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Intersection problem for Droms RAAGs

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Abstract

We solve the subgroup intersection problem (SIP) for any RAAG G of Droms type (i.e., with defining graph not containing induced squares or paths of length 3): there is an algorithm which, given finite sets of generators for two subgroups $H, K \leq G$, decides whether $H \cap K$ is finitely generated or not, and, in the affirmative case, it computes a set of generators for $H \cap K$. Taking advantage of the recursive characterization of Droms groups, the proof consists in separately showing that the solvability of SIP passes through free products, and through direct products with free-abelian groups. We note that most of RAAGs are not Howson, and many (e.g. $\mathbb{F}_2 \times \mathbb{F}_2$) even have unsolvable SIP.

1 Introduction

In group theory, the study of intersections of subgroups has been recurrently considered in the literature. Roughly speaking, the problem is “*given subgroups $H, K \leq G$, find $H \cap K$* ”. However, in the context of Geometric Group Theory, where groups may be infinite, or even non finitely generated, one needs to be more precise about the word *find*, specially if one is interested in the computational point of view.

A group is said to satisfy *Howson’s property* — or to be *Howson*, for short — if the intersection of any two (and so, finitely many) finitely generated subgroups is again finitely generated.

Classical examples of Howson groups include free-abelian, and free groups. In \mathbb{Z}^m Howson’s property is trivial, whereas for free groups it was proved by Howson himself in [15], where he also gave an algorithm to compute generators for the intersection.

Not far from these groups one can find examples without the Howson property: consider the group $\mathbb{F}_2 \times \mathbb{Z} = \langle a, b \mid - \rangle \times \langle t \mid - \rangle$ and the subgroups $H = \langle a, b \rangle$ and $K = \langle ta, b \rangle$; both

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are clearly 2-generated but $H \cap K = \{w(a, b) : |w|_a = 0\} = \langle\langle b \rangle\rangle \leq \mathbb{F}_2$, which is not finitely generated. In this context, it is natural to consider the following decision problems.

Subgroup intersection problem, SIP(G). *Given words $u_1, \dots, u_n, v_1, \dots, v_m$ in the generators of G , decide whether the subgroup intersection $\langle u_1, \dots, u_n \rangle \cap \langle v_1, \dots, v_m \rangle$ is finitely generated or not; and, in the affirmative case, compute a generating set for this intersection.*

Coset intersection problem, CIP(G). *Given a finite set of words $w, w', u_1, \dots, u_n, v_1, \dots, v_m$ in the generators of G , decide whether the coset intersection $w\langle u_1, \dots, u_n \rangle \cap w'\langle v_1, \dots, v_m \rangle$ is empty or not; and in the negative case, compute a coset representative.*

In [8], Delgado–Ventura prove that direct products of free-abelian and free groups have both SIP and CIP solvable. The goal of the present paper is to extend the algebraic arguments given there, in order to achieve similar properties for a much wider family of groups. To this end it is convenient to consider the following variations for a general finitely presented group G .

Twofold intersection problem, TIP(G). *Solve both SIP(G) and CIP(G).*

Extended subgroup intersection problem, ESIP(G). *Given a finite set of words $u_1, \dots, u_n, v_1, \dots, v_m, w, w'$ in the generators of G , decide whether the intersection of the subgroups $H = \langle u_1, \dots, u_n \rangle$ and $K = \langle v_1, \dots, v_m \rangle$ is finitely generated or not; and in the affirmative case: (i) compute a generating set for $H \cap K$, and (ii) decide whether the coset intersection $wH \cap w'K$ is empty or not (denoted by CIP_{fg}), computing a coset representative in case it is not.*

The main result in this paper is about *finitely generated* PC-groups (a.k.a. right-angled Artin groups, or RAAGs). This prominent class of groups is closely linked to some crucial examples of groups, notably Bestvina and Brady’s example of a group which is homologically finite (of type FP) but not geometrically finite (in fact not of type F_2), and Mihailova’s example of a group with unsolvable subgroup membership problem. More recently D.Wise et al. developed a method of showing that a vast amount of groups are virtual subgroups of RAAGs. Wise used this method to solve some well-known problems in group theory, like Baumslag’s conjecture on residual finiteness of one-relator groups with torsion. Furthermore, building on the work of Wise and Kahn–Markovic, I. Agol proved the famous virtually fibred conjecture (the last main open problem in 3-manifold theory due to Thurston), by showing that fundamental groups of closed, irreducible, hyperbolic 3-manifolds are virtual subgroups of RAAGs.

Definition 1.1. A group G is said to be *partially commutative* (a PC-group, for short) if it admits a presentation of the form

$$\langle X \mid [x_i, x_j] = 1, \text{ whenever } \{x_i, x_j\} \in E \rangle, \quad (1)$$

for some (not necessarily finite) simple graph $\Gamma = (X, E)$. In this case, we say that G is presented by the *commutation graph* Γ and write $G = G_\Gamma$. Then, we say that (1) is a *graphical presentation* for G , and X is a *graphical generating set* (or *basis*) for G . In the f.g. case (i.e., when X is finite) we shall refer to G_Γ as a *right-angled Artin group* (a RAAG, for short).

A subgraph Δ of a graph $\Gamma = (X, E)$ is said to be *full* if it has exactly the edges that appear in Γ over its vertex set, say $Y \subseteq X$; in this case, Δ is called the *full subgraph of Γ spanned by Y* and we write $\Delta \leq \Gamma$. When none of the graphs belonging to a certain family \mathcal{F} appear as a full subgraph of Γ , we say that Γ is *\mathcal{F} -free*.

PC-groups can be thought as a family of groups interpolating between two extreme cases: free-abelian groups (presented by complete graphs), and free groups (presented by edgeless graphs); having as *graphical generating set* precisely the standard free-abelian and free bases, respectively. More generally, the PC-group presented by the disjoint union of graphs $\Gamma \sqcup \Delta$ is the free product $G_\Gamma * G_\Delta$, and the PC-group presented by the join of graphs $\Gamma \vee \Delta$ (obtained by adding to $\Gamma \sqcup \Delta$ every edge joining a vertex in Γ to a vertex in Δ) is the direct product $G_\Gamma \times G_\Delta$.

Despite the extreme (free and free-abelian) cases being subgroup-closed, this is not the case for PC-groups. Droms characterized the finitely generated PC-groups having this property in the following well known result.

Theorem 1.2 (C. Droms, 1987, [11]). *Let Γ be a finite graph. Then, every subgroup of G_Γ is again a (possibly non finitely generated) PC-group if and only if Γ is $\{P_4, C_4\}$ -free.* \square

Here, P_n stands for the *path graph* on n vertices; and C_n stands for the *cycle graph* on n vertices. Accordingly, we say that a graph is a *Droms graph* if it is finite and $\{P_4, C_4\}$ -free, and a PC-group is a *Droms group* if it is presented by a Droms graph.

Remark 1.3. We note that finite $\{P_4, C_4\}$ -free graphs have received diverse denominations throughout the literature, including *comparability graphs of forests* (in [29]), *transitive forests* (in [26]), *trivially perfect graphs* (in [4]), and *quasi-threshold graphs* (in [19]).

1.1 Results

The main result in the present paper is the following theorem.

Theorem 1.4. *Every Droms group has solvable ESIP (and, in particular, solvable SIP).* \square

The strategy of the proof arises from the following crucial lemma given by Droms on the way of proving Theorem 1.2: *Every nonempty Droms graph is either disconnected, or it contains a central vertex (i.e., one vertex adjacent to any other vertex).*

This easily provides the following recursive definition of the Droms family (with both the graphical and the algebraic counterparts):

Corollary 1.5 (Droms, [11]). *The family of Droms graphs (resp., Droms groups) can be recursively defined as the smallest family \mathcal{D} (resp., \mathcal{D}) satisfying the following rules:*

- | | |
|--|--|
| [D1] $K_0 \in \mathcal{D}$; | [D1] $\{1\} \in \mathcal{D}$; |
| [D2] $\Gamma_1, \Gamma_2 \in \mathcal{D} \Rightarrow \Gamma_1 \sqcup \Gamma_2 \in \mathcal{D}$; | [D2] $G_1, G_2 \in \mathcal{D} \Rightarrow G_1 * G_2 \in \mathcal{D}$; |
| [D3] $\Gamma \in \mathcal{D} \Rightarrow K_1 \vee \Gamma \in \mathcal{D}$. | [D3] $G \in \mathcal{D} \Rightarrow \mathbb{Z} \times G \in \mathcal{D}$. \square |

Our proof of Theorem 1.4 is based on the following preservability results for the intersection properties we are interested in.

Theorem 1.6. *Let G be a Droms group. If G has solvable SIP, then $\mathbb{Z}^m \times G$ also has solvable SIP.* \square

Theorem 1.7. *Let G be a Droms group. If G has solvable ESIP, then $\mathbb{Z}^m \times G$ also has solvable ESIP.* \square

Theorem 1.8. *If two finitely presented groups G_1 and G_2 have solvable ESIP, then their free product $G_1 * G_2$ also has solvable ESIP.* \square

Theorem 1.9. *If two finitely presented groups G_1 and G_2 have solvable TIP, then their free product $G_1 * G_2$ also has solvable TIP.* \square

To prove Theorems 1.6 and 1.7 we extend the techniques in [8] from free groups to RAAGs; and to prove Theorems 1.8 and 1.9 we use Ivanov’s techniques to understand and work with subgroups of free products (see [16, 17]). Both are relatively long and technical arguments, each requiring several pages of preliminary considerations. For the sake of clarity, we decided to include them instead of writing a shorter preprint but harder to read.

Our main result (Theorem 1.4) can be seen as a partial generalization of Theorem 2.7(i) by Kapovich–Weidmann–Myasnikov in the sense that we prove a stronger thesis than them (namely, ESIP instead of MP, see Figure 1) for a smaller class of groups (Droms instead of coherent PC-groups). In this situation, it is interesting to ask the following questions.

Question 1. *Does the group \mathbb{G}_{P_4} have solvable SIP? Is it true that a RAAG have solvable SIP if and only if it is Droms?*

Particularly suggestive for us is the result from Aalbersberg–Hoogeboom [1] stating that the intersection problem for a partially commutative monoid is solvable if and only if its commutation graph is Droms. The situation is intriguingly similar to that for the MP.

The paper is organized as follows. In Section 2 we establish the necessary background and references about Droms groups and algorithmic issues, and we prove the main result, Theorem 1.4, modulo Theorems 1.6, 1.7, 1.9, and 1.8. Then, in Section 3, we study the direct product case, proving Theorems 1.6 and 1.7; and finally, in Section 4, we consider the free product situation proving Theorems 1.9 and 1.8. See [6, Part III] for a more detailed version of these results.

2 Preliminaries

Below we present the necessary preliminaries on algorithmicity and PC-groups.

2.1 Algorithmic aspects

Similar preserving properties concerning free and direct products were studied for the Membership Problem (MP) by K. A. Mikhailova. In [25] she proved that MP is preserved

under free products; whereas in [24], she showed that $\mathbb{F}_2 \times \mathbb{F}_2$ has unsolvable membership problem, proving that MP (and thus SIP and CIP) *do not* pass to direct products.

Several obvious relations among the already introduced algorithmic problems are summarized in the diagram below:

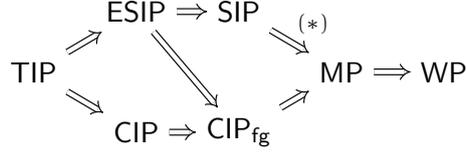


Figure 1: Some dependencies between algorithmic problems

where the starred implication is true with the extra assumption that the involved group is torsion-free, and has solvable word problem. We recall that WP and MP stand for the classical *word problem*, and the *subgroup membership problem* stated below.

(Subgroup) membership problem, $MP(G)$. *Given a finite set of words w, u_1, \dots, u_n in the generators of G , decide whether w represents an element in the subgroup generated by u_1, \dots, u_n ; and in the affirmative case compute an expression of w as a word in the u_i 's.*

Lemma 2.1. *If a torsion-free group satisfies SIP and WP, then it also satisfies MP.*

Proof. Let $G = \langle X \mid R \rangle$. Given words u, v_1, \dots, v_m in X , apply SIP to $H = \langle u \rangle$, and $K = \langle v_1, \dots, v_m \rangle$: since $H \cap K$ is cyclic (and so, finitely generated), SIP will always answer YES, and return a finite set of words w_1, \dots, w_p in X such that $H \cap K = \langle w_1, \dots, w_p \rangle = \langle u^r \rangle$, for some unknown $r \in \mathbb{Z}$.

Now, since each w_i must be a power of u^r (say $w_i = u^{r_i}$), we can compute the exponents $r_1, \dots, r_p \in \mathbb{Z}$ by brute force enumeration (even without using WP). Once we have obtained the integers $r_1, \dots, r_p \in \mathbb{Z}$, we can effectively compute the greatest common divisor $r = \gcd(r_1, \dots, r_p)$ and get $H \cap K = \langle w_1, \dots, w_p \rangle = \langle u^r \rangle$.

Now, it is clear that $u \in K$ if and only if $u \in H \cap K = \langle u^r \rangle$; i.e., if and only if $u = u^{rs}$, for some $s \in \mathbb{Z}$. To decide whether such an s exists, first apply WP to the input word u in order to decide whether $u = 1$ or not. In the affirmative case the answer is obviously YES; otherwise, $u \neq 1$ and torsion-freeness of G tells us that $u \in K$ (and the answer is YES) if and only if $r = \pm 1$. \square

Remark 2.2. Note that $CIP_{fg} \Rightarrow MP$ without any further condition, since $g \in \langle h_1, \dots, h_k \rangle$ if and only if $g \cdot \{1\} \cap 1 \cdot \langle h_1, \dots, h_k \rangle \neq \emptyset$.

Corollary 2.3. *For PC-groups, both SIP and CIP imply MP. In particular, SIP and CIP are unsolvable for $\mathbb{F}_2 \times \mathbb{F}_2 = G_{C_4}$, and hence for any PC-group G_Γ with $C_4 \leq \Gamma$.* \square

Remark 2.4. Note that the difference between properties TIP and ESIP is that the second one says nothing about $wH \cap w'K$ in the case when $H \cap K$ is not finitely generated, while TIP is required to answer about emptiness even in this case; this is a subtlety that will become important along the paper.

Finally, note that $(Hw)^{-1} = w^{-1}H$ and $w_1Hw_2 = w_1w_2(w_2^{-1}Hw_2) = w_1w_2H^{w_2}$. Therefore, the variants of CIP for right, left, and two-sided cosets are equivalent problems.

2.2 PC-groups

Below, we recall some well-known results about PC-groups we will need throughout the paper; we refer the reader to [5, 12, 13, 21] for detailed surveys, and further reference.

Theorem 2.5. (i) *Let Γ_1, Γ_2 be simple graphs. Then, the groups $G_{\Gamma_1}, G_{\Gamma_2}$ are isomorphic if and only if the graphs Γ_1, Γ_2 are isomorphic ([10]). In particular, the isomorphism problem is solvable within RAAGs.*

(ii) *The abelianization of the PC-group G_{Γ} is the free-abelian group of rank $|V\Gamma|$.*

(iii) *The word and conjugacy problems are solvable for RAAG's ([14, 22, 28, 30, 31]).*

(iv) *PC-groups are torsion-free ([2]).*

(v) *RAAGs are residually finite; in particular, they are Hopfian ([21]).*

(vi) *The maximum rank of a free-abelian subgroup of a RAAG G_{Γ} is the size of a largest clique in Γ ([21]).*

(vii) *Disjoint union and graph join correspond, respectively, to free product and direct product of groups: $G_{X \sqcup Y} \simeq G_X * G_Y$ and $G_{X \vee Y} \simeq G_X \times G_Y$.*

(viii) *A PC-group G_{Γ} splits as a nontrivial free product if and only if its defining graph Γ is disconnected.*

(ix) *A PC-group G_{Γ} splits as a nontrivial direct product if and only if its defining graph Γ is a join.*

(x) *The center of a PC-group G_{Γ} is the (free-abelian) subgroup generated by the set of central vertices in Γ .*

(xi) *Let Γ be an arbitrary simple graph, and Y a subset of vertices of Γ . Then, the subgroup of G_{Γ} generated by Y is again a PC-group, presented by the corresponding full subgraph, $\langle Y \rangle \simeq G_{\Gamma[Y]}$. \square*

Besides Droms groups, other subfamilies of PC-groups naturally arise as directly related with the intersection problem. For example, in [7], Delgado characterized the PC-groups satisfying the Howson property precisely as those being fully residually free (or free products of free-abelian groups).

Another interesting subfamily of PC-groups is that of chordal groups, that is the PC-groups presented by a finite chordal graph (i.e., one with no induced cycles of length strictly greater than three).

Clearly, Droms graphs are chordal and not the other way around. From Theorem 1.2, it is clear that Droms groups are *coherent* (every finitely generated subgroup is finitely presented). However, this last class was proved to be bigger, corresponding precisely to chordal groups, which turn out to have some nice algorithmic properties as well.

Theorem 2.6 (Droms, [9]). *Let Γ be a finite graph. Then, the RAAG G_Γ is coherent if and only if Γ is chordal.* \square

Theorem 2.7 (Kapovich–Weidmann–Myasnikov, [20]). *Let Γ be a finite chordal graph (i.e., G_Γ is a coherent RAAG). Then, (i) G_Γ has solvable membership problem; (ii) given a finite subset $S \subseteq G_\Gamma$, we can algorithmically find a presentation for the subgroup $\langle S \rangle \leq G_\Gamma$.* \square

We remark the pertinacious absence of $G_{C_4} = \mathbb{F}_2 \times \mathbb{F}_2$ from any family of algorithmically well behaved groups. However, the exact boundary of the class of RAAGs having solvable MP is not known: chordal groups have it, and G_{C_4} does not. Which of the groups G_{C_n} , for $n \geq 5$, have solvable MP? Is it possible to find a characterization of the RAAG's with solvable MP?

Finally, we recall that for submonoids, the exact border for the corresponding membership problem is already known: In [23] is proved that the submonoid membership problem is solvable in a PC-group if and only if it is Droms. Note that this implies, in particular, that G_{P_4} is a group with solvable subgroup membership problem (it is chordal), but unsolvable submonoid membership problem (it is not Droms).

2.3 Droms groups

Due to the recursive description in Corollary 1.5, any Droms graph Γ decomposes as the join of its central part $Z(\Gamma) \simeq K_m$, and the full subgraph $\Gamma_0 = \Gamma \setminus Z(\Gamma) \leq \Gamma$, that is $\Gamma = K_m \vee \Gamma_0$, where $m \geq 0$ and Γ_0 being either *empty* or a *disconnected* Droms graph; this is called the *primary decomposition* of Γ . In particular: (i) G_Γ is free-abelian $\Leftrightarrow \Gamma$ is complete $\Leftrightarrow Z(\Gamma) = \Gamma \Leftrightarrow \Gamma_0 = \emptyset$. (ii) Γ is connected $\Leftrightarrow Z(\Gamma) \neq \emptyset \Leftrightarrow m \geq 1 \Leftrightarrow \Gamma$ is a cone. (iii) G_Γ is centerless $\Leftrightarrow \Gamma$ is disconnected $\Leftrightarrow Z(\Gamma) = \emptyset \Leftrightarrow m = 0$.

Remark 2.8. However, all subgroups of Droms groups (including the non finitely generated ones) are again $\{P_4, C_4\}$ -free PC-groups. In particular, every finitely generated subgroup of a Droms group is again a Droms group.

Finally, we need the following algorithmic result for later use.

Proposition 2.9. *Let $G_\Gamma = \langle X \mid R \rangle$ be a Droms group. Then, there exists an algorithm which, given words $w_1(X), \dots, w_p(X)$ in the generators of G_Γ , (i) computes a basis for the subgroup $H = \langle w_1, \dots, w_p \rangle \leq G_\Gamma$; (ii) writes the basis elements in terms of the original generators, and vice versa.*

Proof. Since Droms graphs are chordal, by Theorem 2.7, we can effectively compute a finite presentation for H , say $H = \langle Y \mid S \rangle$. Then, one can exhaustively explore the tree of all possible Tietze transformations applied to $\langle Y \mid S \rangle$ until getting one, say $\langle Z \mid Q \rangle$, in graphical form (namely, with all relators being commutators of certain pairs of generators); this will be achieved in finite time because we know in advance that H is indeed a RAAG.

At this point, we know that $H = \langle w_1, \dots, w_p \rangle \simeq \langle Y \mid S \rangle \simeq \langle Z \mid Q \rangle$, and need to compute expressions for the basis elements $Z = \{z_1, \dots, z_r\}$ in terms of the w_i 's and vice versa. We start a brute force search, using the following two parallel procedures:

- (1) enumerate all homomorphisms $\langle Z \mid Q \rangle \rightarrow H = \langle w_1, \dots, w_p \rangle$; this can be done enumerating all possible r -tuples $(v_j)_{j=1}^r$ of words in $\{w_1, \dots, w_p\}^\pm$, and checking whether they determine a well-defined homomorphism.
- (2) for each such homomorphism $z_j \mapsto v_j$, $j = 1, \dots, r$, analyze whether it is onto H by enumerating all words in $\{v_1, \dots, v_r\}^\pm$, and checking whether each of w_1, \dots, w_p appear in the list.

Since we know that $\langle Z \mid Q \rangle \simeq H$, there exists a surjective homomorphism from $\langle Z \mid Q \rangle$ onto H and so, the above procedure will eventually find and output one of them, $z_j \mapsto v_j$, $j = 1, \dots, r$. Finally, since RAAG's are Hopfian (see Theorem 2.5(v)), such a surjective homomorphism is indeed an isomorphism. We then have the z_j 's written as words on the w_i 's, and the w_i 's as words on the z_j 's from the stopping criteria at step (2). \square

2.4 Proof of the main result

The main result in the present paper easily reduces to Theorems 1.7 and 1.8.

Proof of Theorem 1.4. Let Γ be a Droms graph, and let G_Γ be the corresponding Droms PC-group. We will prove Theorem 1.4 by induction on the number of vertices $|\Gamma|$. If $|\Gamma| = 0$, then $G_\Gamma = 1$, and obviously has solvable ESIP.

Now, consider a nonempty Droms graph Γ , and assume that every Droms PC-group with strictly less than $|\Gamma|$ vertices has solvable ESIP. Consider the primary decomposition of Γ , say $\Gamma = K_m \vee \Gamma_0$. If Γ_0 is empty then Γ is complete, $G_\Gamma \simeq \mathbf{Z}^m$ is free-abelian and so, it has solvable ESIP. Otherwise, Γ_0 is disconnected, say $\Gamma_0 = \Gamma_1 \vee \Gamma_2$ with Γ_1 and Γ_2 being Droms again. By induction, both G_{Γ_1} and G_{Γ_2} have solvable ESIP, by Theorem 1.8 $G_{\Gamma_0} = G_{\Gamma_1} * G_{\Gamma_2}$ has solvable ESIP, and by Theorem 1.7 $G_\Gamma = \mathbf{Z}^m \times G_{\Gamma_0}$ also has solvable ESIP. \square

3 The direct product case

This section is devoted to proving Theorems 1.6 and 1.7. To this end, we analyze the Droms groups presented by connected graphs.

3.1 Preparation

For all this subsection, we fix an arbitrary connected non-complete Droms graph Γ and its primary decomposition $\Gamma = K_m \vee \Gamma_0$, where $m \geq 1$ and Γ_0 is a disconnected Droms graph. Algebraically, $G_\Gamma = \mathbf{Z}^m \times G_{\Gamma_0}$, $Z(G_\Gamma) = \mathbf{Z}^m$, and G_{Γ_0} is a nontrivial free product. Let $V\Gamma_0 = X = \{x_1, \dots, x_n\}$ and $VK_m = T = \{t_1, \dots, t_m\}$.

Every element in G_Γ can be written as a word on $\{t_1, \dots, t_m, x_1, \dots, x_n\}$, where the t_i 's are free to move to any position. We will systematically write all these t_i 's ordered on the left, and we will abbreviate them as a vectorial power of a formal symbol 't'. This way, every element in $G_\Gamma = \mathbb{Z}^m \times G_{\Gamma_0}$ can be written in the form $t_1^{a_1} \cdots t_m^{a_m} u(x_1, \dots, x_n) = t^{\mathbf{a}} u(x_1, \dots, x_n)$, where $\mathbf{a} = (a_1, \dots, a_m) \in \mathbb{Z}^m$, and $u = u(x_1, \dots, x_n)$ is a word on the x_i 's. Clearly, the product of elements is then given by the rule $(t^{\mathbf{a}} u) \cdot (t^{\mathbf{b}} v) = t^{\mathbf{a}+\mathbf{b}} uv$.

Quotienting by the center of G_Γ gives rise to the short exact sequence

$$1 \longrightarrow \mathbb{Z}^m \longrightarrow G_\Gamma \xrightarrow{\pi_0} G_{\Gamma_0} \longrightarrow 1, \quad (2)$$

$$t^{\mathbf{a}} u \longmapsto u$$

where π_0 just erases the occurrences of letters in $T^\pm = \{t_1, \dots, t_m\}^\pm$.

Definition 3.1. For a given subgroup $H \leq G_\Gamma$, and an element $u \in G_{\Gamma_0}$, we define the (abelian) completion of u in H (the H -completion of u , for short) to be the set $\mathcal{C}_H(u) = \{\mathbf{a} \in \mathbb{Z}^m \mid t^{\mathbf{a}} u \in H\}$.

Lemma 3.2. *The completion $\mathcal{C}_H(u)$ is either empty (when $u \notin H\pi_0$), or a coset of $\mathbb{Z}^m \cap H$. More precisely, if $u_1, \dots, u_n \in H\pi_0$, and $\omega(u_1, \dots, u_n)$ is an arbitrary word on them, then $\mathcal{C}_H(\omega(u_1, \dots, u_n)) = \sum_{i=1}^n \omega_i \mathcal{C}_H(u_i)$, where $\omega_i = |\omega|_i$ is the total exponent of the variable u_i in ω . \square*

Lemma 3.3. *Let $G_\Gamma = \mathbb{Z}^m \times G_{\Gamma_0}$ be the primary decomposition of a connected Droms group. Then, any subgroup $H \leq G_\Gamma$ splits as $H = (\mathbb{Z}^m \cap H) \times H\pi_0\sigma$, where $\pi_0: G_\Gamma \rightarrow G_{\Gamma_0}$ is the natural projection killing the center of G_Γ , and $\sigma: H\pi_0 \rightarrow H$ is a section of $\pi_0|_H$.*

Proof. Let $X = \{x_1, \dots, x_n\}$ be the (finite) set of vertices of Γ_0 (i.e., let $G_{\Gamma_0} = \langle X \mid R \rangle$, where $R \subseteq [X, X]$); and let $\mathbb{Z}^m = \langle t_1, \dots, t_m \mid [t_i, t_j] \forall i, j \rangle$. Now, consider the restriction to $H \leq \mathbb{Z}^m \times G_{\Gamma_0}$ of the natural short exact sequence (2):

$$1 \longrightarrow \mathbb{Z}^m \simeq Z(G_\Gamma) \longrightarrow G_\Gamma \xrightarrow{\pi_0} G_{\Gamma_0} \longrightarrow 1 \quad (3)$$

$$\quad \quad \quad \vee \quad \quad \quad \vee \quad \quad \quad \vee$$

$$1 \longrightarrow \mathbb{Z}^m \cap H \longrightarrow H \longrightarrow H\pi_0 \longrightarrow 1, \quad (4)$$

Since G_{Γ_0} is Droms, we know that $H\pi_0 \leq G_{\Gamma_0}$ is again a PC-group. Thus, there exists a (not necessarily finite) subset $Y = \{y_j\}_j \subseteq G_{\Gamma_0}$ such that $H\pi_0 \simeq \langle Y \mid S \rangle$, where S is a certain collection of commutators of the y_j 's.

Now, observe that any map $\sigma: Y \rightarrow H$ sending each $y_j \in Y$ back to any of its π_0 -preimages in H will necessarily respect the relations in S : indeed, for each commutator $[y_i, y_j] \in S$, we have $[y_i\sigma, y_j\sigma] = [t^{\mathbf{a}_i} y_i, t^{\mathbf{a}_j} y_j] = [y_i, y_j]$ (for certain abelian completions $\mathbf{a}_i, \mathbf{a}_j \in \mathbb{Z}^m$). Therefore, any such map σ defines a (injective) section of the restriction $\pi_0|_H$. Thus, the short exact sequence (4) splits and, for any such section σ , $H\pi_0 \simeq H\pi_0\sigma \leq H$. Moreover, since the kernel of the subextension (4) lies in the center of G_Γ , the conjugation action is trivial, and the claimed result follows. \square

Corollary 3.4. *In the above situation, the subgroup $H \leq G_\Gamma$ is finitely generated if and only if $H\pi_0 \leq G_{\Gamma_0}$ is finitely generated. \square*

Remark 3.5. For any $H = \langle t^{b_1}, \dots, t^{b_r}, t^{a_1} u_1, \dots, t^{a_s} u_s \rangle \leq G_\Gamma = \mathbb{Z}^m \times G_{\Gamma_0}$, where $u_i \neq 1$ for all $i = 1, \dots, s$, we have $\langle t^{b_1}, \dots, t^{b_r} \rangle \leq \mathbb{Z}^m \cap H = Z(G_\Gamma) \cap H \leq Z(H)$, but these two inclusions are, in general, not equalities: for the first one, a nontrivial product of the last s generators could, in principle, be equal to t^c for some element $c \notin \langle b_1, \dots, b_r \rangle$; and for the second we could have, for example, $u_1 = \dots = u_s \neq 1$ so that $t^{a_1} u_1$ belongs to $Z(H)$ but not to $Z(G_\Gamma)$.

Let us consider now two finitely generated subgroups $H_1, H_2 \leq G_\Gamma$ and analyze when the intersection $H_1 \cap H_2$ is again finitely generated. We will see that the behaviour of the embedding $(H_1 \cap H_2)\pi_0 \leq (H_1)\pi_0 \cap (H_2)\pi_0$ is crucial to this end.

Lemma 3.6. *Let $H_1, H_2 \leq G_\Gamma$. Then,*

- (i) $(H_1 \cap H_2)\pi_0 \leq H_1\pi_0 \cap H_2\pi_0$, sometimes with strict inclusion;
- (ii) $(H_1 \cap H_2)\pi_0 \triangleleft H_1\pi_0 \cap H_2\pi_0$;
- (iii) $[(H_1)\pi_0 \cap (H_2)\pi_0, (H_1)\pi_0 \cap (H_2)\pi_0] \leq (H_1 \cap H_2)\pi_0$.

Proof. (i). This is clear.

(ii). To see normality, consider $u \in (H_1 \cap H_2)\pi_0$, and $v \in (H_1)\pi_0 \cap (H_2)\pi_0$; then, there exist elements $t^a u \in H_1 \cap H_2$, and $t^{b_i} v \in H_i$, for $i = 1, 2$. Now observe that $t^a(v^{-1}uv) = v^{-1}(t^a u)v = (t^{b_i} v)^{-1}(t^a u)(t^{b_i} v) \in H_i$, for $i = 1, 2$. Thus, $t^a v^{-1}uv \in H_1 \cap H_2$ and so, $v^{-1}uv \in (H_1 \cap H_2)\pi_0$.

(iii). Take $u, v \in (H_1)\pi_0 \cap (H_2)\pi_0$; then, there exist elements $t^{a_i} u \in H_i$, $t^{b_i} v \in H_i$, for $i = 1, 2$. Now, observe that $[u, v] = u^{-1}v^{-1}uv = (t^{a_i} u)^{-1}(t^{b_i} v)^{-1}(t^{a_i} u)(t^{b_i} v) \in H_i$, for $i = 1, 2$. Thus, $[u, v]$ belongs to $H_1 \cap H_2$, and to $(H_1 \cap H_2)\pi_0$, as claimed. \square

Lemma 3.7. *If $(H_1 \cap H_2)\pi_0$ (and so, $H_1 \cap H_2$) is finitely generated, then $(H_1)\pi_0 \cap (H_2)\pi_0$ is also finitely generated.*

Proof. Let us assume that $(H_1)\pi_0 \cap (H_2)\pi_0$ is not finitely generated and find a contradiction.

By Remark 2.8, $(H_1)\pi_0 \cap (H_2)\pi_0$ is again a PC-group with infinite $\{P_4, C_4\}$ -free commutation graph, say Δ , and $(H_1 \cap H_2)\pi_0 \leq G_{\Delta'} \leq G_\Delta = (H_1)\pi_0 \cap (H_2)\pi_0$, where Δ' is the full subgraph of Δ determined by the vertices appearing in the reduced expressions of elements in $(H_1 \cap H_2)\pi_0$. Note that the assumption of finite generability for $(H_1 \cap H_2)\pi_0$ implies that Δ' is finite. Note also that, by construction, Δ' is minimal, i.e., for any $x \in V\Delta'$, there exists an element $g \in (H_1 \cap H_2)\pi_0$ such that $g \notin G_{\Delta' \setminus \{x\}}$.

In this situation, $(H_1)\pi_0 \cap (H_2)\pi_0$ cannot be abelian since, if so, we would have a non finitely generated free-abelian group embedded in the finitely generated PC-group G_Γ , which is not possible (see Theorem 2.5(vi)). So, Δ is not complete. Take two non-adjacent vertices, say u, v , from Δ ; since $\mathbb{F}_2 \simeq \langle u, v \rangle \leq G_\Delta$, Lemma 3.6(iii) tells us that $\mathbb{F}_\infty \simeq [\langle u, v \rangle, \langle u, v \rangle] \leq [G_\Delta, G_\Delta] \leq (H_1 \cap H_2)\pi_0$ and thus, $(H_1 \cap H_2)\pi_0$ is not abelian either. Accordingly, neither the infinite graph Δ , nor the finite graph Δ' are complete.

Suppose now there is a missing edge between some vertex $x \in V\Delta'$ and some vertex $y \in V\Delta \setminus V\Delta'$. Take an element $g \in (H_1 \cap H_2)\pi_0$ with $g \notin G(\Delta' \setminus \{x\})$ and Lemma 3.6(ii) would tell us that $y^{-1}gy \in G_{\Delta'}$, which is a contradiction.

Hence, in Δ , every vertex from Δ' is connected to every vertex outside Δ' . But now, take two non-adjacent vertices x_1, x_2 from Δ' and two non-adjacent vertices y_1, y_2 from $\Delta \setminus \Delta'$ (there must also be some since $\Delta \setminus \Delta'$ is infinite and \mathbb{Z}^∞ does not embed into G_Γ). Then, the full subgraph of Δ with vertex set $\{x_1, x_2, y_1, y_2\}$ form a copy of C_4 , a contradiction with Δ being Droms. \square

3.2 Proofs of Theorems 1.6 and 1.7

Proof of Theorem 1.6. First of all, observe that we can restrict ourselves to the case where G is a disconnected Droms group.

Therefore, we consider $\Gamma = K_m \vee \Gamma_0$ (where $m \geq 1$ and Γ_0 is a disconnected Droms graph), we will assume SIP is solvable for G_{Γ_0} , and we will prove it solvable for $G_\Gamma = \mathbb{Z}^m \times G_{\Gamma_0}$. Let $X = V\Gamma_0 = \{x_1, \dots, x_n\}$ and $T = VK_m = \{t_1, \dots, t_m\}$.

Given a finite set of generators for a subgroup, say $H_1 \leq G_\Gamma$, the first step is to improve them: project them to G_{Γ_0} , and then apply Proposition 2.9 to compute a basis for $H_1\pi_0$, say $\{u_1, \dots, u_{n_1}\}$ with commutation graph G_{Δ_1} . The respective completions, say $t^{a_1} u_1, \dots, t^{a_{n_1}} u_{n_1} \in H_1$ can be computed from the words expressing the u_i 's in terms of the projected generators, and recomputing them on the original generators for H_1 .

Now, for each of the original generators of H_1 , say $t^c v$, we can write $v \in H_1\pi_0$ in terms of the basis u_1, \dots, u_{n_1} , say $v = v(u_1, \dots, u_{n_1})$ and compute $v(t^{a_1} u_1, \dots, t^{a_{n_1}} u_{n_1}) = t^d v(u_1, \dots, u_{n_1}) = t^d v$. Since $t^c v, t^d v \in H_1$, we get $t^{c-d} \in H_1 \cap \mathbb{Z}^m$. Repeating this operation for each generator of H_1 , we get a generating set for $H_1 \cap \mathbb{Z}^m$ which is easily reducible to a free-abelian basis, say $\{t^{b_1}, \dots, t^{b_{m_1}}\}$. In this way, we can compute bases for H_1 and H_2 :

$$\{t^{b_1}, \dots, t^{b_{m_1}}, t^{a_1} u_1, \dots, t^{a_{n_1}} u_{n_1}\} \text{ and } \{t^{b'_1}, \dots, t^{b'_{m_2}}, t^{a'_1} u'_1, \dots, t^{a'_{n_2}} u'_{n_2}\}, \quad (5)$$

where $\{t^{b_1}, \dots, t^{b_{m_1}}\}$ and $\{t^{b'_1}, \dots, t^{b'_{m_2}}\}$ are free-abelian bases of $L_1 = H_1 \cap \mathbb{Z}^m$ and $L_2 = H_2 \cap \mathbb{Z}^m$, respectively; and $\{u_1, \dots, u_{n_1}\}$ and $\{u'_1, \dots, u'_{n_2}\}$ are basis of $H_1\pi_0$ and $H_2\pi_0$, with commutation graphs Δ_1, Δ_2 . That is, $H_1\pi_0 \simeq G_{\Delta_1}$ and $H_2\pi_0 \simeq G_{\Delta_2}$.

Now, the solvability of SIP in G_{Γ_0} (assumed by hypothesis) allows us to decide whether $H_1\pi_0 \cap H_2\pi_0$ is finitely generated or not. If not, then (by Lemma 3.7) neither is $(H_1 \cap H_2)\pi_0$, and we are done. Thus, we can assume that $H_1\pi_0 \cap H_2\pi_0$ is finitely generated. Then, the hypothesis provides a finite set of generators and hence a basis — say $W = \{w_1, \dots, w_{n_3}\}$ with commutation graph Δ_3 — for $H_1\pi_0 \cap H_2\pi_0$. That is, $H_1\pi_0 \cap H_2\pi_0 = G_{\Delta_3} = \langle w_1, \dots, w_{n_3} \rangle \leq G_{\Gamma_0}$, where the w_i 's are words on X . After writing each $w_i \in W$ as a word on U and U' respectively — say $w_i = \omega_i(u_1, \dots, u_{n_1})$ and $w_i = \omega'_i(u'_1, \dots, u'_{n_2})$ — we obtain a description of the inclusions $\iota_1: H_1\pi_0 \cap H_2\pi_0 \rightarrow H_1\pi_0$, and $\iota_2: H_1\pi_0 \cap H_2\pi_0 \rightarrow H_2\pi_0$ in terms of the corresponding bases.

Abelianizing ι_1 and ι_2 , we get the integral matrices \mathbf{P}_1 (of size $n_3 \times n_1$), and \mathbf{P}_2 (of size $n_3 \times n_2$) and complete the upper half of Figure 2, where the ρ_i 's are the corresponding abelianization maps. Note that, even though ι_1 and ι_2 are injective, their abelianizations \mathbf{P}_1 and \mathbf{P}_2 need not be (n_3 could very well be bigger than n_1 or n_2).

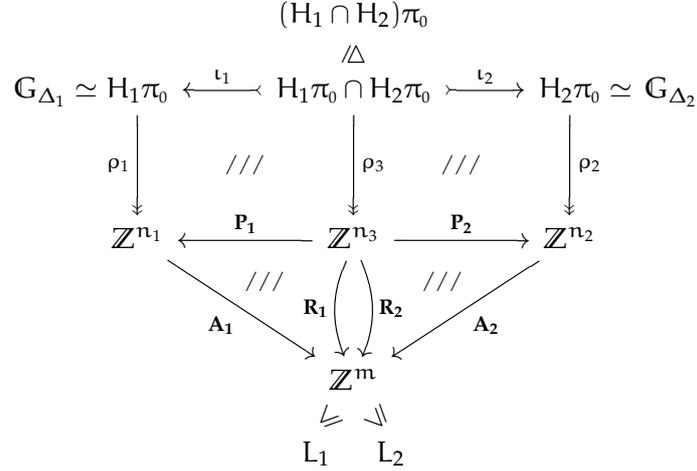


Figure 2: Intersection diagram for subgroups of Droms groups

Now, we can recompute the words ω_i (resp., ω'_i) as words on the $(t^{a_i} u_i)$'s (resp., on the $(t^{a'_i} u'_i)$'s) to get particular preimages of the w_i 's in \mathbf{H}_1 (resp., \mathbf{H}_2). Namely,

$$\begin{aligned}
\omega_i(t^{a_1} u_1, \dots, t^{a_{n_1}} u_{n_1}) &= t^{\omega_i \mathbf{A}_1} \omega_i(u_1, \dots, u_{n_1}) = t^{\omega_i \mathbf{A}_1} w_i \in \mathbf{H}_1, \\
\omega'_i(t^{a'_1} u'_1, \dots, t^{a'_{n_2}} u'_{n_2}) &= t^{\omega'_i \mathbf{A}_2} \omega'_i(u'_1, \dots, u'_{n_2}) = t^{\omega'_i \mathbf{A}_2} w_i \in \mathbf{H}_2,
\end{aligned}$$

where $\omega_i = (\omega_i)^{\text{ab}}$, $\omega'_i = (\omega'_i)^{\text{ab}}$; and $\mathbf{A}_1, \mathbf{A}_2$ are the integral matrices (of sizes $n_1 \times m$ and $n_2 \times m$) having as rows $\{\mathbf{a}_1, \dots, \mathbf{a}_{n_1}\}$ and $\{\mathbf{a}'_1, \dots, \mathbf{a}'_{n_2}\}$ respectively. Hence, the abelian completions of $w_i \in \mathbf{H}_1\pi_0 \cap \mathbf{H}_2\pi_0$ in \mathbf{H}_1 and \mathbf{H}_2 are the linear varieties:

$$\begin{aligned}
\mathcal{C}_{\mathbf{H}_1}(w_i) &= \omega_i \mathbf{A}_1 + \mathbf{L}_1 = w_i \iota_1 \rho_1 \mathbf{A}_1 + \mathbf{L}_1 = w_i \rho_3 \mathbf{R}_1 + \mathbf{L}_1, \\
\mathcal{C}_{\mathbf{H}_2}(w_i) &= \omega'_i \mathbf{A}_2 + \mathbf{L}_2 = w_i \iota_2 \rho_2 \mathbf{A}_2 + \mathbf{L}_2 = w_i \rho_3 \mathbf{R}_2 + \mathbf{L}_2,
\end{aligned}$$

where $\mathbf{L}_j = \mathbb{Z}^m \cap \mathbf{H}_j$, and we have used the commutation $\iota_j \rho_j = \rho_3 \mathbf{P}_j$ together with the definition $\mathbf{R}_j := \mathbf{P}_j \mathbf{A}_j$, for $j = 1, 2$; see Figure 2. Note that all maps and matrices involved in Figure 2 are explicitly computable from the data.

To finish our argument, it suffices to understand which elements of $\mathbf{H}_1\pi_0 \cap \mathbf{H}_2\pi_0$ belong to $(\mathbf{H}_1 \cap \mathbf{H}_2)\pi_0$. They are, precisely, those whose \mathbf{H}_1 -completion and \mathbf{H}_2 -completion intersect:

$$\begin{aligned}
(\mathbf{H}_1 \cap \mathbf{H}_2)\pi_0 &= \{w \in \mathbf{H}_1\pi_0 \cap \mathbf{H}_2\pi_0 \mid (w\rho_3 \mathbf{P}_1 \mathbf{A}_1 + \mathbf{L}_1) \cap (w\rho_3 \mathbf{P}_2 \mathbf{A}_2 + \mathbf{L}_2) \neq \emptyset\} \\
&= (\{\mathbf{d} \in \mathbb{Z}^{n_3} \mid (\mathbf{d}\mathbf{R}_1 + \mathbf{L}_1) \cap (\mathbf{d}\mathbf{R}_2 + \mathbf{L}_2) \neq \emptyset\}) \rho_3^{-1} \\
&= (\{\mathbf{d} \in \mathbb{Z}^{n_3} \mid \mathbf{d}(\mathbf{R}_1 - \mathbf{R}_2) \in \mathbf{L}_1 + \mathbf{L}_2\}) \rho_3^{-1} \\
&= (\mathbf{L}_1 + \mathbf{L}_2)(\mathbf{R}_1 - \mathbf{R}_2)^{-1} \rho_3^{-1} = \mathbf{M} \rho_3^{-1},
\end{aligned} \tag{6}$$

where $M := (L_1 + L_2)(\mathbf{R}_1 - \mathbf{R}_2)^{-1}$ denotes the full preimage of $L_1 + L_2$ by the matrix $\mathbf{R}_1 - \mathbf{R}_2$ for which a basis is clearly computable using linear algebra. At this point, we can decide whether $(H_1 \cap H_2)\pi_0$ is finitely generated or not by distinguishing two cases.

If Δ_3 is complete (this includes the case where Δ_3 is empty and $n_3 = 0$), then $H_1\pi_0 \cap H_2\pi_0 \simeq \mathbb{Z}^{n_3}$ is abelian, ρ_3 is the identity, and $(H_1 \cap H_2)\pi_0 = (L_1 + L_2)(\mathbf{R}_1 - \mathbf{R}_2)^{-1} = M$ is always finitely generated and computable.

So, assume Δ_3 is not complete. Since it is a Droms graph, it will have a primary decomposition, say $\Delta_3 = K_{n_4} \vee \Delta_5$, where $n_4 \geq 0$, and Δ_5 is Droms again, disconnected, and with $|\mathcal{V}\Delta_5| = n_5 = n_3 - n_4 \geq 2$. Let us rename the vertices $\{w_1, \dots, w_{n_3}\}$ of Δ_3 as $\mathcal{V}K_{n_4} = \{z_1, \dots, z_{n_4}\}$, and $\mathcal{V}\Delta_5 = \{y_1, \dots, y_{n_5}\}$, depending on whether they belong to K_{n_4} or Δ_5 . This means that $H_1\pi_0 \cap H_2\pi_0 \simeq \mathbf{G}_{\Delta_3} = \mathbb{Z}^{n_4} \times \mathbf{G}_{\Delta_5}$, where $n_4 \geq 0$, and $\mathbf{G}_{\Delta_5} \neq 1$ decomposes as a nontrivial free product. Furthermore, the normal subgroup $(H_1 \cap H_2)\pi_0 \trianglelefteq \mathbf{G}_{\Delta_3}$ is *not* contained in \mathbb{Z}^{n_4} (taking two vertices, say y_i, y_j , in different components of Δ_5 , Lemma 3.6(iii) tells us that $1 \neq [y_i, y_j] \in (H_1 \cap H_2)\pi_0$). In this situation, the abelianization map $\rho_3: \mathbf{G}_{\Delta_3} \rightarrow \mathbb{Z}^{n_3}$ is the identity on the center \mathbb{Z}^{n_4} of \mathbf{G}_{Δ_3} and so, can be decomposed in the form $\rho_3 = \text{id} \times \rho_5: \mathbf{G}_{\Delta_3} = \mathbb{Z}^{n_4} \times \mathbf{G}_{\Delta_5} \rightarrow \mathbb{Z}^{n_4} \times \mathbb{Z}^{n_5} = \mathbb{Z}^{n_3}$, $(\mathbf{c}, \mathbf{v}) \mapsto (\mathbf{c}, \mathbf{v})$. where \mathbf{v} denotes the abelianization $\mathbf{v} = \mathbf{v}^{\text{ab}} \in \mathbb{Z}^{n_5}$. Of course, if $n_4 = 0$ then $\rho_5 = \rho_3$.

Now consider the image of $(H_1 \cap H_2)\pi_0$ under the projection $\pi_1: \mathbb{Z}^{n_4} \times \mathbf{G}_{\Delta_5} \rightarrow \mathbf{G}_{\Delta_5}$, which is nontrivial since $(H_1 \cap H_2)\pi_0 \not\leq \mathbb{Z}^{n_4}$. We have $1 \neq (H_1 \cap H_2)\pi_0\pi_1 \trianglelefteq \mathbf{G}_{\Delta_5}$, a nontrivial normal subgroup in a group which decomposes as a nontrivial free product. Therefore,

$$\begin{aligned} H_1 \cap H_2 \text{ is f.g.} &\Leftrightarrow (H_1 \cap H_2)\pi_0\pi_1 \text{ is f.g.} \Leftrightarrow (H_1 \cap H_2)\pi_0\pi_1 \trianglelefteq_{\text{fi}} \mathbf{G}_{\Delta_5} \\ &\Leftrightarrow M\rho_3^{-1}\pi_1 \trianglelefteq_{\text{fi}} \mathbf{G}_{\Delta_5} \Leftrightarrow M\pi_1^{\text{ab}}\rho_5^{-1} \trianglelefteq_{\text{fi}} \mathbf{G}_{\Delta_5} \\ &\Leftrightarrow M\pi_1^{\text{ab}} \trianglelefteq_{\text{fi}} \mathbb{Z}^{n_5} \Leftrightarrow \text{rk}(M\pi_1^{\text{ab}}) = n_5. \end{aligned} \quad (7)$$

The first of these equivalences is a (double) application of Corollary 3.4. The second one is an application of the following theorem in [3, Section 6] by B. Baumslag: *Let G be the free product of two nontrivial groups. Let H be a finitely generated subgroup containing a nontrivial normal subgroup of G . Then H is of finite index in G .* The fourth equivalence is correct since $\pi_1\rho_5 = \rho_3\pi_1^{\text{ab}}$ and all of them are surjective maps. Finally, the fifth equivalence is correct because following backwards the epimorphism ρ_5 , a subgroup $M\pi_1^{\text{ab}} \leq \mathbb{Z}^{n_5}$ is of finite index if and only if its full preimage $M\pi_1^{\text{ab}}\rho_5^{-1}$ is of finite index in \mathbf{G}_{Δ_5} , in which case the two indices do coincide, namely, $[\mathbb{Z}^{n_5} : M\pi_1^{\text{ab}}] = [\mathbf{G}_{\Delta_5} : M\pi_1^{\text{ab}}\rho_5^{-1}]$; see Figure 3.

$$\begin{array}{ccccc} (H_1 \cap H_2)\pi_0 \simeq M\rho_3^{-1} \leq & \mathbf{G}_{\Delta_3} & \xrightarrow{\pi_1} & \mathbf{G}_{\Delta_5} & \geq M\pi_1^{\text{ab}}\rho_5^{-1} \\ & \rho_3 \downarrow & & \downarrow \rho_5 & \\ (L_1 + L_2)(\mathbf{R}_1 - \mathbf{R}_2)^{-1} = M \leq & \mathbb{Z}^{n_3} & \xrightarrow{\pi_1^{\text{ab}}} & \mathbb{Z}^{n_5} & \geq M\pi_1^{\text{ab}} \end{array}$$

Figure 3: The map π_1 and its abelianization

Since the map π_1^{ab} is computable, the last condition in (7) can be effectively checked. Hence, we can algorithmically decide whether $H_1 \cap H_2$ is finitely generated or not (ultimately, in terms of some integral matrix having the correct rank). This solves the decision part of SIP.

It only remains to compute a finite set of generators for $H_1 \cap H_2$ assuming it is finitely generated, i.e., assuming the equivalent conditions in (7) are satisfied.

We first use linear algebra to compute a finite family C of coset representatives of \mathbb{Z}^{n_5} modulo $M\pi_1^{\text{ab}}$. Then, choose arbitrary ρ_5 -preimages in \mathbb{G}_{Δ_5} , say $\{v_1, \dots, v_r\}$, where $r = [\mathbb{G}_{\Delta_5} : M\pi_1^{\text{ab}}\rho_5^{-1}] = [\mathbb{Z}^{n_5} : M\pi_1^{\text{ab}}]$ (we can take, for example, $y_1^{a_1} \cdots y_{n_5}^{a_{n_5}} \in \mathbb{G}_{\Delta_5}$ for each vector $\mathbf{a} = (a_1, \dots, a_{n_5}) \in \mathbb{Z}^{n_5}$). Now, construct the Schreier graph of the subgroup $(H_1 \cap H_2)\pi_0\pi_1 = M\rho_3^{-1}\pi_1 = M\pi_1^{\text{ab}}\rho_5^{-1} \leq_{\text{fi}} \mathbb{G}_{\Delta_5}$ with respect to $\mathbb{V}\Delta_5 = \{y_1, \dots, y_{n_5}\}$, in the following way: draw as vertices the cosets $[v_1], \dots, [v_r]$; then, for every $[v_i]$ ($i = 1, \dots, r$), and every y_j ($j = 1, \dots, n_5$), draw an edge labelled y_j from $[v_i]$ to $[v_i y_j]$. Here, we need to algorithmically recognize which is the coset $[v_i y_j]$ from our list of vertices, but this is easy since: $[v_i y_j] = [v_k] \Leftrightarrow v_i y_j v_k^{-1} \in M\pi_1^{\text{ab}}\rho_5^{-1} \Leftrightarrow (v_i y_j v_k^{-1})\rho_5 \in M\pi_1^{\text{ab}}$.

From the Schreier graph of $(H_1 \cap H_2)\pi_0\pi_1 \leq_{\text{fi}} \mathbb{G}_{\Delta_5}$, we can obtain a finite set of generators for $(H_1 \cap H_2)\pi_0\pi_1$ just reading the labels of the closed paths $\mathbf{T}[e]$ corresponding to the arcs, e , outside a chosen maximal tree \mathbf{T} . These will be words on $\mathbb{V}\Delta_5 = \{y_1, \dots, y_{n_5}\}$, i.e., elements of \mathbb{G}_{Δ_5} not using the central vertices $\{z_1, \dots, z_{n_4}\}$.

The next step is to lift the obtained generators to generators of $(H_1 \cap H_2)\pi_0$, pulling them back through π_1 . For each one of them, say $g(y_1, \dots, y_{n_5})$, we look for its preimages in $(H_1 \cap H_2)\pi_0$; they all are of the form $z_1^{\lambda_1} \cdots z_{n_4}^{\lambda_{n_4}} g(y_1, \dots, y_{n_5})$, where the unknowns $\lambda_1, \dots, \lambda_{n_4} \in \mathbb{Z}$ can be found by solving the system of linear equations coming from the fact $z_1^{\lambda_1} \cdots z_{n_4}^{\lambda_{n_4}} g(y_1, \dots, y_{n_5}) \in M\rho_3^{-1}$. That is, $(\lambda_1, \dots, \lambda_{n_4}, |g|_1, \dots, |g|_{n_5})(\mathbf{R}_1 - \mathbf{R}_2) \in L_1 + L_2$.

For each such $g(y_1, \dots, y_{n_5})$, we compute a particular preimage of the previous form and put them all, together with a free-abelian basis for

$$\ker \pi_1 \cap (H_1 \cap H_2)\pi_0 = \left\{ z_1^{\lambda_1} \cdots z_{n_4}^{\lambda_{n_4}} : (\lambda_1, \dots, \lambda_{n_4}, 0, \dots, 0)(\mathbf{R}_1 - \mathbf{R}_2) \in L_1 + L_2 \right\},$$

to constitute a set of generators for $(H_1 \cap H_2)\pi_0$.

Finally, we have to lift these generators for $(H_1 \cap H_2)\pi_0$, to a set of generators for $H_1 \cap H_2$: for each such generator, say h_j , write it as a word $h_j = \omega_j(u_1, \dots, u_{n_1})$ and as a word $h_j = \omega'_j(u'_1, \dots, u'_{n_2})$ in the original bases \mathbf{U} for $H_1\pi_0$ and \mathbf{U}' for $H_2\pi_0$, respectively. Now, reevaluate each ω_j and ω'_j in the corresponding basis elements from (5) for H_1 and H_2 respectively, to obtain vectors $\mathbf{c}_j, \mathbf{c}'_j \in \mathbb{Z}^m$ such that:

$$\begin{aligned} \omega_j(t^{\mathbf{a}^1} u_1, \dots, t^{\mathbf{a}^{n_1}} u_{n_1}) &= t^{\mathbf{c}_j} \omega_j(u_1, \dots, u_{n_1}) = t^{\mathbf{c}_j} h_j \in H_1, \\ \omega'_j(t^{\mathbf{a}'^1} u'_1, \dots, t^{\mathbf{a}'^{n_2}} u'_{n_2}) &= t^{\mathbf{c}'_j} \omega'_j(u'_1, \dots, u'_{n_2}) = t^{\mathbf{c}'_j} h_j \in H_2. \end{aligned}$$

Finally, for each j , compute a vector $\mathbf{d}_j \in (\mathbf{c}_j + L_1) \cap (\mathbf{c}'_j + L_2)$ (note that these intersections of linear varieties must be nonempty because $h_j \in (H_1 \cap H_2)\pi_0$), and consider the element $t^{\mathbf{d}_j} h_j \in H_1 \cap H_2$. All these elements $t^{\mathbf{d}_j} h_j$, together with a free-abelian basis for $H_1 \cap H_2 \cap$

$\mathbb{Z}^m = (H_1 \cap \mathbb{Z}^m) \cap (H_2 \cap \mathbb{Z}^m) = L_1 \cap L_2$ constitute the desired set of generators for $H_1 \cap H_2$, and the proof is completed. \square

Below, we extend the previous arguments to prove Theorem 1.7.

Proof of Theorem 1.7. By exactly the same argument as before, we can reduce to a non-complete connected Droms graph Γ with primary decomposition $\Gamma = K_m \vee \Gamma_0$ (where $m \geq 1$ and Γ_0 is a disconnected Droms graph), we assume ESIP to be solvable for G_{Γ_0} , and we have to solve it for $G_\Gamma = \mathbb{Z}^m \times G_{\Gamma_0}$.

We are given finite sets of generators for two subgroups $H_1, H_2 \leq G_\Gamma$, and two extra elements $t^{\mathbf{a}}u, t^{\mathbf{a}'}u' \in G_\Gamma$. Since the solvability of ESIP implies that of SIP, we can apply Theorem 1.6 to effectively decide whether $H_1 \cap H_2$ is finitely generated or not, and in the affirmative case compute a basis for $H_1 \cap H_2$. We assume all the notation developed along the proof of Theorem 1.6.

Now, if $H_1 \cap H_2$ is not finitely generated there is nothing else to do; otherwise, we can compute a basis, say $\{v_1, \dots, v_p\}$, for $H_1 \cap H_2$ and we have to decide whether the coset intersection $(t^{\mathbf{a}}u)H_1 \cap (t^{\mathbf{a}'}u')H_2$ is empty or not. Note that $(t^{\mathbf{a}}u)H_1 \cap (t^{\mathbf{a}'}u')H_2 = \emptyset$ if and only if $((t^{\mathbf{a}}u)H_1 \cap (t^{\mathbf{a}'}u')H_2)\pi_0 = \emptyset$, and that $((t^{\mathbf{a}}u)H_1 \cap (t^{\mathbf{a}'}u')H_2)\pi_0 \subseteq ((t^{\mathbf{a}}u)H_1)\pi_0 \cap ((t^{\mathbf{a}'}u')H_2)\pi_0 = u(H_1\pi_0) \cap u'(H_2\pi_0)$. Then, since $H_1 \cap H_2$ is finitely generated, we know from Lemma 3.7 that $H_1\pi_0 \cap H_2\pi_0$ is finitely generated as well. Hence, an application of the ESIP solvability hypothesis for G_{Γ_0} , tells us whether the coset intersection $u(H_1\pi_0) \cap u'(H_2\pi_0)$ is empty or not. If it is empty, then $((t^{\mathbf{a}}u)H_1 \cap (t^{\mathbf{a}'}u')H_2)\pi_0$ is empty as well, and we are done.

Otherwise, $u(H_1\pi_0) \cap u'(H_2\pi_0) \neq \emptyset$, and the hypothesis gives us an element $v_0 \in u(H_1\pi_0) \cap u'(H_2\pi_0)$ as a word on $V\Gamma_0 = \{x_1, \dots, x_n\}$; further, $u(H_1\pi_0) \cap u'(H_2\pi_0) = v_0(H_1\pi_0 \cap H_2\pi_0)$.

Observe that $((t^{\mathbf{a}}u)H_1 \cap (t^{\mathbf{a}'}u')H_2)\pi_0$ consists precisely of those elements v_0w , with $w \in H_1\pi_0 \cap H_2\pi_0$, for which there exists a vector $\mathbf{c} \in \mathbb{Z}^m$ such that $t^{\mathbf{c}}v_0w \in (t^{\mathbf{a}}u)H_1 \cap (t^{\mathbf{a}'}u')H_2$; that is, such that $t^{\mathbf{c}-\mathbf{a}}u^{-1}v_0w \in H_1$, and $t^{\mathbf{c}-\mathbf{a}'}(u')^{-1}v_0w \in H_2$. That is, $\mathbf{c} - \mathbf{a} \in \mathcal{C}_{H_1}(u^{-1}v_0w)$, and $\mathbf{c} - \mathbf{a}' \in \mathcal{C}_{H_2}((u')^{-1}v_0w)$. Hence, $((t^{\mathbf{a}}u)H_1 \cap (t^{\mathbf{a}'}u')H_2)\pi_0 = \emptyset$ if and only if, for all $w \in H_1\pi_0 \cap H_2\pi_0$, $(\mathbf{a} + \mathcal{C}_{H_1}(u^{-1}v_0w)) \cap (\mathbf{a}' + \mathcal{C}_{H_2}((u')^{-1}v_0w)) = \emptyset$.

Fix an arbitrary word $w = \omega(w_1, \dots, w_{n_3}) \in H_1\pi_0 \cap H_2\pi_0$ (for its abelianization, write $|\omega|_i = \lambda_i$, for $i = 1, \dots, n_3$). Choose vectors $\mathbf{c} \in \mathcal{C}_{H_1}(u^{-1}v_0)$, $\mathbf{c}' \in \mathcal{C}_{H_2}((u')^{-1}v_0)$, and $\mathbf{d}_i \in \mathcal{C}_{H_1}(w_i)$, $\mathbf{d}'_i \in \mathcal{C}_{H_2}(w_i)$, for $i = 1, \dots, n_3$. By Lemma 3.2:

$$\begin{aligned} (\mathbf{a} + \mathcal{C}_{H_1}(u^{-1}v_0w)) \cap (\mathbf{a}' + \mathcal{C}_{H_2}((u')^{-1}v_0w)) &= \\ &= (\mathbf{a} + \mathbf{c} + \mathcal{C}_{H_1}(w)) \cap (\mathbf{a}' + \mathbf{c}' + \mathcal{C}_{H_2}(w)) \\ &= \left(\mathbf{a} + \mathbf{c} + \sum_{i=1}^{n_3} \lambda_i \mathcal{C}_{H_1}(w_i) \right) \cap \left(\mathbf{a}' + \mathbf{c}' + \sum_{i=1}^{n_3} \lambda_i \mathcal{C}_{H_2}(w_i) \right) \\ &= \left(\mathbf{a} + \mathbf{c} + \sum_{i=1}^{n_3} \lambda_i \mathbf{d}_i + L_1 \right) \cap \left(\mathbf{a}' + \mathbf{c}' + \sum_{i=1}^{n_3} \lambda_i \mathbf{d}'_i + L_2 \right). \end{aligned}$$

Hence, the coset intersection $\left((t^{\mathbf{a}} u)H_1 \cap (t^{\mathbf{a}'} u')H_2 \right) \pi_0$ is empty if and only if for every integer $\lambda_1, \dots, \lambda_{n_3} \in \mathbb{Z}$, $(\mathbf{a} - \mathbf{a}' + \mathbf{c} - \mathbf{c}') + \sum_{i=1}^{n_3} \lambda_i (\mathbf{d}_i - \mathbf{d}'_i) \notin L_1 + L_2$; or equivalently, if and only if

$$\left((\mathbf{a} - \mathbf{a}' + \mathbf{c} - \mathbf{c}') + \langle \mathbf{d}_1 - \mathbf{d}'_1, \dots, \mathbf{d}_{n_3} - \mathbf{d}'_{n_3} \rangle \right) \cap (L_1 + L_2) = \emptyset.$$

This can be effectively decided using linear algebra. And furthermore, in case it is not empty, one can compute an explicit element from $(t^{\mathbf{a}} u)H_1 \cap (t^{\mathbf{a}'} u')H_2$, just following the computations back. This completes the proof. \square

4 The free product case

In this section, we shall consider the free product case and prove Theorems 1.8 and 1.9. We follow the graph-theoretical approach developed by S. Ivanov in [16].

The classical theory of Stallings foldings (see [27]) provides a bijection between the subgroups of a given free group and certain kind of labelled directed graphs (the so-called *Stallings automata*) which when restricted to finitely generated subgroups (corresponding to finite automata) is fully constructive. This geometric approach allowed to solve many algorithmic problems about free groups in a very nice and intuitive way.

S. Ivanov (in [16]) generalized this machinery to free products. Among other applications, this allowed him to give a modern proof of the Kurosh Subgroup Theorem, and of B. Baumslag theorem stating that free products of Howson groups are again Howson (see [16–18]). However, he did not consider algorithmic issues in his approach.

More recently, Kapovich–Weidmann–Myasnikov extended further these folding techniques (see [20]) to fundamental groups of graphs of groups. Their method produces the automaton corresponding to a given subgroup, under some conditions on the edge groups (which automatically hold in the case of free products), and leads to the solution of the membership problem in some cases. However, they do not analyze subgroup intersections.

In order to treat algorithmically intersections of subgroups of free groups, we use Ivanov’s approach with the necessary technical adaptations to make it fully algorithmic. In the sake of clarity, we offer here a self-contained exposition.

The idea is to use generalized folding techniques to algorithmically represent any finitely generated subgroup $H \leq G_1 * G_2$ by a finite graph of certain kind, called a *reduced wedge automaton*, denoted by Γ_H . We note that from such an object Γ_H one can already deduce, algorithmically, a Kurosh decomposition for the subgroup H .

Note also that following Ivanov’s argument it is possible to define Γ_H for arbitrary subgroups $H \leq G_1 * G_2$. On the other hand, Ivanov gives in [16] a generalization of the classical “pullback” technique for free groups: given two subgroups $H_1, H_2 \leq G_1 * G_2$, and having at hand corresponding reduced wedge automata Γ_{H_1} and Γ_{H_2} , he describes a reduced wedge automaton $\Gamma_{H_1} \wedge \Gamma_{H_2}$ for $H_1 \cap H_2$, in terms of Γ_{H_1} and Γ_{H_2} . Note a substantial difference with the free situation: $\Gamma_{H_1} \wedge \Gamma_{H_2}$ may very well be an infinite object even with Γ_{H_1} and Γ_{H_2} being

finite (corresponding to the possible non-Howson situation; i.e., $H_1 \cap H_2$ may very well be non finitely generated, even with H_1 and H_2 being finitely generated).

The main argument in the present section is the following: given finite generating sets for H_1 and H_2 , we are able to construct Γ_{H_1} and Γ_{H_2} , and then start constructing $\Gamma_{H_1} \wedge \Gamma_{H_2}$ (even with the possibility of this being infinite). The crucial point is that, in finite time while the construction is running, we are able to either detect that $H_1 \cap H_2$ is not finitely generated, or to complete the construction of $\Gamma_{H_1} \wedge \Gamma_{H_2}$; in the first case we have algorithmically deduced that the intersection $H_1 \cap H_2$ is not finitely generated, and in the second case we have effectively constructed $\Gamma_{H_1} \wedge \Gamma_{H_2}$, from which we shall be able to extract a finite set of generators for $H_1 \cap H_2$. In order to check whether $H_1 \cap H_2$ is not finitely generated the hypothesis of SIP or ESIP in the factor groups G_1 and G_2 will be crucial.

4.1 Wedge automata

We assume the reader familiar with standard Stallings automata (representing subgroups of free groups, say $\mathbb{F}_2 = \langle a \rangle * \langle b \rangle$), as involutive $\{a, b\}$ -automata recognizing exactly the elements in the corresponding subgroup).

To cover the more general situation of $G_1 * G_2$, we need to encode more information into the arcs. A classical a -labelled arc would correspond to what we call here a G_1 -wedge: an arc subdivided in two halves, admitting a (possibly trivial) label from G_1 on each side, and also admitting a (possibly trivial) subgroup $A \leq G_1$ as a label of the middle (special) vertex between the two halves. Doing the same with the b -arcs (and subgroups of G_2) we get an automaton with two types of vertices, *primary* (the original ones), and *secondary* (the new ones). See Figure 4.



Figure 4: Wedges of first and second kind

In these new automata, walks are going to spell *subsets* of $G_1 * G_2$ (instead of *words* in $\{a, b\}^\pm$), by picking *all* the elements from the label of a secondary vertex when traversing it. Allowing, in addition, vertices to have any degree, we get the new notion of wedge automaton.

Definition 4.1. Let G_1, G_2 be two arbitrary groups. A (*wedge*) (G_1, G_2) -*automaton* is a septuple $\Gamma = (V\Gamma, E\Gamma, \iota, \tau, \ell, \ell^{-1}, \odot)$, where:

- (i) $\Gamma = (V\Gamma, E\Gamma, \iota, \tau, \ell, \ell^{-1})$ is an involutive digraph (called *underlying digraph* of Γ) with three disjoint types of vertices, $V\Gamma = V_0\Gamma \sqcup V_1\Gamma \sqcup V_2\Gamma$; namely, *primary* (those in $V_0\Gamma$, denoted by \bullet), *1-secondary* (those in $V_1\Gamma$, denoted by \circ), and *2-secondary* (those in $V_2\Gamma$, denoted by \bigcirc); and with all arcs in $E\Gamma$ joining (in either direction) a primary vertex with a

secondary one; i.e., $E\Gamma = E_1\Gamma \sqcup E_2\Gamma$, where the arcs in E_j are called j -arcs and connect primary vertices with j -secondary vertices (for $j = 1, 2$).

- (ii) ℓ is a twofold *label map*: for $\nu = 1, 2$, $\ell: E_\nu\Gamma \rightarrow G_\nu$, $e \mapsto \ell_e$ compatible with the involution in Γ ; and $\ell: V_\nu\Gamma \rightarrow \text{Sgp}(G_\nu)$, $q \mapsto \ell_q$.
- (iii) \bullet is a distinguished primary vertex called the *basepoint* of Γ .

We say that a wedge automaton Γ is *connected* (resp., *finite*) if the underlying undirected graph is so; note that, by definition, it is always primary-secondary bipartite. We will also say that a wedge automaton Γ is of *finite type* if the underlying digraph is finite, and the subgroups labelling the secondary vertices are all finitely generated. This will always be the situation when we consider computational issues; in this case, the labels of vertices will usually be given by finite sets of generators. We say that a vertex label is *trivial* if $\ell_q = \{1\}$. If not stated otherwise all the wedge automata appearing from this point will be assumed to be finite.

Recall that wedge automata are involutive (as automata): for every arc $e \equiv p \rightarrow q$ reading g , there exists a unique inverse arc $e^{-1} \equiv q \rightarrow p$ reading g^{-1} . Hence, a wedge automata Γ can always be represented by one of its (say positive, denoted by $E^+\Gamma$) arc orientations. Then $E\Gamma = E^\pm\Gamma := E^+\Gamma \sqcup E^-\Gamma$, where $E^-\Gamma$ is the set of inverses of the arcs in $E^+\Gamma$. A *walk* in Γ is a sequence of alternating and successively incident vertices and arcs, starting and ending at primary vertices, $\gamma = p_0(e_1^{-1}q_1e_1')p_1(e_2^{-1}q_2e_2')p_2 \cdots p_{r-1}(e_r^{-1}q_re_r')p_r$, where p_0, \dots, p_r are (not necessarily distinct) primary vertices, q_1, \dots, q_r are (not necessarily distinct) secondary vertices, and for every $k = 1, \dots, r$, all three of e_k, q_k, e_k' are simultaneously of the same type ν_k ($\nu_k = 1, 2$). A ν -*elementary walk* is a walk of length 2 visiting a secondary vertex of ν -type, $\nu = 1, 2$; it is *degenerate* if it consists of two mutually inverse arcs; otherwise it is called *non-degenerate*. Every walk γ decomposes as a product of elementary walks (either degenerate or nondegenerate, and with possible repetitions) in a unique way, corresponding to the brackets in the expression above: this is called the *elementary decomposition* of γ (for convention, we take $r = 0$ when the walk γ is trivial). We say that a walk γ is *alternating* if its elementary decomposition sequence $\gamma_1, \gamma_2, \dots, \gamma_r$ alternates between types 1 and 2.

The *length* of a walk is the number of arcs in the sequence defining it, i.e., twice the number of elementary walks r in its elementary decomposition.

Remark 4.2. Note that γ involves no backtracking if and only if the γ_i 's in its elementary decomposition $\gamma = \gamma_1 \cdot \gamma_2 \cdots \gamma_r$ are all nondegenerate, and there is no backtracking in the consecutive products $\gamma_i \cdot \gamma_{i+1}$.

Definition 4.3. The *label* of a walk $\gamma = p_0(e_1^{-1}q_1e_1')p_1(e_2^{-1}q_2e_2')p_2 \cdots p_{r-1}(e_r^{-1}q_re_r')p_r$, denoted by ℓ_γ , is the subset $\ell_\gamma = (\ell_{e_1}^{-1}\ell_{q_1}\ell_{e_1'}) (\ell_{e_2}^{-1}\ell_{q_2}\ell_{e_2'}) \cdots (\ell_{e_r}^{-1}\ell_{q_r}\ell_{e_r'}) \subseteq G_1 * G_2$. That is, while travelling along γ , when we traverse an arc e , we *pick* its label ℓ_e , and when we traverse a secondary vertex q we take all labels $c \in \ell_q$ (primary vertices have no contribution to ℓ_γ). Picking always the trivial element when visiting a secondary vertex, we obtain the so-called *basic label* of γ , $\ell_\gamma^\bullet = (\ell_{e_1}^{-1}\ell_{e_1'}) (\ell_{e_2}^{-1}\ell_{e_2'}) \cdots (\ell_{e_r}^{-1}\ell_{e_r'}) \in \ell_\gamma \subseteq G_1 * G_2$.

It is clear that $\ell_{\alpha^{-1}} = \ell_\alpha^{-1}$ and $\ell_\alpha \cdot \ell_\beta = \ell_{\alpha\beta}$.

Remark 4.4. Note also that if $\gamma = p_0(e_1^{-1}q_1e_1')p_1(e_2^{-1}q_2e_2')p_2 \cdots p_{r-1}(e_r^{-1}q_re_r')p_r$ is alternating, and we take elements $c_i \in \ell_{q_i}$ such that $\ell_{e_i}^{-1}c_i\ell_{e_i'} \neq 1$ for all $i = 1, \dots, r$, then the brackets in the expression $(\ell_{e_1}^{-1}c_1\ell_{e_1'}) (\ell_{e_2}^{-1}c_2\ell_{e_2'}) \cdots (\ell_{e_r}^{-1}c_r\ell_{e_r'})$ indicate, precisely, the syllable decomposition of the element in $G_1 * G_2$ read by γ ; otherwise, some consecutive pairs of brackets may merge into the same syllable. This remark will be crucial later.

Let Γ be a (G_1, G_2) -automaton and let p, p' be two primary vertices. We define the *coset recognized* by Γ relative to (p, p') to be the set $\langle \Gamma \rangle_{(p, p')} := \cup_{\gamma} \ell_{\gamma}$, where the union runs over all walks in Γ from p to p' . When $p = p'$, then we abbreviate $\langle \Gamma \rangle_p := \langle \Gamma \rangle_{(p, p)}$. Moreover, if $p = p' = \bullet$ then we simply write $\langle \Gamma \rangle = \langle \Gamma \rangle_{\bullet}$, and we call it the *subgroup recognized* by Γ . The lemma below, which is straightforward to prove, justifies this terminology.

Lemma 4.5. *Let Γ be a (G_1, G_2) -wedge automaton, and let $p, p' \in V_0 \Gamma$. Then,*

- (i) $\langle \Gamma \rangle_p$ is a subgroup of $G_1 * G_2$;
- (ii) $\langle \Gamma \rangle_p$ and $\langle \Gamma \rangle_{p'}$ are conjugate to each other; viz. $\langle \Gamma \rangle_{p'} = (\langle \Gamma \rangle_p)^g$, for every $g \in \langle \Gamma \rangle_{(p, p')}$;
- (iii) $\langle \Gamma \rangle_{(p, p')}$ is a right coset of $\langle \Gamma \rangle_{(p, p)}$; viz. $\langle \Gamma \rangle_{(p, p')} = \langle \Gamma \rangle_{(p, p)} \cdot g$, for every $g \in \langle \Gamma \rangle_{(p, p')}$. \square

Proposition 4.6. *For every subgroup $H \leq G_1 * G_2$, there exists a (G_1, G_2) -automaton Γ recognizing H . Furthermore, if H is finitely generated, one such Γ is of finite type and algorithmically constructible from a finite set of generators for H given in normal form.*

Proof. Let $H = \langle W \rangle$, where $W = \{w_1, w_2, \dots\}$ is a set of generators for H . For every non-trivial generator in W , say w , consider its normal form as an element of $G_1 * G_2$, say $w = a_1b_1 \cdots a_sb_s$, with $s \geq 1$, $a_i \in G_1$, $b_i \in G_2$, $a_i \neq 1$ for all $i = 2, \dots, s$, and $b_i \neq 1$ for all $i = 1, \dots, s-1$. Let $\text{Fl}(w)$ denote the (G_1, G_2) -automaton depicted in Fig. 5, and called the *petal automaton* corresponding to w . Clearly, $\langle \text{Fl}(w) \rangle = \langle w \rangle$.

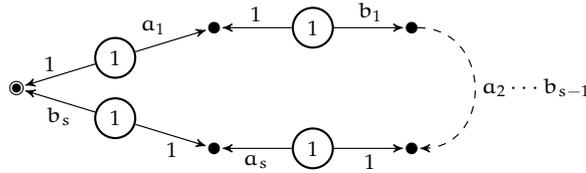


Figure 5: A wedge petal

Now consider $\text{Fl}(W)$ the disjoint union of all the $\text{Fl}(w_i)$'s identifying the basepoints into a single primary vertex (declared as basepoint); the resulting object is a (G_1, G_2) -automaton called the *flower automaton* corresponding to W . Clearly, $\langle \text{Fl}(W) \rangle = \langle W \rangle = H$. Moreover, if $|W| < \infty$ then $\text{Fl}(W)$ is of finite type and constructible. \square

4.2 Reduced wedge automata

In the same vein as in the classical Stallings graphs, we will ask our wedge graphs to be 'deterministic' (in a precise sense specified below). Similar constructions are called

‘irreducible graphs’ by Ivanov in [16], and are particular cases of the so-called ‘folded graphs’ in Kapovich–Weidman–Miasnikov [20].

Definition 4.7. Let G_1, G_2 be two groups, and let Γ be a finite (G_1, G_2) -automaton. We say that Γ is *reduced* if the following conditions are satisfied: (i) Γ is connected; (ii) every primary vertex of Γ is incident with at most one arc from $E_1\Gamma$, and at most one arc from $E_2\Gamma$; (iii) no nondegenerate elementary walk reads the trivial element; that is, for $\nu = 1, 2$, every ν -secondary vertex $q \in V_\nu \Gamma$, and every pair of *different* ν -arcs e_1, e_2 with $\iota e_1 = \iota e_2 = q$, we have that $1 \notin \ell_{e_1}^{-1} \ell_q \ell_{e_2}$ (equivalently, $\ell_{e_1} \ell_{e_2}^{-1} \notin \ell_q$).

In a reduced wedge automaton Γ , remark 4.2 can be restated in the following way: γ presents no backtracking if and only if it is alternating and the elementary walks on its elementary decomposition are all nondegenerate. In this case, additionally, property (iii) from Definition 4.7 ensures that the elementary decomposition of γ gives the syllable decomposition of every $g \in \ell_\gamma$ as element from $G_1 * G_2$ (since nondegenerate elementary walks do not admit the trivial element as a label). However, this is not the whole story: even with some of the γ_i ’s being degenerate, we can still get the syllable decomposition of $g \in \ell_\gamma$ assuming that the elements picked from the labels of the backtracking vertices (if any) are nontrivial. This motivates the following lemma, definition and the subsequent important technical lemma.

Lemma 4.8. For any reduced (G_1, G_2) -automaton Γ , and any nontrivial walk γ between two primary vertices and without backtracking, $1 \notin \ell_\gamma$.

Proof. This is clear since, by the previous remark, the label of γ gives its syllable decomposition as element of $G_1 * G_2$. \square

Definition 4.9. Let γ be a walk in a wedge automaton Γ , with elementary decomposition $\gamma = \gamma_1 \gamma_2 \cdots \gamma_r = p_0(e_1^{-1} q_1 e_1') p_1(e_2^{-1} q_2 e_2') p_2 \cdots p_{r-1}(e_r^{-1} q_r e_r') p_r$. We define the *reduced label* of γ as

$$\tilde{\ell}_\gamma = \left\{ (\ell_{e_1}^{-1} c_1 \ell_{e_1'}) (\ell_{e_2}^{-1} c_2 \ell_{e_2'}) \cdots (\ell_{e_r}^{-1} c_r \ell_{e_r'}) \mid \begin{array}{l} c_i \in \ell_{q_i} \\ c_i \neq 1 \text{ if } \gamma_i \text{ is degenerate} \end{array} \right\} \subseteq \ell_\gamma.$$

Lemma 4.10. For a reduced (G_1, G_2) -automaton Γ , we have $\langle \Gamma \rangle = \cup_\gamma \ell_\gamma = \cup_{\hat{\gamma}} \tilde{\ell}_{\hat{\gamma}}$, where the first union runs over all \bullet -walks γ of Γ , and the second one only over the alternating \bullet -walks $\hat{\gamma}$ of Γ .

Proof. The inclusion ‘ \supseteq ’ is clear, since the first union is over more sets than the second one, and $\ell_\gamma \supseteq \tilde{\ell}_\gamma$. To see ‘ \subseteq ’, fix a \bullet -walk $\gamma = \gamma_1 \cdots \gamma_r$, and an element $g \in \ell_\gamma$, and let us find an alternating \bullet -walk $\hat{\gamma}$ such that $g \in \tilde{\ell}_{\hat{\gamma}}$.

In fact, if γ is not alternating then there are $\nu = 1, 2$ and $i = 1, \dots, r-1$ such that $\gamma_i = p_{i-1} e_i^{-1} q_i e_i' p_i$ and $\gamma_{i+1} = p_i e_{i+1}^{-1} q_{i+1} e_{i+1}' p_{i+1}$ are both of type ν ; so, by condition (ii) in Definition 4.7, $q_i = q_{i+1}$ and $e_i' = e_{i+1}$. Replacing $\gamma_i \gamma_{i+1}$ by $= p_{i-1} e_i^{-1} q_i e_{i+1}' p_{i+1}$, we get a new \bullet -walk $\gamma_{(1)}$, with shorter elementary decomposition and such that $g \in \ell_{\gamma_{(1)}}$ as well, since $(\ell_{e_i}^{-1} c_i \ell_{e_i'}) (\ell_{e_{i+1}}^{-1} c_{i+1} \ell_{e_{i+1}'}) = \ell_{e_i}^{-1} (c_i c_{i+1}) \ell_{e_{i+1}'}$, for all $c_i, c_{i+1} \in \ell_{q_i} \leq G_\nu$. Repeating this

operation a finite number of times, say k , we can assume that $\hat{\gamma} = \gamma_{(k)}$ is alternating and $g \in \ell_{\hat{\gamma}}$.

It remains to prove that, maybe simplifying $\hat{\gamma}$, $g \in \tilde{\ell}_{\hat{\gamma}}$: if q_j is the secondary vertex in the degenerate elementary walk $\gamma_j = p_{j-1}e_j^{-1}q_je_jp_{j-1}$, and the corresponding c_j picked in the formation of g is trivial, then just ignore $(\ell_{e_j}^{-1}c_j\ell_{e_j}) = 1$, and realize g in the label of $\hat{\gamma}_{(1)} = \gamma_1 \cdots \gamma_{j-1}\gamma_{j+1} \cdots \gamma_s$, a \bullet -walk with shorter elementary decomposition, which will be again alternating after repeating the operation in the above paragraph. Repeating this operations a finite number of times, until no trivial choices are made at the degenerate vertices, we obtain the desired result. \square

Remark 4.11. The usefulness of the previous lemma is the following: when realizing an element $g \in \langle \Gamma \rangle$ from the subgroup recognized by a reduced automaton X as $g \in \tilde{\ell}_{\hat{\gamma}}$ for some alternating \bullet -walk $\hat{\gamma}$, the elementary decomposition of $\hat{\gamma}$ automatically provides the syllable decomposition of g as element of $G_1 * G_2$. This is a crucial bridge between the algebraic and the geometric aspects of the theory.

One of the most useful applications of reduced (G_1, G_2) -automata is that they naturally encode the Kurosh free product decomposition (the induced splitting) of their recognized subgroups as subgroups of $G_1 * G_2$. With some technical differences, our exposition follows [16, 20], but with special emphasis on the algorithmic point of view. The theorem below appears as Lemma 4 in Ivanov's [18] and, in a more general setting, as Proposition 4.3 in Kapovich–Weidmann–Myasnikov [20]. A detailed proof (including the algorithmicity in the the finite type case) can be found in [6].

Notation 4.12. Let Γ be a reduced (G_1, G_2) -wedge automaton. Fix a maximal subtree \mathbf{T} in Γ^+ , and let $E = E\Gamma^+ \setminus E\mathbf{T}$ be the set of arcs of Γ^+ outside \mathbf{T} . For every two vertices $u, v \in V\Gamma$, let $\mathbf{T}[u, v]$ denote the unique walk without backtracking from u to v along the tree \mathbf{T} . Now, for every vertex $u \in V\Gamma$, let z_u denote the basic label $z_u = \ell_{\mathbf{T}[u, \bullet]}^{\bullet} \in G_1 * G_2$ and, for every arc $e \in E^{\pm}$, let x_e denote the element $x_e = z_{te}^{-1} \ell_e z_{te} \in G_1 * G_2$.

Theorem 4.13. *Let Γ be a reduced (G_1, G_2) -automaton. Then, with the above notations, the subgroup recognized by Γ is*

$$\langle \Gamma \rangle = \mathbb{F} * \left(\bigstar_{q \in V_1\Gamma} z_q^{-1} \ell_q z_q \right) * \left(\bigstar_{q \in V_2\Gamma} z_q^{-1} \ell_q z_q \right), \quad (8)$$

where \mathbb{F} is the free subgroup of $\langle \Gamma \rangle$ freely generated by the set $\{x_e \mid e \in E\}$. Moreover, if Γ is of finite type, then the subgroup $\langle \Gamma \rangle$ is finitely generated, and we can algorithmically compute a Kurosh decomposition like (8) for $\langle \Gamma \rangle \leq G_1 * G_2$. \square

Corollary 4.14. *For a reduced (G_1, G_2) -automaton Γ , the group $\langle \Gamma \rangle$ is finitely generated if and only if Γ is of finite type (i.e., the underlying graph of Γ has finite rank, and all vertex labels of Γ are finitely generated). In this case, a set of generators (in the form of the Kurosh decomposition theorem) for $\langle \Gamma \rangle$ is computable. \square*

Theorem 4.13 and Corollary 4.14 are easily seen to be false if we substitute reduced automata by general wedge automata.

4.3 Effective reduction of wedged automata

Following [16] and [20], the next step is to show that every finitely generated subgroup $H \leq G_1 * G_2$ is the subgroup recognized by some reduced (G_1, G_2) -automaton of finite type. We introduce several elementary operations on wedge automata which will not change their recognized subgroup (to simplify notation, we shall work with positive arcs and assume that everything done to an arc e will also be done accordingly to e^{-1}).

Definition 4.15. Let us consider the following elementary transformations on a (G_1, G_2) -wedge automata:

- (i) *Adjustment*: replacing the label of any arc e leaving a secondary vertex q , by $c \cdot \ell_e$, for any $c \in \ell_q$; see Figure 6.



Figure 6: Adjustment

- (ii) *Conjugation*: (for $\nu = 1, 2$) replacing, given $g \in G_\nu$, the label ℓ_q of a ν -secondary vertex q , by $(\ell_q)^{g^{-1}} = g\ell_q g^{-1}$; and replacing the label ℓ_{e_i} of every arc e_i incident from q , by the respective $g\ell_{e_i}$; see Figure 7.

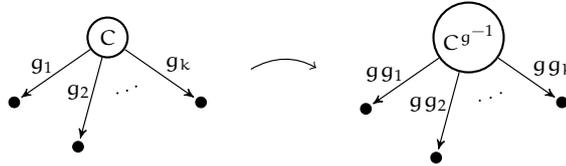


Figure 7: Conjugation

- (iii) *Isolation*: removing from Γ all the connected components not containing the basepoint; see Figure 8.

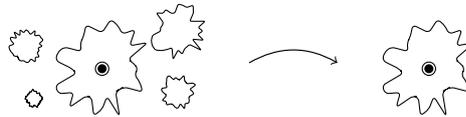


Figure 8: Isolation

- (iv) *Primary open folding*: for $\nu = 1, 2$, given two ν -secondary vertices q_1, q_2 adjacent to the same primary vertex through respective arcs e_1, e_2 with the same label $g \in G_\nu$, identify q_1 and q_2 into a new secondary vertex with label $\langle \ell_{q_1}, \ell_{q_2} \rangle$, and identify the arcs e_1, e_2 into a new arc with the same label g ; see Figure 9.

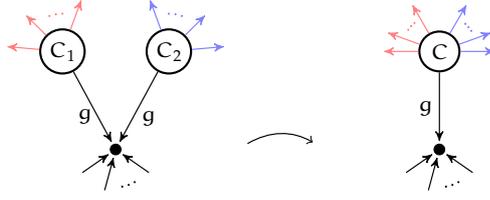


Figure 9: Primary open folding, where $C = \langle C_1 \cup C_2 \rangle$

- (v) *Secondary open folding*: for $\nu = 1, 2$, given a ν -secondary vertex q adjacent to two different primary vertices p_1, p_2 through arcs e_1, e_2 having the same label $g \in G_\nu$, identify the vertices p_1 and p_2 , and the arcs e_1, e_2 into an arc maintaining the label g ; see Figure 10.

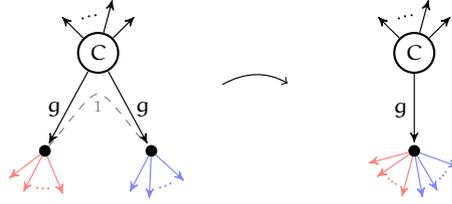


Figure 10: Secondary open folding

- (vi) *Closed folding*: for $\nu = 1, 2$, given a primary vertex p adjacent to a ν -secondary vertex q by two non mutually inverse arcs e_1, e_2 , consists in identifying e_1 and e_2 into a single arc with label ℓ_{e_1} , and change the label of q from ℓ_q to $\langle \ell_q, \ell_{e_1} \ell_{e_2}^{-1} \rangle \leq G_\nu$; see Figure 11.

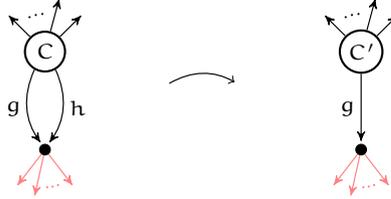


Figure 11: Closed folding, where $C' = \langle C, gh^{-1} \rangle$

Note that the (folding) transformations (iv), (v), and (vi) in Definition 4.15 decrease the number of arcs in the automata exactly by 1. The following result is also straightforward.

Lemma 4.16. *Let Γ be a reduced (G_1, G_2) -automaton, and let $\Gamma \rightsquigarrow \Gamma'$ be any of the elementary transformations in Definition 4.15. Then, the recognized subgroups of Γ and Γ' coincide. \square*

Theorem 4.17 (S. V. Ivanov, 1999, [16]). *For any groups G_1, G_2 , and any finitely generated subgroup $H \leq G_1 * G_2$, there exists a reduced (G_1, G_2) -automaton recognizing H . Moreover, if both G_1 and G_2 have solvable membership problem, then given a finite set of generators of a subgroup $H \leq G_1 * G_2$, one can algorithmically obtain a reduced (G_1, G_2) -automaton of finite type recognizing H .*

Proof. Since the trivial automaton recognizes the trivial subgroup, we can assume $H \neq 1$. Starting with the corresponding flower automaton for H , successively do the following:

- (I) if, for $\nu = 1, 2$, Γ has a primary vertex adjacent to two different ν -secondary vertices, then apply a suitable conjugation move to one of the ν -vertices and then a primary open folding.
- (II) if, for $\nu = 1, 2$, Γ has a primary vertex and a ν -secondary vertex connected to each other by two non mutually inverse arcs then apply a closed folding.
- (III) if Γ has a secondary vertex q adjacent to two different primary vertices through arcs e_1 and e_2 , $\iota e_1 = \iota e_2 = q$, such that $\ell_{e_1} \ell_{e_2}^{-1} \in \ell_q$, apply a suitable adjustment move to one of the two arcs, and then an elementary open folding.

Note that this can be done algorithmically (we use MP in G_1 and G_2 for (III)). Since the number of edges decreases at each step, this process will eventually stop. It is clear that the resulting automaton is reduced, of finite type, and it recognizes H by Lemma 4.16. \square

4.4 A reduced automaton for the intersection

Recall that if G_1 and G_2 are Howson, then $G_1 * G_2$ is Howson as well; see [3] and [16], the second proof being essentially the one we present here.

The goal of this section is to describe, following [16], a reduced automaton $\Gamma_H \wedge \Gamma_K$ recognizing the intersection $H \cap K$, in terms of given reduced automata Γ_H and Γ_K , with $\langle \Gamma_H \rangle = H$ and $\langle \Gamma_K \rangle = K$, where $H, K \leq G_1 * G_2$ are finitely generated subgroups. This construction is not algorithmic, in general, since $\Gamma_H \wedge \Gamma_K$ may very well be not of finite type, even when both Γ_H and Γ_K are so (corresponding to the case where H, K are both finitely generated but $H \cap K$ is not, see Corollary 4.14).

Later, in Section 4.6, we shall give an effective procedure which starts constructing, locally, the aforementioned automaton $\Gamma_H \wedge \Gamma_K$. While running, there will be an alert observing the construction: if some (algorithmically checkable) specific situation occurs, then the intersection $H \cap K$ is not finitely generated. We shall achieve our goal by proving that, in finite time, either the alert sounds or the procedure terminates providing the finite type reduced automaton $\Gamma_H \wedge \Gamma_K$ as output.

So, suppose we are given two finite type reduced (G_1, G_2) -wedge automata Γ_H and Γ_K with recognized subgroups $\langle \Gamma_H \rangle = H$ and $\langle \Gamma_K \rangle = K$, where $H, K \leq G_1 * G_2$ are finitely generated subgroups. Along the following paragraphs, we will first define the *product automaton* $\tilde{\Gamma}$ of Γ_H and Γ_K whose main connected component will be the junction $\Gamma_H \wedge \Gamma_K$.

Define the set of *primary vertices of the product* $\tilde{\Gamma}$ as the cartesian product $V_0 \tilde{\Gamma} = V_0 \Gamma_H \times V_0 \Gamma_K$. and the basepoint \odot of $\tilde{\Gamma}$ to be the pair of basepoints, i.e., $\odot = (\odot_H, \odot_K)$. Now, for $\nu = 1, 2$, we consider the subsets of primary vertices (in Γ_H and Γ_K) adjacent to some ν -secondary vertex, i.e., $V_{0 \leftarrow \nu} \Gamma_H = \{p \in V_0 \Gamma_H \mid p \text{ is adj. to a } \nu\text{-secondary in } \Gamma_H\} \subseteq V_0 \Gamma_H$ (and idem for $V_{0 \leftarrow \nu} \Gamma_K$) and define the relation \equiv_ν on the set $V_{0 \leftarrow \nu} \Gamma_H \times V_{0 \leftarrow \nu} \Gamma_K \subseteq V_0 \tilde{\Gamma}$ to be: $(p_1, p'_1) \equiv_\nu (p_2, p'_2)$ if

and only if there exist two ν -elementary walks: γ from p_1 to p_2 in Γ_H (say $\gamma = p_1 e_1^{-1} q e_2 p_2$, with $q \in V_\nu \Gamma_H$), and γ' from p'_1 to p'_2 in Γ_K (say $\gamma' = p'_1 e_1'^{-1} q' e_2' p'_2$, with $q' \in V_\nu \Gamma_K$), such that the intersection of their labels is nonempty, $l_\gamma \cap l_{\gamma'} \neq \emptyset$.

The following result is not hard to see with the natural reasoning.

Lemma 4.18. *The relations \equiv_1 and \equiv_2 are equivalence relations.* \square

Now, define the ν -secondary vertices of $\tilde{\Gamma}$ to be the equivalence classes modulo \equiv_ν , namely $V_\nu \tilde{\Gamma} = (V_{0 \leftarrow \nu} \Gamma_H \times V_{0 \leftarrow \nu} \Gamma_K) / \equiv_\nu$. Finally, define the ν -arcs of $\tilde{\Gamma}$ as $E_\nu \tilde{\Gamma} = \{\mathbf{q} \rightarrow \mathbf{p} \mid \mathbf{q} \in V_\nu \tilde{\Gamma}, \mathbf{p} \in \mathbf{q}\}$, i.e., for each secondary vertex $\mathbf{q} \in V_\nu \tilde{\Gamma}$, and each primary vertex $\mathbf{p} = (p, p') \in \mathbf{q}$, add a ν -arc $\mathbf{q} \rightarrow \mathbf{p}$. This finishes the definition of the underlying digraph of $\tilde{\Gamma}$. Observe that $\tilde{\Gamma}$ may not be connected in general, even with Γ_H and Γ_K being so.

By construction, it is clear that, for $\nu = 1, 2$, every primary vertex of $\tilde{\Gamma}$ is adjacent to at most one ν -secondary vertex of $\tilde{\Gamma}$ through at most one arc. We define now natural projections (digraph homomorphisms) from $\tilde{\Gamma}$ to Γ_H and Γ_K . Define $\pi: \tilde{\Gamma} \rightarrow \Gamma_H$ as follows: for primary vertices, take the projection to the first coordinate, $\pi: V_0 \tilde{\Gamma} \rightarrow V_0 \Gamma_H, (p, p') \mapsto p$; for ν -secondary vertices, assign to every vertex $\mathbf{q} \in V_\nu \tilde{\Gamma}$ the only ν -secondary vertex in Γ_H adjacent to every $p = (p, p') \pi \in V_0 \Gamma_H$ for $(p, p') \in \mathbf{q}$; finally, for ν -arcs, assign to every $\mathbf{q} \rightarrow \mathbf{p}$ (for $\mathbf{q} \in V_\nu \tilde{\Gamma}$ and $(p, p') \in \mathbf{q}$) the unique ν -arc in Γ_H from $q = \mathbf{q}\pi$ to $p = (p, p')\pi$. Clearly, π is a well defined digraph homomorphism, called the *projection* to Γ_H . The projection to Γ_K , denoted by $\pi': \tilde{\Gamma} \rightarrow \Gamma_K$, is defined in an analogous way.

It remains to establish the labels for its vertices and arcs. For $\nu = 1, 2$, and for every ν -secondary vertex $\mathbf{q} \in V_\nu \tilde{\Gamma}$, choose a distinguished primary vertex $\mathbf{p}_q = (p_q, p'_q) \in \mathbf{q}$, and let \mathbf{e}_q be the only arc in $\tilde{\Gamma}$ from \mathbf{q} to the representative \mathbf{p}_q . This means that, in Γ_H , there is a ν -arc $\mathbf{e}_q := (\mathbf{e}_q)\pi \in E_\nu \Gamma_H$ from $q = \mathbf{q}\pi \in V_\nu \Gamma_H$ to p_q , and, in Γ_K , there is a ν -arc $\mathbf{e}'_q := (\mathbf{e}_q)\pi' \in E_\nu \Gamma_K$ from $q' = \mathbf{q}\pi' \in V_\nu \Gamma_K$ to p'_q . Then, we define the label of vertex \mathbf{q} as

$$l_{\mathbf{q}} := l_{\mathbf{e}_q}^{-1} l_q l_{\mathbf{e}_q} \cap l_{\mathbf{e}'_q}^{-1} l_{q'} l_{\mathbf{e}'_q} = l_q^{l_{\mathbf{e}_q}} \cap l_{q'}^{l_{\mathbf{e}'_q}} \leq G_\nu. \quad (9)$$

Finally, for any ν -arc $\mathbf{e} \in E_\nu \tilde{\Gamma}$ from \mathbf{q} to a primary vertex $\mathbf{p} = (p, p') \in \mathbf{q}$, call $e := \mathbf{e}\pi$, $e' := \mathbf{e}\pi'$, and define the label of \mathbf{e} as an arbitrary element from the coset intersection $l_{\mathbf{e}} \in l_{\mathbf{e}_q}^{-1} l_q l_{\mathbf{e}} \cap l_{\mathbf{e}'_q}^{-1} l_{q'} l_{\mathbf{e}'}$, which is nonempty since $(p_q, p'_q) \equiv_\nu (p, p')$ by construction; see Figure 12. Note that, in particular, we can take $l_{\mathbf{e}_q} = 1$.

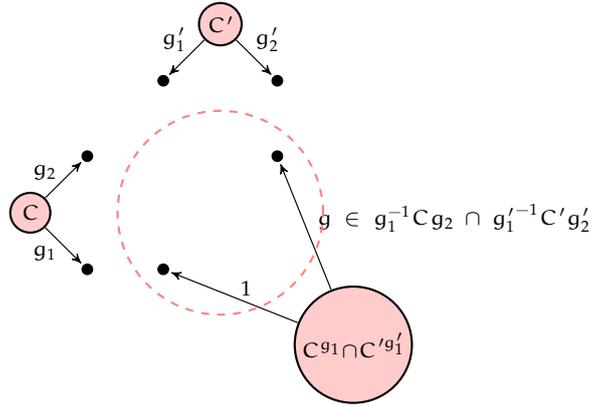


Figure 12: Secondary vertex and related labels in $\tilde{\Gamma}$

This completes the definition of the product automaton $\tilde{\Gamma}$ (technically depending on some choices made on the way). The main property of the labels defined for $\tilde{\Gamma}$ is expressed in the following lemma.

Lemma 4.19. *Let $\tilde{\Gamma}$ be the product of two reduced wedge automata Γ_H and Γ_K . Then, for every ν -elementary walk γ in $\tilde{\Gamma}$, the projected walks $\gamma\pi$ and $\gamma\pi'$ are ν -elementary in Γ_H and Γ_K , respectively, and we have $l_\gamma = l_{\gamma\pi} \cap l_{\gamma\pi'}$. Furthermore, γ is degenerate if and only if both $\gamma\pi$ and $\gamma\pi'$ are degenerate.*

Proof. Clearly, the projections by π and π' of ν -elementary walks in $\tilde{\Gamma}$ are also ν -elementary walks (in Γ_H and Γ_K , respectively), with the original one being degenerate if and only if both projections are degenerate (note that γ may be nondegenerate with one (and only one) of $\gamma\pi$ and $\gamma\pi'$ being degenerate; see Figure 12).

To see the equality in labels, let $\gamma = \mathbf{p}_1 \mathbf{e}_1^{-1} \mathbf{q} \mathbf{e}_2 \mathbf{p}_2$ be a ν -elementary walk in $\tilde{\Gamma}$, where $\mathbf{p}_1 = (p_1, p_1')$ and $\mathbf{p}_2 = (p_2, p_2')$, and let $\gamma\pi = p_1 e_1^{-1} q e_2 p_2$ and $\gamma\pi' = p_1' (e_1')^{-1} q' e_2' p_2'$ be the corresponding ν -elementary walks in Γ_H and Γ_K , respectively. We have $l_{\gamma\pi} = l_{e_1}^{-1} l_q l_{e_2}$ and $l_{\gamma\pi'} = l_{e_1'}^{-1} l_{q'} l_{e_2'}$.

Now, consider the distinguished ν -arc \mathbf{e}_q incident to \mathbf{q} (possibly equal to \mathbf{e}_1 and/or \mathbf{e}_2). According to the above definitions, we have $l_q = l_q^{e_q} \cap l_{q'}^{e_q'}$ and $l_{e_i} \in l_{e_q}^{-1} l_q l_{e_i} \cap l_{e_q'}^{-1} l_{q'} l_{e_i'}$, for $i = 1, 2$; see Figure 12. Therefore,

$$l_\gamma = l_{e_1}^{-1} l_q l_{e_2} = l_{e_1}^{-1} \cdot \left(l_q^{e_q} \cap l_{q'}^{e_q'} \right) \cdot l_{e_2} = l_{e_1}^{-1} l_q l_{e_2} \cap l_{e_1'}^{-1} l_{q'} l_{e_2'} = l_{\gamma\pi} \cap l_{\gamma\pi'}. \quad (10)$$

This completes the proof. \square

We can now state the definition of junction automaton $\Gamma_H \wedge \Gamma_K$ and show that it is a reduced (G_1, G_2) -wedge automaton such that $\langle \Gamma_H \wedge \Gamma_K \rangle = H \cap K$.

Definition 4.20. With the above notation, we define the *junction automaton* of Γ_H and Γ_K , denoted by $\Gamma_H \wedge \Gamma_K$, as the connected component of $\tilde{\Gamma}$ containing the basepoint. (Recall that

there are some arbitrary choices made on the way so, $\Gamma_H \wedge \Gamma_K$ is not canonically associated to Γ_H and Γ_K .)

Proposition 4.21. *The junction automaton $\Gamma_H \wedge \Gamma_K$ is a (G_1, G_2) -reduced automaton recognizing the subgroup $H \cap K$.*

Proof. It is clear that the junction automaton satisfies the properties (i) and (ii) in Definition 4.7.

To see property (iii), take $\nu = 1, 2$, let $\mathbf{q} \in V_\nu \tilde{\Gamma}$ be a ν -secondary vertex of $\tilde{\Gamma}$, let $\mathbf{e}_1, \mathbf{e}_2$ be two different ν -arcs from \mathbf{q} to $\mathbf{p}_1 = (p_1, p'_1)$ and $\mathbf{p}_2 = (p_2, p'_2)$, respectively, and consider the nondegenerate elementary walk $\gamma = \mathbf{p}\mathbf{e}_1^{-1}\mathbf{q}\mathbf{e}_2\mathbf{p}_2$. By symmetry, we can assume $p_1 \neq p_2$, i.e., that $\gamma\pi = p_1\mathbf{e}_1^{-1}\mathbf{q}\mathbf{e}_2p_2$ is a nondegenerate elementary walk in Γ_H . Since Γ_H is reduced, $1 \notin \ell_{\mathbf{e}_1}^{-1}\ell_{\mathbf{q}}\ell_{\mathbf{e}_2} = \ell_{\gamma\pi}$ and by Lemma 4.19, $1 \notin \ell_\gamma = \ell_{\gamma\pi} \cap \ell_{\gamma\pi'}$.

It remains to show that $\langle \Gamma_H \wedge \Gamma_K \rangle = H \cap K$. Indeed, let γ be an arbitrary \odot -walk in $\tilde{\Gamma}$, and let $\gamma = \gamma_1 \cdots \gamma_r$ be its elementary decomposition. Clearly, the elementary decompositions of $\gamma\pi$ in Γ_H , and $\gamma\pi'$ in Γ_K , are $\gamma\pi = (\gamma_1\pi) \cdots (\gamma_r\pi)$ and $\gamma\pi' = (\gamma_1\pi') \cdots (\gamma_r\pi')$, respectively. Then, by Lemma 4.19, $\ell_\gamma = \ell_{\gamma_1} \cdots \ell_{\gamma_r} \subseteq \ell_{\gamma_1\pi} \cdots \ell_{\gamma_r\pi} = \ell_{(\gamma_1\pi) \cdots (\gamma_r\pi)} = \ell_{\gamma\pi} \subseteq H$. Since this is true for every γ , we deduce $\langle \Gamma_H \wedge \Gamma_K \rangle \leq H$; and, by the symmetric argument, also $\langle \Gamma_H \wedge \Gamma_K \rangle \leq K$.

For the other inclusion, take an element $g \in H \cap K$, and let $g = g_1 \cdots g_r$ be its syllable decomposition in $G_1 * G_2$. Since Γ_H is a reduced automaton and $\langle \Gamma_H \rangle = H$, Lemma 4.10 ensures us that $g \in \tilde{\ell}_\gamma$ for some *alternating* \odot -walk γ from Γ_H ; in this situation, its elementary decomposition, $\gamma = \gamma_1 \cdots \gamma_r$, corresponds to the syllable decomposition $g = g_1 \cdots g_r$, i.e., $g_i \in \tilde{\ell}_{\gamma_i}$, for $i = 1, \dots, r$. Similarly, there exists an *alternating* \odot -walk γ' from Γ_K , whose elementary decomposition $\gamma' = \gamma'_1 \cdots \gamma'_r$ again corresponds to the syllable decomposition $g = g_1 \cdots g_r$, i.e., $g_i \in \tilde{\ell}_{\gamma'_i}$, for $i = 1, \dots, r$.

Write $\gamma_i = \mathbf{p}_{i-1}\mathbf{e}_i^{-1}\mathbf{q}_i\mathbf{f}_i\mathbf{p}_i$ and $\gamma'_i = \mathbf{p}'_{i-1}(\mathbf{e}'_i)^{-1}\mathbf{q}'_i\mathbf{f}'_i\mathbf{p}'_i$. Then, for each $i = 1, \dots, r$, we have $g_i \in \ell_{\gamma_i} = \ell_{\mathbf{e}_i}^{-1}\ell_{\mathbf{q}_i}\ell_{\mathbf{f}_i}$ and $g_i \in \ell_{\gamma'_i} = \ell_{\mathbf{e}'_i}^{-1}\ell_{\mathbf{q}'_i}\ell_{\mathbf{f}'_i}$; so, $\emptyset \neq \ell_{\mathbf{e}_i}^{-1}\ell_{\mathbf{q}_i}\ell_{\mathbf{f}_i} \cap \ell_{\mathbf{e}'_i}^{-1}\ell_{\mathbf{q}'_i}\ell_{\mathbf{f}'_i} \subseteq G_\nu$. This means that $\mathbf{p}_{i-1} = (p_{i-1}, p'_{i-1}) \equiv_{\nu_i} (p_i, p'_i) = \mathbf{p}_i$, where ν_i is the common type of the vertices \mathbf{q}_i (in Γ_H) and \mathbf{q}'_i (in Γ_K). Therefore, \mathbf{p}_{i-1} and \mathbf{p}_i are both incident to a common ν_i -secondary vertex in $\Gamma_H \wedge \Gamma_K$. In other words, there is a ν_i -elementary walk in $\Gamma_H \wedge \Gamma_K$, say γ_i , from \mathbf{p}_{i-1} to \mathbf{p}_i . Finally, by Lemma 4.19, $g_i \in \ell_{\gamma_i} \cap \ell_{\gamma'_i} = \ell_{\gamma_i}$. Therefore, $g = g_1 \cdots g_r \in \ell_{\gamma_1} \cdots \ell_{\gamma_r} = \ell_{\gamma_1 \cdots \gamma_r} \subseteq \langle \Gamma_H \wedge \Gamma_K \rangle$, concluding the proof. \square

Corollary 4.22. *In the above situation, $H \cap K$ is finitely generated if and only if all the vertex labels of $\Gamma_H \wedge \Gamma_K$ are finitely generated.* \square

4.5 Understanding intersections of cosets

According to Lemma 4.5, given a wedge automaton Γ_H , the union of labels of all the walks in Γ_H from the basepoint to a primary vertex $\mathbf{p} \in V_0 \Gamma_H$, denoted by $\langle \Gamma_H \rangle_{(\odot, \mathbf{p})}$, constitutes a coset of the recognized subgroup $\langle \Gamma_H \rangle = H$. In general, though, this does not reflect all the cosets of H (consider, for example the cases when Γ_H has only finitely many primary

vertices, but H has infinite index in $G_1 * G_2$). We can slightly modify the automaton Γ_H to fix this issue.

Let $u = a_1 b_1 \cdots a_s b_s \in G_1 * G_2$, written in normal form. Consider the (G_1, G_2) -wedge automaton (also denoted by u) consisting on a chain spelling the normal form for g , and having trivial vertex label; let us call it the *thread* for u . Attach this thread to Γ_H by identifying the basepoint \bullet_H with τu , and then apply the folding process until no more foldings are possible (see the proof of Theorem 4.17). Observe that operation (II) will not be used, and the triviality of the vertex labels in the thread implies that the vertex groups already present in Γ_H will not be changed along the process. So, the output is the exact same graph Γ_H with a terminal segment of the thread attached somewhere and sticking out; denote this new automaton by Γ_{Hu} . Clearly, Γ_{Hu} is a reduced automaton, like Γ_H , and furthermore, since the new secondary vertices out of Γ_H have trivial label, $\langle \Gamma_{Hu} \rangle = \langle \Gamma_H \rangle = H$.

By Lemma 4.5 (iii), $\langle \Gamma_{Hu} \rangle_{(\bullet_H, \tau u)} = \langle \Gamma_H \rangle \cdot u$ (the situation where this coset could already be represented by a vertex in Γ_H corresponds to the fact that the thread happens to fold completely and so, $\Gamma_{Hu} = \Gamma_H$).

Now let us go back to the graph $\Gamma_H \wedge \Gamma_K$. It is useful to understand the intersection of H and K but also, adding the corresponding hairs, it will be useful to understand the intersection of two arbitrary cosets Hu and Kv .

Given elements $u, v \in G_1 * G_2$, consider the reduced automata Γ_{Hu} and Γ_{Kv} , and consider the junction automaton $\Gamma_{Hu} \wedge \Gamma_{Kv}$.

Lemma 4.23. *With the above notation,*

- (i) $\Gamma_H \wedge \Gamma_K$ is a reduced subgraph of $\Gamma_{Hu} \wedge \Gamma_{Kv}$;
- (ii) $Hu \cap Kv \neq \emptyset$ if and only if the vertex $(\tau u, \tau v)$ belongs to $\Gamma_{Hu} \wedge \Gamma_{Kv}$;
- (iii) for any walk γ in $\Gamma_{Hu} \wedge \Gamma_{Kv}$ from (\bullet_H, \bullet_K) to $(\tau u, \tau v)$, and any $g \in \ell_\gamma$, we have $Hu \cap Kv = \langle \Gamma_{Hu} \wedge \Gamma_{Kv} \rangle_{(\bullet_H, \bullet_K), (\tau u, \tau v)} = (H \cap K)g$.

Proof. Note that the initial set of primary vertices for $\Gamma_{Hu} \wedge \Gamma_{Kv}$, namely $V_0 \Gamma_{Hu} \times V_0 \Gamma_{Kv}$, contains as a subset $V_0 \Gamma_H \times V_0 \Gamma_K$, the initial set of primary vertices for $\Gamma_H \wedge \Gamma_K$. And two old vertices $(p_1, p'_1), (p_2, p'_2) \in \Gamma_H \wedge \Gamma_K$ are \equiv_v -equivalent in $\Gamma_H \wedge \Gamma_K$ if and only if they are \equiv_v -equivalent as vertices in $\Gamma_{Hu} \wedge \Gamma_{Kv}$ (since vertices of $\Gamma_{Hu} \wedge \Gamma_{Kv}$ outside $\Gamma_H \wedge \Gamma_K$ have always trivial labels). This proves (i).

Suppose first that $(\tau u, \tau v) \in V \Gamma_{Hu} \wedge \Gamma_{Kv}$, let γ be a walk in $\Gamma_{Hu} \wedge \Gamma_{Kv}$ from (\bullet_H, \bullet_K) to $(\tau u, \tau v)$, and consider its basic label $\ell_\gamma^\bullet \in G_1 * G_2$. By the same argument as in Proposition 4.21, $\ell_\gamma^\bullet \in \ell_{\gamma\pi} \cap \ell_{\gamma\pi'}$. But $\gamma\pi$ (resp., $\gamma\pi'$) is a walk in Γ_{Hu} from \bullet_H to τu (resp., a walk in Γ_{Kv} from \bullet_K to τv) hence, by Lemma 4.5 (iii), $\ell_\gamma^\bullet \in \langle \Gamma_{Hu} \rangle_{(\bullet_H, \tau u)} \cap \langle \Gamma_{Kv} \rangle_{(\bullet_K, \tau v)} = Hu \cap Kv$, concluding that $Hu \cap Kv \neq \emptyset$.

Conversely, suppose that $Hu \cap Kv \neq \emptyset$ and let $g \in Hu \cap Kv$. Again by Lemma 4.5 (iii), there exist walks γ in Γ_{Hu} from \bullet_H to τu , and γ' in Γ_{Kv} from \bullet_K to τv , such that $g \in \ell_\gamma \cap \ell_{\gamma'}$. Again, with an argument like in the proof of Proposition 4.21, there exists a walk γ in

$\Gamma_{H\mathbf{u}} \wedge \Gamma_{K\mathbf{v}}$ from (\bullet_H, \bullet_K) to $(\tau\mathbf{u}, \tau\mathbf{v})$ such that $g \in \ell_{\mathbf{v}}$. In particular, $(\tau\mathbf{u}, \tau\mathbf{v}) \in \Gamma_{H\mathbf{u}} \wedge \Gamma_{K\mathbf{v}}$. This proves (ii) and (iii). \square

4.6 Proofs of Theorems 1.8 and 1.9

Let us now address the algorithmic aspects of this construction. For all the present section, assume the two starting reduced (G_1, G_2) -automata Γ_H and Γ_K to be of finite type (namely, H and K are finitely generated subgroups of $G_1 * G_2$). Recall that although the underlying graph of $\Gamma_H \wedge \Gamma_K$ is finite, the labels of the vertices in $\Gamma_H \wedge \Gamma_K$ may very well be non finitely generated as a result of intersections of finitely generated subgroups of G_1 and of G_2 .

A first easy observation is that, under the assumption that both G_1 and G_2 are Howson, then $\Gamma_H \wedge \Gamma_K$ will always be of finite type. This proves that the free product of two Howson groups is again Howson, recovering a classical result originally proved by B. Baumslag in [3].

Proof of Theorem 1.8. Assume that both G_1 and G_2 have solvable ESIP; and suppose we are given two finitely generated subgroups $H, K \leq G_1 * G_2$ by finite sets of generators, and two extra elements $\mathbf{u}, \mathbf{v} \in G_1 * G_2$, all of them in normal form. By Remark 2.2, both G_1 and G_2 also have solvable membership problem; and by Theorem 4.17, we can compute reduced (G_1, G_2) -automata Γ_H and Γ_K such that $\langle \Gamma_H \rangle = H$, and $\langle \Gamma_K \rangle = K$.

Now let us keep constructing $\Gamma_H \wedge \Gamma_K$: we start looking at the basepoint $\bullet = (\bullet_H, \bullet_K)$, with the whole set $V_0 \tilde{\Gamma} = V_0 \Gamma_H \times V_0 \Gamma_K$ in the background. We have to keep adding ν -secondary vertices (with their labels), and ν -arcs (with their labels too) connecting them to certain primaries, until getting $\Gamma_H \wedge \Gamma_K$, the full connected component of $\tilde{\Gamma}$ containing the basepoint \bullet .

We start checking whether there exists $\nu \in \{1, 2\}$, such that both \bullet_H and \bullet_K have nonempty ν -neighborhoods. If not, then the basepoint \bullet is not adjacent to any secondary vertex in $\tilde{\Gamma}$, and we are done (namely, the product is the trivial automaton, and $H \cap K = 1$). Otherwise, let $\mathbf{q} \in V_{\nu} \Gamma_H$, and $\mathbf{e} \in E_{\nu} \Gamma_H$ with $\iota\mathbf{e} = \mathbf{q}$, $\tau\mathbf{e} = \bullet_H$; and let $\mathbf{q}' \in V_{\nu} \Gamma_K$ and $\mathbf{e}' \in E_{\nu} \Gamma_K$ with $\iota\mathbf{e}' = \mathbf{q}'$, $\tau\mathbf{e}' = \bullet_K$, and enlarge our picture by drawing a new ν -secondary vertex, say \mathbf{q} , and a new ν -arc, say $\mathbf{e} = (\mathbf{e}, \mathbf{e}')$, from \mathbf{q} to \bullet . According to (9) — and with respect to the choice $(\mathbf{p}_{\mathbf{q}}, \mathbf{p}'_{\mathbf{q}}) = (\bullet_H, \bullet_K)$ — we know that the label of \mathbf{q} is $\ell_{\mathbf{q}} = \ell_{\mathbf{q}}^{\ell_{\mathbf{e}}} \cap \ell_{\mathbf{q}'}^{\ell_{\mathbf{e}'}} \leq G_{\nu}$.

Applying SIP for G_{ν} to the (finitely generated) subgroups $\ell_{\mathbf{q}}^{\ell_{\mathbf{e}}}$ and $\ell_{\mathbf{q}'}^{\ell_{\mathbf{e}'}}$, we can decide whether $\ell_{\mathbf{q}}$ is finitely generated or not. In case it is not, kill the whole process and declare $H \cap K$ to be non finitely generated. Otherwise, compute a finite set of generators for $\ell_{\mathbf{q}}$, assign $\ell_{\mathbf{e}} = 1$, and check which other primary vertices from $\tilde{\Gamma}$ are adjacent to \mathbf{q} : $\mathbf{p} = (\mathbf{p}, \mathbf{p}') \in V_0 \tilde{\Gamma}$ is adjacent to \mathbf{q} if and only if $(\mathbf{p}, \mathbf{p}') \equiv_{\nu} (\mathbf{p}_{\mathbf{q}}, \mathbf{p}'_{\mathbf{q}})$, which happens if and only if there exists $f \in E_{\nu} \Gamma_H$ from \mathbf{q} to \mathbf{p} , and $f' \in E_{\nu} \Gamma_K$ from \mathbf{q}' to \mathbf{p}' , such that $\ell_{\mathbf{e}}^{-1} \ell_{\mathbf{q}} \ell_f \cap \ell_{\mathbf{e}'}^{-1} \ell_{\mathbf{q}'} \ell_{f'} \neq \emptyset$. So, run over every $\mathbf{p} \in V_0 \Gamma_H$ adjacent to \mathbf{q} , and every $\mathbf{p}' \in V_0 \Gamma_K$ adjacent to \mathbf{q}' and, for each such pair, check whether the intersection of (right) cosets

$$\ell_{\mathbf{q}}^{\ell_{\mathbf{e}}} \cdot (\ell_{\mathbf{e}}^{-1} \ell_f) \cap \ell_{\mathbf{q}'}^{\ell_{\mathbf{e}'}} \cdot (\ell_{\mathbf{e}'}^{-1} \ell_{f'}) = \ell_{\mathbf{e}}^{-1} \ell_{\mathbf{q}} \ell_f \cap \ell_{\mathbf{e}'}^{-1} \ell_{\mathbf{q}'} \ell_{f'} \quad (11)$$

is empty or not; this can be done using the above call to ESIP from G_ν , since they are right cosets of $\ell_q^{\ell_e}, \ell_q^{\ell_{e'}} \leq G_\nu$, whose intersection happens to be finitely generated.

In case this intersection is not empty, add a ν -arc, say $\mathbf{f} = (f, f')$, from \mathbf{q} to \mathbf{p} , and $\ell_{(f, f')}$ arbitrarily chosen from that nonempty intersection. After this procedure, we have a complete picture of the 1-elementary and 2-elementary walks in $\Gamma_H \wedge \Gamma_K$ starting at the basepoint \odot .

Now, for every $\nu = 1, 2$, and every primary vertex $\mathbf{p} = (p, p')$ added to the picture and not yet explored, repeat the same process (with \mathbf{p} in place of \odot). Since the underlying graph of $\tilde{\Gamma}$ is finite, this procedure will either find a non finitely generated vertex label, or will finish the complete construction of $\Gamma_H \wedge \Gamma_K$ in finite time. In the first case, we deduce the non finitely generated type of $H \cap K$; in the second case, we can compute generators for $H \cap K$ (in fact, a Kurosh decomposition) applying Theorem 4.13.

Hence, so far, we have solved SIP($G_1 * G_2$). To finish the proof, let us place ourselves in the case where $H \cap K$ is finitely generated (and so, with the junction automaton $\Gamma_H \wedge \Gamma_K$ fully constructed), and let us decide whether the intersection of right cosets $Hu \cap Kv$ is empty or not. We can extend the computation of $\Gamma_H \wedge \Gamma_K$ to that of $\Gamma_{Hu} \wedge \Gamma_{Kv}$; or alternatively construct directly $\Gamma_{Hu} \wedge \Gamma_{Kv}$ from the beginning.

It only remains to check whether the vertex $(\tau u, \tau v)$ appears in $\Gamma_{Hu} \wedge \Gamma_{Kv}$, or not. Namely, using Lemma 4.23 (ii): $(\tau u, \tau v)$ is connected to (\odot_H, \odot_K) if and only if the intersection $Hu \cap Kv$ is nonempty; and, if so, any element g from the label of any walk from (\odot_H, \odot_K) to $(\tau u, \tau v)$ belongs to such intersection, $g \in Hu \cap Kv = (H \cap K)g$. This concludes the proof. \square

Finally, we complement the arguments in the last proof to prove that TIP also passes through free products.

Proof of Theorem 1.9. Since solvability of TIP implies that of ESIP, Theorem 1.8 already gives us ESIP for $G_1 * G_2$. It remains to solve CIP for $G_1 * G_2$ in the case where the given finitely generated subgroups H, K have a non finitely generated intersection.

Given $H, K \leq G_1 * G_2$ finitely generated, and $u, v \in G_1 * G_2$, run the same algorithm as in the proof of Theorem 1.8: construct Γ_{Hu} and Γ_{Kv} and start building the junction $\Gamma_{Hu} \wedge \Gamma_{Kv}$; when we encounter a secondary vertex \mathbf{q} whose label $\ell_{\mathbf{q}} = \ell_q^{\ell_e} \cap \ell_q^{\ell_{e'}} \leq G_\nu$ is not finitely generated, instead of computing a set of generators for it (which is not possible), we just put the trivial subgroup as a label in place of $\ell_{\mathbf{q}}$. Then, when analyzing which other primary vertices are adjacent to \mathbf{q} , we need to decide if the intersection of cosets from equation (11) are empty or not: even though $\ell_q^{\ell_e} \cap \ell_q^{\ell_{e'}} \leq G_\nu$ is not finitely generated, the decision can be made effective using CIP from G_ν . This way, we can algorithmically complete the description of $\Gamma_{Hu} \wedge \Gamma_{Kv}$ except that, for some secondary vertices \mathbf{q} , instead of having generators for $\ell_{\mathbf{q}}$, we just have the trivial element labelling them.

Of course, this is not enough information for computing a set of generators for $H \cap K$. But it suffices for deciding whether the vertices (\odot_H, \odot_K) and $(\tau w, \tau w')$ belong to the same connected component of $\Gamma_{Hu} \wedge \Gamma_{Kv}$. By Lemma 4.23, this allows us to decide whether the intersection of cosets $Hu \cap Kv$ is empty or not; and in case it is not, we can compute an element from it, just choosing a walk γ from (\odot_H, \odot_K) to $(\tau w, \tau w')$, and then picking an

element from ℓ_γ (if γ traverses some secondary vertex with a non finitely generated label, we just recorded the trivial element from it for this purpose). This completes the proof. \square

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