create(new Access_Node'(Create_Node(In_The_Container)));
    SP.Value(SP.Pointer(shp)).all.Elem := Elem;
    SP.Value(SP.Pointer(shp)).all.Previous := The_Shortcut.Position;
    SP.Value(SP.Pointer(shp)).all.Next :=
      SP.Value(SP.Pointer(The_Shortcut.Position)).all.Next;
    SP.Value(SP.Pointer(shp)).all.Deleted := False;
    SP.Value(SP.Pointer(SP.Pointer(SP.Pointer(The_shortcut.Position)).all.Next)).all.Previous := shp;
    In_The_Container.Last_Item_Added := shp;
    In_The_Container.Cardinality := In_The_Container.Cardinality + 1;
  end if;
end Add_After;

procedure Remove (The_Shortcut : Shortcut) is
  Dummy : Shortcut_Pointer;
begin
  if not Defined(The_Shortcut) then
    raise Undefined_Shortcut;
  end if;
  if The_Shortcut.For_The_Container.First_Item = The_Shortcut.Position then
    The_Shortcut.For_The_Container.First_Item :=
      SP.Value(SP.Pointer(The_Shortcut.Position)).all.Next;
  end if;
  if The_Shortcut.For_The_Container.Last_Item = The_Shortcut.Position then
    The_Shortcut.For_The_Container.Last_Item :=
      SP.Value(SP.Pointer(The_Shortcut.Position)).all.Previous;
  end if;
  if The_Shortcut.For_The_Container.Last_Item_Added = The_Shortcut.Position then
    The_Shortcut.For_The_Container.Last_Item_Added :=
      SP.Value(SP.Pointer(The_Shortcut.Position)).all.Previous;
  end if;
  if SP.Value(SP.Pointer(SP.Pointer(SP.Pointer(The_shortcut.Position)).all.Next)) /= null then
    SP.Value(SP.Pointer(SP.Pointer(SP.Pointer(The_shortcut.Position)).all.Next)).all.Previous :=
      SP.Value(SP.Pointer(The_shortcut.Position)).all.Previous;
  end if;
  if SP.Value(SP.Pointer(SP.Pointer(SP.Pointer(The_shortcut.Position)).all.Previous)) /= null then
    SP.Value(SP.Pointer(SP.Pointer(SP.Pointer(The_shortcut.Position)).all.Previous)).all.Next :=
      SP.Value(SP.Pointer(The_shortcut.Position)).all.Next;
  end if;
  The_Shortcut.For_The_Container.all.Cardinality :=
    The_Shortcut.For_The_Container.all.Cardinality - 1;
  SP.Value(SP.Pointer(The_Shortcut.Position)).all.Deleted := True;
  SP.Value(SP.Pointer(The_Shortcut.Position)).all.Next := Dummy;
  SP.Value(SP.Pointer(The_Shortcut.Position)).all.Previous := Dummy;
  end Remove;

function Item_Of (The_Shortcut : Shortcut) return Item is
begin
  if not Defined(The_Shortcut) then
    raise Undefined_Shortcut;
  end if;
  return SP.Value(SP.Pointer(The_Shortcut.Position)).all.Elem;
end Item_Of;
function Defined (The_Shortcut: Shortcut) return Boolean is
begin
    if SP.Value(SP.Pointer(The_Shortcut.Position)) = null then
        return false;
    elsif The_Shortcut.For_The_Container = null then
        return false;
    else return not SP.Value(SP.Pointer(The_Shortcut.Position)).all.Deleted;
    end if;
end Defined;

function New_Iterator (For_The_Container : Container) return Iterator is
    it : Iterator;
    p : Address_Conversions.Object_Pointer :=
        Address_Conversions.To_Pointer (For_The_Container'Address);
begin
    it.Position := For_The_Container.First_Item;
    it.For_The_Container := Access_Container(p);
    return it;
end New_Iterator;

procedure Reset (Obj : in out Iterator) is
begin
    if Obj.For_The_Container = null then
        raise Iterator_Is_Not_Bound;
    end if;
    Obj.Position := Obj.For_The_Container.all.First_Item;
end Reset;

procedure Reverse_Reset (Obj : in out Iterator) is
begin
    if Obj.For_The_Container = null then
        raise Iterator_Is_Not_Bound;
    end if;
    Obj.Position := Obj.For_The_Container.all.Last_Item;
end Reverse_Reset;

procedure Next (Obj : in out Iterator) is
begin
    if Is_Done(Obj) then
        raise Iterator_Out_Of_Range;
    end if;
    Obj.Position := SP.Value (SP.Pointer (Obj.Position)).all.Next;
end Next;

procedure Previous (Obj : in out Iterator) is
begin
    if Is_Done(Obj) then
        raise Iterator_Out_Of_Range;
    end if;
    Obj.Position := SP.Value (SP.Pointer (Obj.Position)).all.Previous;
end Previous;

function Is_Done (Obj : Iterator) return Boolean is
begin
    if Obj.For_The_Container = null then
        raise Iterator_Is_Not_Bound;
    end if;
    return SP.Value (SP.Pointer (Obj.Position)).all.Deleted;
end Is_Done;

function Current_Item (Obj : Iterator) return Item is
begin
    if Is_Done(Obj) then
        raise Iterator_Out_Of_Range;
    end if;
    return Item_Of(Obj);
end Current_Item;

procedure Access_Current_Item is
begin
    Apply (SP.Value(SP.Pointer(In_The_Iterator.Position)).all.Elem);
end Access_Current_Item;
procedure Visit (Using : in out Iterator) is
  Success : Boolean;
begin
  Reset (Using);
  while not Is_Done (Using) loop
    Apply (Current_Item (Using), Success);
    exit when not Success;
  end loop;
end Visit;

procedure Modify (Using : in out Iterator) is
  Success : Boolean;
  procedure Caller (I : in out Item) is
    begin
      Apply (I, Success);
    end Caller;
  procedure Call_Apply is new Access_Current_Item (Caller, Using);
begin
  Reset (Using);
  while not Is_Done (Using) loop
    Call_Apply;
    exit when not Success;
  end loop;
end Modify;

function Cardinality (C : Container) return Natural is
begin
  return C.Cardinality;
end Cardinality;

procedure Purge (C : in out Container) is
  it1,it2 : Iterator := New_Iterator(C);
begin
  Reset(it1);
  while not Is_Done(it1) loop
    it2 := it1;
    Next(it2);
    Remove(Shortcut(it1));
    it1 := it2;
  end loop;
end Purge;

function Node_Of(S : Shortcut) return Access_Node is
begin
  return SP.Value(SP.Pointer(S.Position)).all;
end Node_Of;

procedure Finalize(C : in out Container) is
  it1,it2 : Iterator := New_Iterator(C);
begin
  Reset(it1);
  while not Is_Done(it1) loop
    it2 := it1;
    Next(it2);
    Remove(Shortcut(it1));
    it1 := it2;
  end loop;
end Finalize;

procedure Adjust(C : in out Container) is
  it : Iterator := New_Iterator(C);
  Dummy : Shortcut_Pointer;
begin
  Reset(it);
  C.Cardinality := 0;
  C.First_Item := Dummy;
  C.Last_Item := Dummy;
  C.Last_Item_Added := Dummy;
  while not Is_Done(it) loop
    Add(C,Current_Item(it));
    Next(it);
  end loop;
end Adjust;

end BC.Containers;
package specification of Containers.Bags

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-- $Id: bc-containers-bags.ads,v 1.2 1999/04/10 14:38:18 simon Exp $

generic
package BC.Containers.Bags is

-- A bag denotes a collection of items, drawn from some well-defined
-- universe. A bag may contain duplicate items. A bag actually owns only
-- one copy of each unique item: duplicates are counted, but are not
-- stored with the bag.
--
-- The parameter Item denotes the universe from which the bag draws its
-- items. Items may be a primitive type or user-defined.

type Bag is abstract new Container with private;

function Are_Equal (L, R : Bag'Class) return Boolean;
-- Return True if and only if both bags have the same number of distinct
-- items, and the same items themselves, each with the same count; return
-- False otherwise.
-- Can't call this "=" because of the standard one for Bag.

procedure Clear (B : in out Bag) is abstract;
-- Empty the bag of all items.

procedure Add (B : in out Bag; I : Item; Added : out Boolean) is abstract;
-- Add the item to the bag. If the item is not already a distinct member
-- of the bag, copy the item and add it to the bag and set Added to
-- True. If the item already exists, then increment the number of that
-- item and set Added to False.

procedure Add (B : in out Bag'Class; I : Item);
-- Add the item to the bag. If the item is not already a distinct member
-- of the bag, copy the item and add it to the bag; if it is, increment
-- the number of that item.

procedure Remove (B : in out Bag; I : Item) is abstract;
-- If the item is not a member of the bag, raise BC.Not_Found. Otherwise,
-- if there is exactly one of the item in the bag, remove the item in the
-- bag; if there is more than one of the item in the bag, simply decrement
-- its number.

procedure Union (B : in out Bag'Class; O : Bag'Class);
-- Perform a logical bag union; at the completion of this operation, the
-- bag B contains the items and counts found in its original state
-- combined with the bag O. For each item in the bag O, if the item is
-- not already a distinct member of the bag S, copy the item and add it
-- and its count to the bag S. If the item already is a member, increment
-- its count in S.

procedure Intersection (B : in out Bag'Class; O : Bag'Class);
-- Perform a logical bag intersection; at the completion of this
-- operation, the bag B contains the items found both in its original
-- state and in the bag O. For each item in the bag O, if the item is not
-- already a distinct member of the bag S, do nothing. If the item
-- already is a member of S, set its count to the lower of the two
-- counts. Items in the bag S but not in the bag O are also removed.
procedure Difference (B : in out Bag'Class; O : Bag'Class);
-- Perform a logical bag difference; at the completion of this operation,
-- the bag S contains the items found in its original state, less those
-- found in the bag O. For each item in the bag O, if the item is not
-- already a distinct member of the bag S, do nothing. If the item is a
-- member of the bag S with a count less than that in the bag O, remove
-- the item from the bag S. If the item is a member of the bag S with a
-- count more than that in the bag O, decrement the count in the bag S by
-- the count in the bag O.

function Extent (B : Bag) return Natural is abstract;
-- Return the number of distinct items in the bag.

function Total_Size (B : Bag'Class) return Natural;
-- Return the total number of items in the bag.

function Count (B : Bag; I : Item) return Natural is abstract;
-- Return the number of times the item occurs in the bag.

function Is_Empty (B : Bag) return Boolean is abstract;
-- Return True if and only if there are no items in the bag.

function Is_Member (B : Bag; I : Item) return Boolean is abstract;
-- Return True if and only if the item exists in the bag.

function Is_Subset (B : Bag'Class; O : Bag'Class) return Boolean;
-- Return True if and only if the bag B has the same or fewer distinct
-- items than in the bag O and equal or less numbers of each such item
-- than in the bag O.

function Is_Proper_Subset (B : Bag'Class; O : Bag'Class) return Boolean;
-- Return True if and only if
-- all the distinct items in the bag B are also in the bag O, and
-- either at least one of the items in the bag B has a lower number
-- than the number in the bag O,
-- or there is at least one distinct item in the bag O that is not
-- in the bag B.

private

type Bag is abstract new Container with null record;

procedure Attach (B : in out Bag; I : Item; C : Positive);
procedure Detach (B : in out Bag; I : Item);

procedure Set_Value (B : in out Bag; I : Item; C : Positive);
function Multiplicity (B : Bag'Class) return Natural;
function Number_Of_Buckets (B : Bag) return Natural;
function Length (B : Bag; Bucket : Positive) return Natural;

function Exists (B : Bag; I : Item) return Boolean;
function Item_At (B : Bag; Bucket, Index : Positive) return Item_Ptr;
function Value_At (B : Bag; Bucket, Index : Positive) return Positive;

type Bag_Iterator (B : access Bag'Class)
is new Actual_Iterator (B) with record
   Bucket_Index : Natural := 0;
   Index : Natural := 0;
end record;

procedure Initialize (It : in out Bag_Iterator);
-- Overriding primitive subprograms of the concrete actual Iterator.

procedure Reset (It : in out Bag_Iterator);
procedure Next (It : in out Bag_Iterator);
function Is_Done (It : Bag_Iterator) return Boolean;
function Current_Item (It : Bag_Iterator) return Item;
function Current_Item (It : Bag_Iterator) return Item_Ptr;
end BC.Containers.Bags;
package body of Containers.Bags

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-- $Id: bc-containers-bags.adb,v 1.2 1999/04/10 14:38:18 simon Exp $

with BC.Support.Exceptions;
with System.Address_To_Access_Conversions;

package body BC.Containers.Bags is

package BSE renames BC.Support.Exceptions;
procedure Assert
is new BSE.Assert ("BC.Containers.Bags");

function Are_Equal (L, R : Bag'Class) return Boolean is
  It : Iterator := New_Iterator (L);
begin
-- XXX left out the optimisation which checks whether L, R are
-- identical.
  if Cardinality (L) /= Cardinality (R) then
    return False;
  end if;
  while not Is_Done (It) loop
    if not Exists (R, Current_Item (It)) then
      return False;
    end if;
    if Count (L, Current_Item (It)) /= Count (R, Current_Item (It)) then
      return False;
    end if;
    Next (It);
  end loop;
  return True;
end Are_Equal;

procedure Add (B : in out Bag'Class; I : Item) is
  Dummy : Boolean;
begin
  Add (B, I, Added => Dummy);
end Add;

procedure Union (B : in out Bag'Class; O : Bag'Class) is
  It : Iterator := New_Iterator (O);
begin
-- XXX left out the optimisation which checks whether L, R are
-- identical.
  while not Is_Done (It) loop
    This_Item : Item renames Current_Item (It);
    This_Count : Positive := Count (O, This_Item);
    begin
      if not Is_Member (B, This_Item) then
        Attach (B, This_Item, This_Count);
      else
        Set_Value (B, This_Item, Count (B, This_Item) + This_Count);
      end if;
    end;
    Next (It);
  end loop;
end Union;
procedure Intersection (B : in out Bag'Class; O : Bag'Class) is
  It : Iterator := New_Iterator (B);
  Aux : Iterator;
begin
  -- XXX left out the optimisation which checks whether L, R are
  -- identical.
  while not Is_Done (It) loop
    declare
      This_Item : Item renames Current_Item (It);
      B_Count : Positive := Count (B, This_Item);
    begin
      if not Exists (O, This_Item) then
        Aux := It;    -- IS NEW
        Next (Aux);   -- IS NEW
        Detach (B, This_Item);
        It := Aux;    -- IS NEW
      else
        declare
          O_Count : Positive := Count (O, This_Item);
        begin
          if B_Count > O_Count then
            Set_Value (B, This_Item, O_Count);
          end if;
        end;
        Next (It);
      end if;
    end;
  end loop;
end Intersection;

procedure Difference (B : in out Bag'Class; O : Bag'Class) is
  It : Iterator := New_Iterator (O);
begin
  -- XXX left out the optimisation which checks whether L, R are
  -- identical.
  while not Is_Done (It) loop
    declare
      This_Item : Item renames Current_Item (It);
    begin
      if Exists (B, This_Item) then
        declare
          B_Count : Positive := Count (B, This_Item);
          O_Count : Positive := Count (O, This_Item);
        begin
          if B_Count <= O_Count then
            Detach (B, This_Item);
          else
            Set_Value (B, This_Item, B_Count - O_Count);
          end if;
        end;
        Next (It);
      end if;
    end;
  end loop;
end Difference;

function Total_Size (B : Bag'Class) return Natural is
  It : Iterator := New_Iterator (B);
  Result : Natural := 0;
begin
  while not Is_Done (It) loop
    Result := Result + Count (B, Current_Item (It));
    Next (It);
  end loop;
  return Result;
end Total_Size;
function Is_Subset (B : Bag'Class; O : Bag'Class) return Boolean is
    It : Iterator := New_Iterator (B);
begin
    -- XXX left out the optimisation which checks whether L, R are
    -- identical.
    if Cardinality (B) > Cardinality (O) then
        return False;
    end if;
    while not Is_Done (It) loop
        declare
            This_Item : Item := Current_Item (It);
        begin
            -- why don't I just do "or else Count (B, This_Item) > Count (O,
            -- This_Item)"? ... because it triggered a compiler bug in GNAT
            -- 3.11p (or was it 3.11b2?)
            if not Exists (O, This_Item) then
                return False;
            else
                declare
                    B_Count : Positive := Count (B, This_Item);
                    O_Count : Positive := Count (O, This_Item);
                begin
                    if B_Count > O_Count then
                        return False;
                    end if;
                end;
                end if;
            end;
        Next (It);
    end loop;
    return True;
end Is_Subset;

function Is_Proper_Subset (B : Bag'Class; O : Bag'Class) return Boolean is
    It : Iterator := New_Iterator (B);
    Is_Proper : Boolean := False;
begin
    -- XXX left out the optimisation which checks whether L, R are
    -- identical.
    if Cardinality (B) > Cardinality (O) then
        return False;
    end if;
    while not Is_Done (It) loop
        declare
            This_Item : Item := Current_Item (It);
        begin
            if not Exists (O, This_Item) then
                return False;
            else
                declare
                    B_Count : Positive := Count (B, This_Item);
                    O_Count : Positive := Count (O, This_Item);
                begin
                    if B_Count > O_Count then
                        return False;
                    elsif B_Count < O_Count then
                        Is_Proper := True;
                    end if;
                end;
                end if;
            end;
        Next (It);
    end loop;
    return Is_Proper or else Cardinality (B) < Cardinality (O);
end Is_Proper_Subset;

-- Subprograms to be overridden
procedure Attach (B : in out Bag; I : Item; C : Positive) is
begin
    raise Should_Have_Been_Overridden;
end Attach;

procedure Detach (B : in out Bag; I : Item) is
begin
    raise Should_Have_Been_Overridden;
end Detach;
procedure Set_Value (B : in out Bag; I : Item; C : Positive) is
begin
    raise Should_Have_Been_Overridden;
end Set_Value;

function Multiplicity (B : Bag'Class) return Natural is
    It : Iterator := New_Iterator (B);
    Result : Natural := 0;
begin
    while not Is_Done (It) loop
        Result := Result + Count (B, Current_Item (It));
        Next (It);
    end loop;
return Result;
end Multiplicity;

function Number_Of_Buckets (B : Bag) return Natural is
begin
    raise Should_Have_Been_Overridden;
return 0;
end Number_Of_Buckets;

function Length (B : Bag; Bucket : Positive) return Natural is
begin
    raise Should_Have_Been_Overridden;
return 0;
end Length;

function Exists (B : Bag; I : Item) return Boolean is
begin
    raise Should_Have_Been_Overridden;
return False;
end Exists;

function Value_Of (B : Bag; I : Item) return Positive is
begin
    raise Should_Have_Been_Overridden;
return 1;
end Value_Of;

-- function Item_At (B : Bag; Bucket, Index : Positive) return Item_Ptr is
-- begin
--    raise Should_Have_Been_Overridden;
--    return null;
-- end Item_At;

function Value_At (B : Bag; Bucket, Index : Positive) return Positive is
begin
    raise Should_Have_Been_Overridden;
return 1;
end Value_At;

end BC.Containers.Bags;
Package specification of Containers.Bags.Bounded

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--
-- $Id: bc-containers-bags-bounded.ads,v 1.2 1999/04/10 14:38:18 simon Exp $

with BC.Support.Bounded;
with BC.Support.Hash_Tables;
generic
    with function Hash (V : Item) return Positive is <>;
    Buckets : Positive;
    Size : Positive;
package BC.Containers.Bags.Bounded is
    -- A bag denotes a collection of items, drawn from some well-defined
    -- universe. A bag may contain duplicate items. A bag actually owns only
    -- one copy of each unique item: duplicates are counted, but are not
    -- stored with the bag.
    -- The hash function (the generic parameter Hash) determines the
    -- allocation of pairs to hash buckets. The value returned must not
    -- change during the lifetime of a given Item. The range of hash values
    -- need not be constrained to the number of buckets in the bag.
    -- The hash function must satisfy the condition that, for objects A and
    -- B, if A = B, then Hash (A) must equal Hash (B). The hash function
    -- should attempt to spread the set of possible items uniformly across
    -- the number of buckets. The quality of the hash function has a
    -- significant impact upon performance.

type Bounded_Bag is new Bag with private;

procedure Clear (B : in out Bounded_Bag);
    -- Empty the bag of all items.

procedure Add (B : in out Bounded_Bag; I : Item; Added : out Boolean);
    -- Add the item to the bag. If the item is not already a distinct member
    -- of the bag, copy the item and add it to the bag and set Added to
    -- True. If the item already exists, then increment the number of that
    -- item and set Added to False.

function Available (B : Bounded_Bag) return Natural;
    -- Return the number of unused slots in the bag. Note, since hash buckets
    -- are of fixed size in the bounded bag it will probably not be possible
    -- to use all these slots.

procedure Remove (B : in out Bounded_Bag; I : Item);
    -- If the item is not a member of the bag, raise BC.Not_Found. Otherwise,
    -- if there is exactly one of the item in the bag, remove the item in the
    -- bag; if there is more than one of the item in the bag, simply decrement
    -- its number.

function Extent (B : Bounded_Bag) return Natural;
    -- Return the number of distinct items in the bag.

function Count (B : Bounded_Bag; I : Item) return Natural;
    -- Return the number of times the item occurs in the bag.

function Is_Empty (B : Bounded_Bag) return Boolean;
    -- Return True if and only if there are no items in the bag.
function Is_Member (B : Bounded_Bag; I : Item) return Boolean;
-- Return True if and only if the item exists in the bag.

private

function "=" (L, R : Shortcut) return Boolean;

type Shortcut_Ptr is access all Shortcut;
-- for Shortcut_Ptr'Storage_Size use 0;
package IC is new BC.Support.Bounded (Item => Shortcut,
Item_Ptr => Shortcut_Ptr,
Maximum_Size => Size);

use IC;

package VC is new BC.Support.Bounded (Item => Shortcut,
Item_Ptr => Shortcut_Ptr,
Maximum_Size => Size);

use VC;

type Node_Bag is new Node with
record
  Value : Positive;
end record;
function Create_Node(B: Bounded_Bag) return Access_Node;

function New_Hash(S : Shortcut) return Positive;

package Tables is new BC.Support.Hash_Tables
(Item => Shortcut,
Hash => New_Hash,
Value => Shortcut,
Value_Ptr => Shortcut_Ptr,
Buckets => Buckets,
Item_Container => IC.Bnd_Node,
Item_Container_Ptr => IC.Bnd_Node_Ref,
Value_Container => VC.Bnd_Node,
Value_Container_Ptr => VC.Bnd_Node_Ref);

type Bounded_Bag is new Bag with record
  Rep : Tables.Table;
end record;

procedure Adjust(B : in out Bounded_Bag);
procedure Purge (B : in out Bounded_Bag);

procedure Attach (B : in out Bounded_Bag; I : Item; C : Positive);

procedure Detach (B : in out Bounded_Bag; I : Item);

procedure Set_Value (B : in out Bounded_Bag; I : Item; C : Positive);
-- function Cardinality (B : Bounded_Bag) return Natural;

function Number_Of_Buckets (B : Bounded_Bag) return Natural;

function Length (B : Bounded_Bag; Bucket : Positive) return Natural;

function Exists (B : Bounded_Bag; I : Item) return Boolean;

end BC.Containers.Bags.Bounded;
package body of Containers.Bags.Bounded

with BC.Support.Exceptions;
with System.Address_To_Access_Conversions;
--with Ada.Text_IO; use Ada.Text_IO;

package body BC.Containers.Bags.Bounded is

    package BSE renames BC.Support.Exceptions;
    procedure Assert
        is new BSE.Assert ("BC.Containers.Bags.Bounded");

    function Create_Node (B : Bounded_Bag) return Access_Node is
        A : Access_Node;
        begin
            A := new Node_Bag;
            Node_Bag(A.all).Value := 1;
            return A;
        end Create_Node;

    function "=" (L, R : Shortcut) return Boolean is
        begin
            return Item_Of(L) = Item_Of(R);
        end "=";

    function New_Hash (S : Shortcut) return Positive is
        begin
            return Hash(Item_Of(S));
        end New_Hash;

    procedure Clear (B : in out Bounded_Bag) is
        begin
            Purge(Container(B));
            Tables.Clear (B.Rep);
        end Clear;

    procedure Add (B : in out Bounded_Bag; I : Item; Added : out Boolean) is
        sh : Shortcut;
        A : Access_Node;
        begin
            Containers.Add(B,I);
            sh := Shortcut_To_The_Last_Item_Added(B);
            if Tables.Is_Bound(B.Rep,sh) then
                sh := Tables.Value_Of (B.Rep, sh);
                A := Node_Of(sh);
                Node_Bag(A.all).Value := Node_Bag(A.all).Value + 1;
                Added := False;
                Remove(Shortcut_To_The_Last_Item_Added(B));
                B.Last_Item_Added := sh.Position;
            else
                Tables.Bind (B.Rep, sh, sh);
                Added := True;
            end if;
        end Add;

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--
-- $Id: bc-containers-bags-bounded.adb,v 1.2 1999/04/10 14:38:18 simon Exp $
procedure Remove (B : in out Bounded_Bag; I : Item) is
  sh1, sh2 : Shortcut;
  A : Access_Node;
  Caux : Container;
begin
  Containers.Add(Caux,I);
  sh1 := Shortcut_To_The_Last_Item_Added(Caux);
  Assert (Tables.Is_Bound (B.Rep, sh1),
          BC.Not_Found'Identity,
          "Remove", BS2.Missing);
  sh2 := Tables.Value_Of (B.Rep, sh1);
  Remove(sh1);
  A := Node_Of(sh2);
  if Node_Bag(A.all).Value = 1 then
    Tables.Unbind (B.Rep, sh2);
    Remove(sh2);
  else
    Node_Bag(A.all).Value := Node_Bag(A.all).Value - 1;
  end if;
end Remove;

function Available (B : Bounded_Bag) return Natural is
  Count : Natural := 0;
begin
  for P in 1 .. Buckets loop
    Count := Count + IC.Available (Tables.Item_Bucket (B.Rep, P).all);
  end loop;
  return Count;
end Available;

function Extent (B : Bounded_Bag) return Natural is
begin
  return Tables.Extent (B.Rep);
end Extent;

function Count (B : Bounded_Bag; I : Item) return Natural is
  sh : Shortcut;
  Co : Natural;
  Caux : Container;
begin
  Containers.Add(Caux,I);
  sh := Shortcut_To_The_Last_Item_Added(Caux);
  if not Tables.Is_Bound (B.Rep, sh) then
    Co := 0;
  else
    Co := Node_Bag(Node_Of(Tables.Value_Of (B.Rep, sh)).all).Value;
  end if;
  Remove(sh);
  return Co;
end Count;

function Is_Empty (B : Bounded_Bag) return Boolean is
begin
  return Tables.Extent (B.Rep) = 0;
end Is_Empty;

function Is_Member (B : Bounded_Bag; I : Item) return Boolean is
  Is_M : Boolean;
  sh : Shortcut;
  Caux : Container;
begin
  Containers.Add(Caux,I);
  sh := Shortcut_To_The_Last_Item_Added(Caux);
  Is_M := Tables.Is_Bound (B.Rep, sh);
  Remove(sh);
  return Is_M;
end Is_Member;

-- Private implementations

procedure Purge (B : in out Bounded_Bag) is
begin
  Purge(Container(B));
  Tables.Clear (B.Rep);
end Purge;
procedure Attach (B : in out Bounded_Bag; I : Item; C : Positive) is
sh : Shortcut;
begin
  Containers.Add(B,I);
  sh := Shortcut_To_The_Last_Item_Added(B);
  Node_Bag(Node_Of(sh).all).Value := C;
  Tables.Bind (B.Rep, sh, sh);
end Attach;

procedure Detach (B : in out Bounded_Bag; I : Item) is
sh1, sh2 : Shortcut;
begin
  Containers.Add(Baux,I);
  sh1 := Shortcut_To_The_Last_Item_Added(Caux);
  sh2 := Tables.Value_Of (B.Rep, sh1);
  Tables.Unbind (B.Rep, sh2);
  Remove(sh1);
  Remove(sh2);
end Detach;

procedure Set_Value (B : in out Bounded_Bag; I : Item; C : Positive) is
sh1, sh2 : Shortcut;
begin
  Containers.Add(Caux,I);
  sh1 := Shortcut_To_The_Last_Item_Added(Caux);
  sh2 := Tables.Value_Of (B.Rep, sh1);
  Node_Bag(Node_Of(sh2).all).Value := C;
  Remove(sh1);
end Set_Value;

--  function Cardinality (B : Bounded_Bag) return Natural is
--  begin
--    return Tables.Extent (B.Rep);
--  end Cardinality;

function Number_Of_Buckets (B : Bounded_Bag) return Natural is
begin
  return Buckets;
end Number_Of_Buckets;

function Length (B : Bounded_Bag; Bucket : Positive) return Natural is
begin
  return IC.Length (Tables.Item_Bucket (B.Rep, Bucket).all);
end Length;

function Exists (B : Bounded_Bag; I : Item) return Boolean is
Ex : Boolean;
sh : Shortcut;
begin
  Containers.Add(Caux,I);
  sh := Shortcut_To_The_Last_Item_Added(Caux);
  Ex := Tables.Is_Bound (B.Rep, sh);
  Remove(sh);
  return Ex;
end Exists;

procedure Adjust(B : in out Bounded_Bag) is
it1, it2 : Iterator;
begin
  it1 := New_Iterator(B);
  Reset(it1);
  Containers.Adjust(Container(B));
  it2 := New_Iterator(B);
  Reset(it2);
  while not Is_Done(it2) loop
    Tables.Bind (B.Rep, Shortcut(it2), Shortcut(it2));
    Node_Bag(Node_Of(Shortcut(it2)).all).Value :=
      Node_Bag(Node_Of(Shortcut(it1)).all).Value;
    Next(it1);
    Next(it2);
  end loop;
end Adjust;
end BC.Containers.Bags.Bounded;