Linearly dependent vectorial decomposition of clutters

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Abstract. This paper deals with the question of completing a monotone increasing family of subsets $\Gamma$ of a finite set $\Omega$ to obtain the linearly dependent subsets of a family of vectors of a vector space. Specifically, we demonstrate that such vectorial completions of the family of subsets $\Gamma$ exist and, in addition, we show that the minimal vectorial completions of the family $\Gamma$ provide a decomposition of the clutter $\Lambda$ of the inclusion-minimal elements of $\Gamma$. The computation of such vectorial decomposition of clutters is also discussed in some cases.

Key words: Clutter, Antichain, Hypergraph, Matroid, Decomposition.

1 Introduction

A monotone increasing family of subsets $\Gamma$ of a finite set $\Omega$ is a collection of subsets of $\Omega$ such that any superset of a set in the family $\Gamma$ must be in $\Gamma$. All the inclusion-minimal elements of $\Gamma$ determine a clutter $\Lambda$, that is, a collection of subsets of $\Omega$ none of which is a proper subset of another. Clutters are also known as antichains, Sperner systems or simple hypergraphs.

A wide variety of examples of monotone increasing families exist, among which we find the collection of the linearly dependent subsets of vectors in a vector space. We say that a clutter $\Lambda$ is LD-vectorial if its elements are the inclusion-minimal linearly dependent subsets of an indexed family of vectors of a vector space. In other words, the LD-vectorial clutters are exactly those corresponding to the set of circuits of representable matroids.

In some cases it is convenient to use clutters that are either LD-vectorial or are closed to be LD-vectorial. Examples of this situation can be found in the context of secret-sharing schemes [3,5], or in the framework of algebraic combinatorics and commutative algebra [1,6]. For instance, in the context of secret-sharing schemes, the LD-vectorial clutters become a crucial issue

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Lemma 1. Let $\Omega$ be a finite set. The following statements hold:
1. If \( \Lambda_1, \Lambda_2 \) are two clutters on \( \Omega \) then, \( \Lambda_1 \preceq \Lambda_2 \) if and only if for all \( A_1 \in \Lambda_1 \) there exists \( A_2 \in \Lambda_2 \) such that \( A_2 \subseteq A_1 \).

2. The binary relation \( \preceq \) is a partial order on the set of clutters of \( \Omega \).

There are many interesting families of clutters that can be considered. However, because of their applications, we are interested in those clutters that are LD-vectorial.

Let \( \Omega = \{x_1, \ldots, x_n\} \) be a finite set of \( n \) elements. A monotone increasing family \( \Gamma \) of subsets of \( \Omega \) is said to be an LD-vectorial family if there exists an indexed family of vectors \( v_1, \ldots, v_n \) of a \( K \)-vector space (that can be \( v_i = v_j \)) such that \( \{x_{i_1}, \ldots, x_{i_r}\} \in \Gamma \) if and only if \( \{v_{i_1}, \ldots, v_{i_r}\} \) is a linearly dependent set of vectors. A clutter \( \Lambda \) on \( \Omega \) is said to be an LD-vectorial clutter if the monotone increasing family \( \Lambda^+ \) is an LD-vectorial family.

In other words, a monotone increasing family of subsets \( \Gamma \) is LD-vectorial if \( \Gamma \) is the family of the dependent sets of a representable matroid \( M \) with ground set \( \Omega \); whereas a clutter \( \Lambda \) is LD-vectorial if the clutter \( \Lambda \) is the set of circuits of a representable matroid \( M \) with ground set \( \Omega \). (The reader is referred to [4,7] for general references on matroid theory). Observe that since the binary relation \( \preceq \) is a partial order on the set of clutters of \( \Omega \), it is also a partial order on the set of LD-vectorial clutters, and therefore \( \preceq \) is a partial order on the set of representable matroids. In matroid theory, this is equivalent to the weak order (see [4, Proposition 7.3.11]).

There are clutters on a finite set \( \Omega \) that are not LD-vectorial (in fact, there are matroids that are not representable matroids). So, a natural question that arises at this point is to determine how to complete a clutter \( \Lambda \) to obtain an LD-vectorial clutter. In order to look for LD-vectorial completions, it is important to take into account the binary relation \( \preceq \) rather than the inclusion \( \subseteq \). This is due to the fact that, as the following example shows, there exist clutters \( \Lambda \) such that \( \Lambda \not\preceq \Lambda' \) for any LD-vectorial clutter \( \Lambda' \).

**Example 1.** Let us consider the clutter \( \Lambda = \{\{1,2\}, \{1,3\}, \{2,3,4\}\} \) on the finite set \( \Omega = \{1,2,3,4\} \). Observe that \( \{\{1,2\} \cup \{1,3\}\} \setminus \{1\} = \{2,3\} \subset \{2,3,4\} \). Hence it follows that \( \Lambda \) is not an LD-vectorial clutter and, moreover, \( \Lambda \not\preceq \Lambda' \) for any LD-vectorial clutter \( \Lambda' \). However, we have that \( \Lambda \preceq \Lambda'' \), where \( \Lambda'' \) is the LD-vectorial clutter \( \Lambda'' = \{\{1\}, \{2,3\}\} \) (an LD-vectorial realization of \( \Lambda' \) is given by the set of vectors \( \{v_1, v_2, v_3, v_4\} \) where \( v_1 = (0,0,0), v_2 = (1,0,0), v_3 = (0,1,0) \) and \( v_4 = (0,0,1) \)). Furthermore, if \( \Lambda''' \) is the clutter on \( \Omega \) defined by \( \Lambda''' = \{\{1,2\}, \{1,3\}, \{2,3\}\} \), then we have that \( \Lambda \preceq \Lambda''' \) and that the clutter \( \Lambda''' \) is also an LD-vectorial clutter (an LD-vectorial realization of \( \Lambda''' \) is given by the set of vectors \( \{w_1, w_2, w_3, w_4\} \) where \( w_1 = (1,1), w_2 = (1,1), w_3 = (1,1) \) and \( w_4 = (0,1) \)). Notice that now the clutter \( \Lambda \) can be obtained from the LD-vectorial clutters \( \Lambda' \) and \( \Lambda'' \). Indeed, it is easy to check that \( \Lambda = \min \{\Lambda' \cup \Lambda'' \} \) where \( \Lambda' \in \Lambda' \) and \( \Lambda'' \in \Lambda'' \). Therefore, the clutters \( \Lambda' \) and \( \Lambda'' \) in some way provide a decomposition of \( \Lambda \).
The above example leads us to the following definition. Let $\Lambda$ be a clutter on a finite set $\Omega$. An \textit{LD-vectorial completion} of the clutter $\Lambda$ is an LD-vectorial clutter $\Lambda'$ on the finite set $\Omega$ such that $\Lambda \leq \Lambda'$.

The set of all the LD-vectorial completions of a clutter $\Lambda$ is denoted by $\text{LD-_vect} (\Lambda)$. Observe that if $\emptyset \in \Lambda$, then $\Lambda = \{\emptyset\}$, and thus $\text{LD-_vect} (\Lambda) = \emptyset$. So, from now on, throughout the paper we assume that $\emptyset \notin \Lambda$ if $\Lambda$ is a clutter. As shown in the next section, this assumption guarantees that $\text{LD-vect} (\Lambda) \neq \emptyset$ for all clutters and, in addition, we demonstrate that suitable clutters in the non-empty set $\text{LD-vect} (\Lambda)$ provide a decomposition of the clutter $\Lambda$ in the same way as in Example 1.

3 Two results on LD-vectorial decompositions

The aim of this section is to present two theoretical results concerning the “decomposition” of a clutter $\Lambda$ into LD-vectorial clutters $\Lambda_1, \ldots, \Lambda_r$ (Theorem 1 and Theorem 2). The general case is considered in the first theorem, while the second deals with those “decompositions” of $\Lambda$ whose LD-vectorial components $\Lambda_1, \ldots, \Lambda_r$ admit vectorial realizations over a fixed field $K$.

Let $\Lambda$ be a clutter on a finite set $\Omega$. Our first result, Theorem 1, states that the set $\text{LD-vect} (\Lambda)$ of its LD-vectorial completions is a non-empty set and that its minimal elements provide a decomposition of $\Lambda$ (in the sense that the elements $A$ of the clutter $\Lambda$ can be obtained from the elements $A_i$ of its minimal LD-vectorial completions $\Lambda_1, \ldots, \Lambda_r$).

**Theorem 1.** Let $\Lambda$ be a clutter on a finite set $\Omega$. Then, $\text{LD-vect} (\Lambda) \neq \emptyset$ and $\Lambda = \min \{A_1 \cup \cdots \cup A_r \mid A_i \in \Lambda_i\}$ where $A_1, \ldots, A_r$ are the minimal elements of the poset $(\text{LD-vect} (\Lambda), \leq)$ of the LD-vectorial completions of $\Lambda$. In particular, the clutter $\Lambda$ has a unique minimal LD-vectorial completion if, and only if, $\Lambda$ is an LD-vectorial clutter.

Observe that the previous theorem, Theorem 1, deals with LD-vectorial completions and decompositions in the case where no field restrictions are assumed. The next theorem, Theorem 2, states that a similar result occurs if we consider only the case in which the vector spaces of the LD-vectorial completions are over a fixed field $K$. Before stating the theorem, we introduce some notations.

Let $\Lambda$ be a clutter on a finite set $\Omega$ and let $K$ be a field. Let us denote by $\text{LD-vect}_K (\Lambda)$ the set whose elements are the LD-vectorial completions of $\Lambda$ over $K$; that is, the elements of $\text{LD-vect}_K (\Lambda)$ are the LD-vectorial clutters $\Lambda'$ over $K$ with $\Lambda \leq \Lambda'$. Therefore, $\text{LD-vect} (\Lambda) = \bigcup_K \text{LD-vect}_K (\Lambda)$.

The next theorem states that the set $\text{LD-vect}_K (\Lambda)$ is a non-empty set and that its minimal elements provides a decomposition of the clutter $\Lambda$ (in the sense that the elements of $\Lambda$ can be obtained from the elements of its minimal LD-vectorial completions over $K$).
Theorem 2. Let \( \Lambda \) be a clutter on a finite set \( \Omega \) and let \( \mathbb{K} \) be a field. Then, \( LD-Vect_{\mathbb{K}}(\Lambda) \neq \emptyset \) and \( \Lambda = \min \{ A_1 \cup \cdots \cup A_r \mid A_i \in A \} \) where \( A_1, \ldots, A_r \) are the minimal elements of the poset \( (LD-Vect_{\mathbb{K}}(\Lambda), \leq) \). In particular, the clutter \( \Lambda \) has a unique minimal LD-vectorial completion over \( \mathbb{K} \) if, and only if, \( \Lambda \) is an LD-vectorial clutter over \( \mathbb{K} \).

4 Computing LD-vectorial decompositions

This section is devoted to the computation of the LD-vectorial decomposition of clutters. First we present a partial result (Proposition 1). After this, several examples are given in order to illustrate this proposition (Example 2 and Example 3). Finally, an example where the proposition cannot be applied is analyzed (Example 4).

Our result, Proposition 1, provides a complete description of the minimal LD-vectorial completion of a clutter \( \Lambda \) on a finite set \( \Omega \) of size at most seven. In order to state this proposition, we need to introduce two transformations of clutters, the \( I \)-transformation and the \( T \)-transformation (see [2]).

Let \( \Lambda \) be a clutter on a finite set \( \Omega \). For a subset \( X \subseteq \Omega \), we denote by \( I_\Lambda(X) \) the intersection of the subsets \( A \) in \( \Lambda \) contained in \( X \), that is, \( I_\Lambda(X) = \bigcap_{A \in \Lambda} A \) where \( A \in \Lambda \) and \( A \subseteq X \), (this intersection is the one involved in the characterization of the set of circuits in connected matroids, see [4, Theorem 4.3.2]). We say that a clutter \( \Lambda' \) is an \( I \)-transformation of the clutter \( \Lambda \) if \( \Lambda' = \min \{ A \cup \{ A_1 \cap A_2 \} \} \) where \( A_1, A_2 \in A \) are two different subsets with \( I_\Lambda(A_1 \cup A_2) \neq \emptyset \).

The definition of \( T \)-transformation is more involved. Let \( \Lambda \) be a clutter. We define the elementary transformations \( T^{(1)}(\Lambda) \) and \( T^{(2)}(\Lambda) \) of the clutter \( \Lambda \) as the clutters \( T^{(1)}(\Lambda) = \min \{ A \cup \{ (A_1 \cup A_2) \setminus \{ x \} \} \mid A_1, A_2 \in A \text{ are different} \} \) and \( T^{(2)}(\Lambda) = \min \{ A \cup \{ (A_1 \cup A_2) \setminus I_\Lambda(A_1 \cup A_2) \mid A_1, A_2 \in A \text{ are different} \} \}. \) Since \( T^{(1)}(\Lambda) \) and \( T^{(2)}(\Lambda) \) are clutters, we can apply the elementary transformations again. Hence, for \( (i_1, i_2) \in \{1, 2\} \times \{1, 2\} \) we can consider the clutter \( T^{(i_2)}(T^{(i_1)}(\Lambda)) \). At this point we proceed in a recursive way. Let \( r \geq 2 \) be a non-negative integer and let \( (i_1, \ldots, i_r) \in \{1, 2\}^r \) be an \( r \)-tuple. Then we define the clutter \( T^{(i_1, \ldots, i_r)}(\Lambda) \) by the recursion formula \( T^{(i_1, \ldots, i_{r-1})}(\Lambda) = T^{(i_r)}(T^{(i_1, \ldots, i_{r-1})}(\Lambda)) \); that is, \( T^{(i_1, \ldots, i_r)}(\Lambda) \) is the \( i_r \) elementary transformation of \( T^{(i_1, \ldots, i_{r-1})}(\Lambda) \). We say that a clutter \( \Lambda' \) is a \( T \)-transformation of the clutter \( \Lambda \) if it is obtained from \( \Lambda \) in this way, that is, if \( \Lambda' = T^{(i_1, \ldots, i_r)}(\Lambda) \) for some \( r \)-tuple \( (i_1, \ldots, i_r) \).

Proposition 1. Let \( \Lambda \) be a non-LD-vectorial clutter on a finite set \( \Omega \) of size \( |\Omega| = n \leq 7 \). Let \( \Lambda' \) be a clutter such that \( \Lambda \leq \Lambda' \). Then the following statements hold:
1. The clutter \( \Lambda' \) is an LD-vectorial completion of \( \Lambda \) if, and only if, \( \Lambda' \) is the unique clutter which can be obtained from \( \Lambda' \) by applying I-transformations or \( T \)-transformations.

2. If the clutter \( \Lambda' \) is a minimal element of the poset \( (LD-Vect(\Lambda), \leq) \), then there is a monotone increasing sequence of clutters \( \Lambda = \Lambda_0 < \Lambda_1 < \ldots < \Lambda_r = \Lambda' \) such that for \( i \geq 1 \), either \( \Lambda_i \) is an I-transformation of \( \Lambda_{i-1} \), or \( \Lambda_i \) is a \( T \)-transformation of \( \Lambda_{i-1} \).

We now give two examples to illustrate the above proposition.

**Example 2.** First, let us consider the clutter \( \Lambda = \{(1,2),\{1,3\},\{2,3,4\}\} \) on the finite set \( \Omega = \{1,2,3,4\} \). In this case only two clutters are obtained by using or by combining I-transformations and \( T \)-transformations; namely, the clutters \( \Lambda_1 = \{(1),\{2,3,4\}\} \) and \( \Lambda_2 = \{(1,2),\{1,3\},\{2,3\}\} \). Therefore, from Proposition 1 it follows that the minimal LD-vectorial completions of the clutter \( \Lambda \) are the minimal elements of \( \{A_1, A_2\} \). In this case, \( A_1 \not\leq A_2 \) and \( A_2 \not\leq A_1 \), and so \( \min(\text{LD-Vect}(\Lambda)) = \{A_1, A_2\} \). Observe that now the LD-vectorial decomposition of \( \Lambda \) given in Example 1 can be stated by applying Theorem 1.

**Example 3.** Now, on the finite set \( \Omega = \{1,2,3,4,5\} \), we consider the clutter \( \Lambda = \{(1,2,3),\{1,2,4\},\{1,5\},\{4,5\}\} \). In such a case, it is a straightforward calculation to check that by using or by combining I-transformations and \( T \)-transformations, eleven clutters \( \Lambda_1, \ldots, \Lambda_{11} \) can be obtained. Specifically, by using only I-transformations we obtain the clutters \( \Lambda_1 = \{\{5\},\{1,2\}\} \) and \( \Lambda_2 = \{\{1\},\{4,5\}\} \). The clutters obtained by using only \( T \)-transformations are the clutters \( \Lambda_3 = \{\{1,4\},\{1,5\},\{4,5\},\{1,2,3\},\{2,3,4\},\{2,3,5\}\} \) and \( \Lambda_4 = \{\{1,3\},\{1,4\},\{1,5\},\{3,4\},\{3,5\},\{4,5\}\} \), whereas the clutters obtained by combining the I-transformations and the \( T \)-transformations are the clutters \( \Lambda_5 = \{\{1\},\{5\}\} \), \( \Lambda_6 = \{\{1\},\{2,4\},\{2,5\},\{4,5\}\} \), \( \Lambda_7 = \{\{5\},\{1,2\},\{1,4\},\{2,4\}\} \), \( \Lambda_8 = \{\{1,2\},\{1,4\},\{1,5\},\{2,4\},\{2,5\},\{4,5\}\} \), \( \Lambda_9 = \{\{5\},\{1,2,3\},\{1,2,4\},\{1,3,4\},\{2,3,4\}\} \), \( \Lambda_{10} = \{\{5\},\{3,4\},\{1,2,3\},\{1,2,4\}\} \), and \( \Lambda_{11} = \{\{4\},\{1,5\},\{1,2,3\},\{2,3,5\}\} \). Therefore, by applying Proposition 1 we obtain that the set of the minimal LD-vectorial completions of the clutter \( \Lambda \) is \( \text{min}(\text{LD-Vect}(\Lambda)) = \{\Lambda_1, \ldots, \Lambda_{11}\} = \{\Lambda_1, \Lambda_2, \Lambda_3, \Lambda_9\} \). So, from Theorem 1 we conclude that \( \Lambda \) admits an LD-vectorial decomposition with four components.

To conclude we give an example where Proposition 1 cannot be applied. In addition, the clutter in this example reveals the difference between Theorem 1 and Theorem 2.

**Example 4.** On the finite set \( \Omega = \{1,2,3,4\} \) of four points, we consider the clutter \( \Lambda = \{\{1,2,3\},\{1,2,4\},\{1,3,4\},\{2,3,4\}\} \). It is easy to prove that the clutter \( \Lambda \) is an LD-vectorial clutter over any field \( \mathbb{K} \neq \mathbb{Z}/(2) \). Therefore, by
applying Theorem 1 and Theorem 2 we obtain that \( \min(\text{LD-Vect}(A)) = \{A\} \); that \( \min(\text{LD-Vect}_K(A)) = \{A\} \) if \( K \neq \mathbb{Z}/(2) \), and that \( \min(\text{LD-Vect}_{\mathbb{Z}/(2)}(A)) \) has at least two elements. In this example, by an extensive exploration it is not hard to show that \( \min(\text{LD-Vect}_{\mathbb{Z}/(2)}(A)) \) has six elements; namely \( \min(\text{LD-Vect}_{\mathbb{Z}/(2)}(A)) = \{A_{1,2}, A_{1,3}, A_{1,4}, A_{2,3}, A_{2,4}, A_{3,4}\} \), where if \( 1 \leq i_1 < i_2 \leq 4 \) and if \( \{i_3, i_4\} = \{1, 2, 3, 4\} \setminus \{i_1, i_2\} \), then \( A_{i_1, i_2} = \{\{i_1, i_2\}, \{i_1, i_3, i_4\}, \{i_2, i_3, i_4\}\} \). Now, from Theorem 2 it follows that the clutter \( A \) admits an LD-vectorial decomposition over \( \mathbb{Z}/(2) \) with six components.

References


